

byte: 8-bits

nibble: 4-bits

128 64 32 16 8 4 2 1
1 1 1 1 1 1 1 1 = $2^8 - 1 = 255$
87 | 65₁₀

Other Representations: Octal & Hexadecimal

53 - 1 **LSB**
26 - 0
13 - 1
6 - 0
3 - 1
1 - 1 **MSB**

↑

base 8

binary 3 bits
at a time

↑

base 16

4 bits at a time

Octal:

- Base 8 representation
- Binary 3-bits at a time

32 16 8 4 2 1

53₁₀

... 64 8 1
6 5₈

53₁₀ = 110 } 101
6 5

e.s.

$$\begin{array}{r} 53 \\ - 48 = 6 \times 8 \\ \hline 5 \end{array}$$

Hexadecimal :

- Base 16 rep.
- Binary 4-bits at a time

53₁₀

256	16	1
	3	5 ₁₆

$$\begin{array}{r} 53 \\ - 48 = 16 \times 3 \\ \hline 5 \end{array}$$

--- 32 16 8 4 2 1
0011 0101
3 5₁₆

47₁₀

256	16	1
	2	F ₁₆

$$\begin{array}{r} 47 \\ 32 \\ \hline 15 \end{array}$$

0, 1, 2, ..., 9, A, B, C, D, E, F

10 11 - - - 15

Binary mathematical operations

Addition

$$\begin{array}{r} 10110 \rightarrow 22_{10} \\ + 10101 \rightarrow 21_{10} \\ \hline 101011 \rightarrow 43_{10} \quad \checkmark \end{array}$$

32 16 8 4 2 1

Subtraction

	16	8	4	2	1	
	1	0	1	0	0	26 ₁₀
	1	1	0	1	0	
-	0	0	1	1	1	7 ₁₀
	1	0	0	1	1	19 ₁₀ ✓

Representing negative #'s : signed magnitude, one's complement, two's complement

Signed magnitude :

+ 5
- 5

↑ ↑
sign magnitude

4-bit signed mag. rep.

