# Machine Learning Homework Solution 

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## Question:

$X$ and $Y$ are two sets containing finite number of points. $C_{1}$ and $C_{2}$ are convex hulls of $X$ and $Y$, respectively. Prove that if $C_{1}$ and $C_{2}$ don't intersect, there must exist a hyper plane that separates $C_{1}$ and $C_{2}$.

## Answer:

First we need a few lemmas:
Lemma 1: In a triangle ABC (Figure 1), if $\angle A>\angle B$, then $d(B, C)>d(A, C)$, where $d$ is the Euclidean distance.


Figure 1: Lemma 1
Proof of the Lemma 1: $\frac{d(B, C)}{d(A, C)}=\frac{\sin (\angle A)}{\sin (\angle B)}$.
Lemma 2: A convex hull generated by a finite number of points is a closed set.
Proof of the Lemma 2: Omitted.
Proof the the main results:
Since $C_{1}$ and $C_{2}$ don't intersect and they are closed sets, there exist $A \in C_{1}$, $B \in C_{2}$ such that
$d(A, B)=\min _{M \in C_{1}, N \in C_{2}} d(M, N)=\inf _{M \in C_{1}, N \in C_{2}} d(M, N)=d\left(C_{1}, C_{2}\right)>0$

Define the hyper plane $P$ to be the unique plane which contains the middle point of $A$ and $B$ and is perpendicular to the line $A B$. It's easy to verify that for any point $G$ on the plane $P, d(G, A)=d(G, B)$. For any point $G$ on the $A$ side of the plane, $d(G, A)<d(G, B)$. For any point $G$ on the $B$ side of the plane, $d(G, A)>d(G, B)$. We now prove that the hyper plane $P$ defined this way separates $C_{1}$ and $C_{2}$. It suffices to show that $\forall C \in C_{2}, C$ must lie on the $B$ side of the the plane. The other direction of the proof can be shown similarly.

Assume to the contrary that $C$ lies on the $A$ side of the plane or on the plane. Then $d(A, C) \leq d(B, C)$. There are two major cases:

- case 1: $A, B, C$ are in a line. There are two sub-cases.
- case 1.1: $C$ lies on the line segment $A B$ (Figure 2). Clearly, $d(A, C)<$ $d(A, B)$. This contradicts $d(A, B)=\min _{M \in C_{1}, N \in C_{2}} d(M, N)$.


Figure 2: case 1.1

- case 1.2: $A$ lies on the line segment $B C$ (Figure 3). By the definition of convex hull, $A \in C_{2}$. This contradicts the fact that $C_{1}$ and $C_{2}$ don't intersect.
- case 2: $A, B, C$ are not in a line. Consider the triangle $A B C$. We know $d(A, C) \leq d(B, C)$. By Lemma $1, \angle B \leq \angle A$. Also it is always the case that $\angle A+\angle B+\angle C=\pi$ and $\angle C>0$. Therefore, $\angle B<\pi / 2$. Let's consider two sub-cases.
- case 2.1: $\angle C \geq \pi / 2$ (Figure 4). Therefore $\angle C>\angle B$. By Lemma 1, $d(A, B)>d(A, C)$. This contradicts $d(A, B)=\min _{M \in C_{1}, N \in C_{2}} d(M, N)$.


Figure 3: case 1.2


Figure 4: case 2.1

- case 2.2: $\angle C<\pi / 2$ (Figure 5). Since $\angle B<\pi / 2$, there exists a point $D$ on the line segment $B C$ such that $A D \perp B C$. Therefore, $d(A, D)<d(A, B)$. According to the definition of convex hull, $D \in$ $C_{2}$. This contradicts $d(A, B)=\min _{M \in C_{1}, N \in C_{2}} d(M, N)$.


Figure 5: case 2.2

