

- What is theory?
- Theory is when you make claims that are either True or False, but not both

Example of claims:

$$1+1 = 2$$

there is a graph with  $> 56$  edges

all prime numbers are between 57 and 59

all regular languages are context-free

- More complicated claims are made up with logical connectives

# Logical connectives

- not A            also written  $\neg A$ ,  $\overline{A}$ ,  $\neg A$ ,  ~~$\bar{A}$~~
- A or B            also written  $A \vee B$ ,  $A \cup B$ , ...
- A and B           also written  $A \wedge B$ ,  $A \& B$ , ...
- A implies B      also written  $A \Rightarrow B$ , if A then B, B if A
  
- You should be familiar with these, but let's clear some doubts

Or

A or B means A or B, possibly both

- Different use in everyday language:  
“We shall triumph or perish”
- Intended meaning is:  
“We shall triumph **exclusive-or** perish”
- Do not confuse or with exclusive or!

# Implication

<i>A</i>	<i>B</i>	<i>A</i> implies <i>B</i>
False	False	True
False	True	True
True	False	False
True	True	True

$A \Rightarrow B$  means if A then B

Only False when A True and B False, True otherwise

“ $1 = 0 \Rightarrow$  the earth is flat” is True (False  $\Rightarrow$  False)

- $A \Rightarrow B$  same as: (not A) or B  
(not B)  $\Rightarrow$  (not A) (contrapositive)

## Different meaning in everyday language:

- “You go out if you finish your homework”

A

B

- Logically means  $B \Rightarrow A$ , can go out and not having finished homework!

- Intended meaning:  $A \Rightarrow B$

“You go out **only if** you finish your homework”

- Do not confuse  $A \Rightarrow B$  with  $B \Rightarrow A$  !

Do you understand implication?



- Know for true: Each card has a number on one side and a letter on the other.
- Suppose I claim: If a card has a vowel on one side, then it has an even number on the other side
- Which cards **must** you turn to know if I lie or not?

## De Morgan's Laws:

$\neg(A \wedge B)$  is equivalent to  $(\neg A) \vee (\neg B)$

$\neg(A \vee B)$  is equivalent to  $(\neg A) \wedge (\neg B)$

There are two quantifiers:

$\exists$

there exists

same thing as OR

$\forall$

for all

same thing as AND

Usually:

OR, AND

few things

$\exists, \forall$

many (infinite) things

Example:

$\exists$  a prime  $x > 5$

same as

6 is prime OR 7 is prime OR 8 is prime OR ...

$\forall x, x < y$

same as

$1 < y$  AND  $2 < y$  AND  $3 < y$  AND ...

## De Morgan's Laws for quantifiers:

$\neg \exists x A(x)$  is equivalent to  $\forall x \neg A(x)$

$\neg \forall x A(x)$  is equivalent to  $\exists x \neg A(x)$

## Sets, Functions:

- Sets are just different notation to express the same claims we construct using logical connectives and quantifiers.

This redundant notation turns out to be useful.

$$(x = 1) \vee (x = 16) \vee (x = 23) \Leftrightarrow x \in \{1, 16, 23\}$$

$$x \text{ is even} \Leftrightarrow x \in \{x \mid x \text{ is even}\}$$

$$A(x) \Leftrightarrow x \in \{x \mid A(x)\}$$

With this in mind, sets become straightforward.

- When are two sets equal?

When the defining claims are equivalent:

$$\{x \mid A(x)\} = \{x \mid B(x)\} \quad \text{same as} \quad A(x) \Leftrightarrow B(x)$$

This shows that order and repetitions do not matter,

for example  $\{b, a, a\} = \{a, b\}$ ,

because  $(x = b) \vee (x = a) \vee (x = a)$  and

$(x = a) \vee (x = b)$  are equivalent claims

- When is a set contained in another?

When its defining claim implies the defining claim of the latter:

$$\{x|A(x)\} \subseteq \{x|B(x)\} \Leftrightarrow A(x) \Rightarrow B(x)$$

$$\{x|A(x)\} \supseteq \{x|B(x)\} \Leftrightarrow B(x) \Rightarrow A(x)$$

$$\{x \mid A(x)\} \cup \{x \mid B(x)\} = \{x \mid A(x) \vee B(x)\}$$

$$\{x \mid A(x)\} \cap \{x \mid B(x)\} = \{x \mid A(x) \wedge B(x)\}$$

$$\overline{\{x \mid A(x)\}} = \{x \mid \neg A(x)\}$$

$$\bigcup_i \{x \mid A_i(x)\} = \{x \mid \exists i A_i(x)\}$$

$$\bigcap_i \{x \mid A_i(x)\} = \{x \mid \forall i A_i(x)\}$$

The empty set is denoted  $\emptyset$

It can be defined as  $\emptyset = \{x : 1+1=3\}$

The empty set is a subset of any set:

$$\emptyset \subseteq \{x : A(x)\} \quad \text{always}$$

because  $1+1=3 \Rightarrow A$  for any  $A$

Powerset(A): Set of all subsets of A.

