# Average Case Analysis of Quicksort 

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We assume that all elements are equally likely to be chosen as the pivot element in Partition. When partitioning an array of size $n$ into two subarrays, we have the following possible sizes of the subarrays.

| left array | right array |
| :---: | :---: |
| $n-1$ | 0 |
| $n-2$ | 1 |
| $n-3$ | 2 |
| $\cdot$ | $\cdot$ |
| $\cdot$ | $\cdot$ |
| $\cdot$ | $\cdot$ |
| 1 | $n-2$ |
| 0 | $n-1$ |

So we have

$$
\begin{align*}
T(n) & =\frac{1}{n} \sum_{j=0}^{n-1}(T(j)+T(n-j-1))+n  \tag{1}\\
& =\frac{2}{n} \sum_{k=0}^{n-1} T(k)+n  \tag{2}\\
& =\frac{2}{n} \sum_{k=1}^{n-1} T(k)+n \quad \text { (no recursive call for 0-sized subproblem) } \tag{3}
\end{align*}
$$

Replacing $n$ with $n-1$, we get

$$
\begin{equation*}
T(n-1)=\frac{2}{n-1} \sum_{k=1}^{n-2} T(k)+n-1 \tag{4}
\end{equation*}
$$

Multiplying (3) by $n$ and (4) by $n-1$ :

$$
\begin{align*}
n T(n) & =2 \sum_{k=1}^{n-1} T(k)+n^{2}  \tag{5}\\
(n-1) T(n-1) & =2 \sum_{k=1}^{n-2} T(k)+(n-1)^{2} \tag{6}
\end{align*}
$$

Subtracting (6) from (5), we get

$$
\begin{align*}
& n T(n)-(n-1) T(n-1)=2 T(n-1)+2 n-1  \tag{7}\\
\Longrightarrow & n T(n)=(n+1) T(n-1)+2 n-1  \tag{8}\\
\Longrightarrow & T(n)=\frac{n+1}{n} T(n-1)+2-\frac{1}{n} . \tag{9}
\end{align*}
$$

Ignoring the $1 / n$ term in (9), we have

$$
\begin{equation*}
T(n) \leq \frac{n+1}{n} T(n-1)+2 \tag{10}
\end{equation*}
$$

Method 1: We solve (10) by iteration:

$$
\begin{aligned}
T(n) & \leq 2+\frac{n+1}{n} T(n-1) \\
& \leq 2+\frac{n+1}{n}\left(2+\frac{n}{n-1} T(n-2)\right) \\
& =2+2 \frac{n+1}{n}+\frac{n+1}{n-1}+T(n-2) \\
& \leq 2+2 \frac{n+1}{n}+\frac{n+1}{n-1}\left(2+\frac{n-1}{n-2} T(n-3)\right) \\
& =2+2 \frac{n+1}{n}+2 \frac{n+1}{n-1}+\frac{n+1}{n-2} T(n-3) \\
& =2 \frac{n+1}{n+1}+2 \frac{n+1}{n}+2 \frac{n+1}{n-1}+\frac{n+1}{n-2} T(n-3) \\
. & . \\
& =2(n+1) \sum_{i=0}^{k-1} \frac{1}{(n+1)-i}+\frac{n+1}{n-(k-1)} T(n-k) \\
& =2(n+1) \sum_{i=0}^{n-2} \frac{1}{(n+1)-i}+\frac{n+1}{2} T(1) \\
& =2(n+1) \sum_{j=3}^{n+1} \frac{1}{j}+\frac{n+1}{2} \Theta(1) \\
& =2(n+1) \Theta(\ln n)+\Theta(n) \\
& =\Theta(n \log n)
\end{aligned}
$$

Method 2: Dividing (10) by $n+1$ :

$$
\frac{T(n)}{n+1} \leq \frac{T(n-1)}{n}+\frac{2}{n+1}
$$

Let $R(n)=\frac{T(n)}{n+1}$ (and thus, $R(n-1)=\frac{T(n-1)}{n}$ ). We have

$$
R(n) \leq R(n-1)+\frac{2}{n+1}
$$

Note that $R(1)=\frac{T(1)}{1+1}=\Theta(1)$. We get

$$
\begin{aligned}
R(n) & \leq \frac{2}{n+1}+\frac{2}{n}+\frac{2}{n-1}+\cdots+\frac{2}{3}+R(1) \\
& =2 \sum_{k=3}^{n+1} \frac{1}{k}+\Theta(1) \\
& =\Theta(\ln n)
\end{aligned}
$$

Therefore,

$$
\begin{aligned}
T(n) & =(n+1) R(n) \\
& =(n+1) \Theta(\ln n) \\
& =\Theta(n \ln n)
\end{aligned}
$$