Example:

$$\text{Powerset}(\{1,2,3\}) = \{\emptyset, \{1\}, \{2\}, \{3\}, \{1,2\}, \{2,3\}, \{1,3\}, \{1,2,3\}\}$$

Size of a set A:  $|A|$  = number of elements in it

Example:  $|\{1,2,3\}| = 3$

Fact:  $|\text{Powerset}(A)| = 2^{|A|}$

Example:  $|\text{Powerset}(\{1,2,3\})| = 2^3 = 8$

Important sets:

$\mathbb{N} = \{0, 1, 2, 3, \dots\}$  Natural numbers

$\mathbb{Z} = \{\dots, -3, -2, -1, 0, 1, 2, 3, \dots\}$  Integer numbers

$\mathbb{R} = \{0, 2.5748954, \pi, \sqrt{2}, -17, \dots\}$  Real numbers

These are all infinite sets: contain an infinite number of elements

A **function**  $f$  from set  $A$  to set  $B$  is written  $f : A \rightarrow B$   
is a way to associate to EVERY element  $a \in A$   
ONE element  $f(a) \in B$

$A$  is called domain,  $B$  range

Example:  $f : \{0, 1\} \rightarrow \{a, b, c\}$  defined as  $f(0)=a$ ,  $f(1)=c$

$f : \mathbb{N} \rightarrow \mathbb{N}$  defined as  $f(n) = n+1$

$f : \mathbb{Z} \rightarrow \mathbb{Z}$  defined as  $f(n) = n^2$

Some  $b \in B$  may not be 'touched,' but every  $a \in A$  must be

Tuples: **Ordered** sequences of elements

Example: (5, 2)	2-tuple, or pair
(7, 8, -1)	3-tuple, or triple
( $\emptyset$ , {4,5}, 8, 21)	4-tuple

Order matters:  $(a, b) \neq (b, a)$

By contrast,  $\{a, b\} = \{b, a\}$

# Construct tuples from sets via Cartesian product

$$\begin{aligned} A \times B &= \text{set of pairs } (a, b) : a \in A \text{ and } b \in B \\ &= \{(a, b) : a \in A \text{ and } b \in B\} \end{aligned}$$

$$A \times B \times C = \{(a, b, c) : a \in A \text{ and } b \in B \text{ and } c \in C\}$$

$$A^k = A \times A \times \dots \times A \quad (k \text{ times})$$

## Example

$$\{q, r, s\} \times \{0, 1\} = \{(q, 0), (q, 1), (r, 0), (r, 1), (s, 0), (s, 1)\}$$

$$\begin{aligned} \{a, b\}^3 &= \{(a, a, a), (a, a, b), (a, b, a), (a, b, b), \\ &\quad (b, a, a), (b, a, b), (b, b, a), (b, b, b)\} \end{aligned}$$

# Strings

Strings are like tuples,  
but written without brackets and commas

Example: (h, e, l, l, o) is written as hello  
(0, 1, 0) is written as 010

# Strings

An **alphabet**  $\Sigma$  is a finite, non-empty set.

We call its elements symbols.

Example:  $\Sigma = \{0,1\}$  (the binary alphabet)

$\Sigma = \{a,b,\dots, z\}$  (English language alphabet)

A **string** over an alphabet  $\Sigma$  is a finite, ordered sequence of symbols from  $\Sigma$

Example: 010101000 a string over  $\Sigma = \{0,1\}$

hello a string over  $\Sigma = \{a,b,\dots, z\}$

A string  $w$  is a **substring** of a string  $x$  if the symbols in  $w$  appears consecutively in  $x$

Example: **aba** is a substring of `aaabbaaaabbbb`  
**00** is a substring of `111100010010100`

The **length** of a string  $w$  is the number of symbols in it

Length is denoted  $|w|$

Example:  $|\text{hello}| = 5$        $|001| = 3$

We denote by  $\Sigma^i$  the set of strings of length  $i$

Example:  $\{0,1\}^2 = \{00, 01, 10, 11\}$

$\text{hello} \in \{a,b,\dots, z\}^5$

$001 \in \{0,1\}^3$

The **empty string** is denoted  $\varepsilon$       (never in  $\Sigma$ )

Its length is 0:  $|\varepsilon| = 0$

We denote by  $\Sigma^*$  the set of all strings over  $\Sigma$  of any length, including  $\varepsilon$

Example:  $\{0, 1\}^* = \{\varepsilon, 0, 1, 001, 10101010, \dots\}$   
= all binary strings

$\{a\}^* = \{\varepsilon, a, aa, aaa, aaaa, \dots\}$   
= all strings containing only a

$\emptyset^* = \{\varepsilon\}$

Note:  $\Sigma^* = \{\varepsilon\} \cup (\cup_i \Sigma^i) = \{\varepsilon\} \cup \Sigma^1 \cup \Sigma^2 \cup \Sigma^3 \cup \dots$

$\Sigma^*$  is an infinite set

A **language** over  $\Sigma$  is a set of strings over  $\Sigma$

Example:  $\{ w : w \in \{0,1\}^* \text{ and ends with } 1 \}$   
 $= \{ 01, 1, 11111, 010101011, \dots \}$

$\{ w : w \in \{a,b, \dots, z\}^* \text{ and } |w| > 3 \}$   
 $= \{ aaaa, abab, zytr, \dots \}$

What are we going to compute?

In this class we ask how to compute functions

$$f : \Sigma^* \rightarrow \{\text{accept, reject}\}$$

We do this for:

- Simplicity
- Sufficient to study fundamental questions, such as:  
are there functions that computers cannot compute?