sign 4 2 1

+ 5 : 0 1 0 1

- 5 : 1 1 0 1

~~~~~

sign      mag.

[-7, +7]

+ 0 : 0 0 0 0

- 0 : 1 0 0 0

# One's complement

4-bit 1's comp.

① Positive #s  $\rightarrow$  represented in straight binary

e.g.  $5_{10} \rightarrow \boxed{0101}$

② Negative #s  $\rightarrow$  write the magnitude in straight binary  
and flip all the bits

e.g.  $-5_{10} \rightarrow$  mag: 0101  
flip:  $\boxed{1010}$

largest pos #  
0111  $\rightarrow +7$   
most neg #  
1000  $\rightarrow -7$   
[-7, +7]

+0 0000  
-0 1111

leading bit is sign

0  $\rightarrow$  +

1  $\rightarrow$  -

0110  $\rightarrow +6$   
↑  
+ 6

1011  $\rightarrow -4$   
- 4  
100 = 4

Observation:

what is "flip all the bits"

0101  $\xrightarrow{\text{flip}}$  1010

subtract from 1111

$$\begin{array}{r} 1 \ 1 \ 1 \ 1 \\ - \ 0 \ 1 \ 0 \ 1 \\ \hline 1 \ 0 \ 1 \ 0 \end{array} \quad \begin{array}{l} = 2^4 - 1 = 15 \\ = -5 \\ \hline 10 \end{array}$$

neg #'s in 1's comp  $-x : (2^4 - 1) - x$  in straight bin

2's complement: pos #'s  $\rightarrow$  rep. in straight bin  $+5 \rightarrow 0101$

neg #'s  $\rightarrow$  rep. mag. in bin  $-5 \rightarrow$  <sup>mag.</sup> 0101

• flip all bits  $\rightarrow 1010$

• add 1  $\rightarrow$  1011

$+x \rightarrow x$  in straight binary

$-x \rightarrow 2^4 - x$  in straight binary

# Examples

2's complement: subtraction is addition

①  $4_{10} - 3_{10} = 1_{10}$

$$\begin{array}{r} 4 \quad 0100 \\ + -3 \quad +1101 \\ \hline +1 \quad 10001 \\ \hline \end{array}$$

✓

|       |      |                          |
|-------|------|--------------------------|
| mag 3 | 0011 |                          |
| flip  | 1100 |                          |
| +1    | 1101 | → 2's comp<br>rep. of -3 |

$3_{10} - 4_{10} = -1_{10}$

$$\begin{array}{r} 3 \quad 0011 \\ + -4 \quad +1100 \\ \hline 1111 \end{array}$$

✓

|       |      |                      |
|-------|------|----------------------|
| mag 4 | 0100 |                      |
| flip  | 1011 |                      |
| +1    | 1100 | → 2's comp<br>for -4 |

-1:

|      |      |   |
|------|------|---|
| mag  | 0001 |   |
| flip | 1110 |   |
| +1   | 1111 | ✓ |

Range of 4-bit 2's comp.

biggest pos #: 0111  $\rightarrow$  +7

Smallest neg #: -8  
(most) 1000

|      |      |
|------|------|
| neg: | 1000 |
| flg: | 0111 |
| +1   | 1000 |

$\{-8, +7\}$

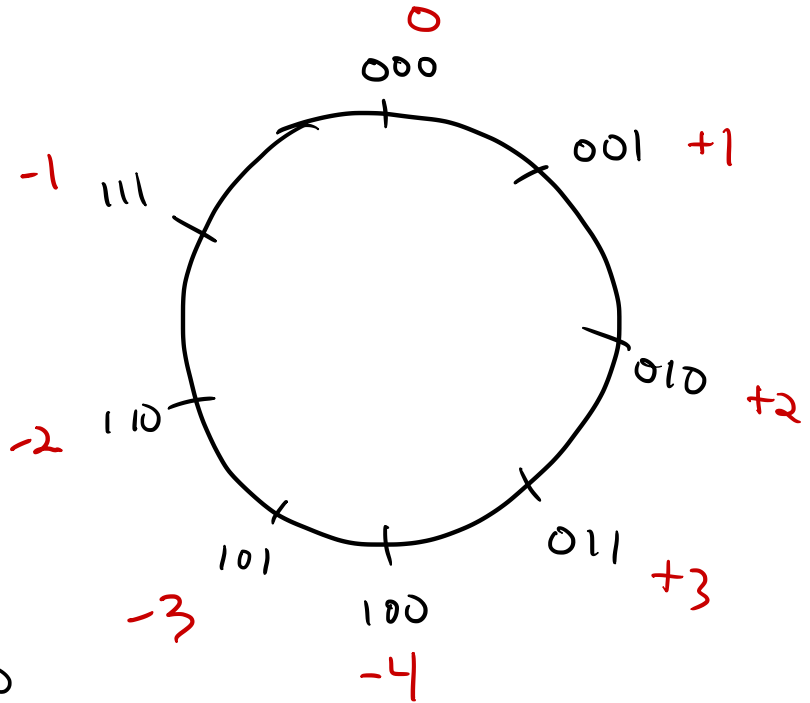


Virtual 2's comp.

3-bit 2's comp

| #  |     |
|----|-----|
| +0 | 000 |
| +1 | 001 |
| +2 | 010 |
| +3 | 011 |
| -4 | 100 |
| -3 | 101 |
| -2 | 110 |
| -1 | 111 |

|    |    |     |
|----|----|-----|
| -2 | 2  | 010 |
|    | 4p | 101 |
|    | +1 | 110 |



Digression:

## Logic Gates

motivating  
example

addition

$$\begin{array}{r}
 a: \quad 1 \ 0 \ 1 \\
 b: + 1 \ 1 \ 0 \\
 \hline
 \end{array}$$

carry bit (0)  
 $a_0$  (1)  
 $b_0$  (0)  
 Sum bit  $S_0$  (1)

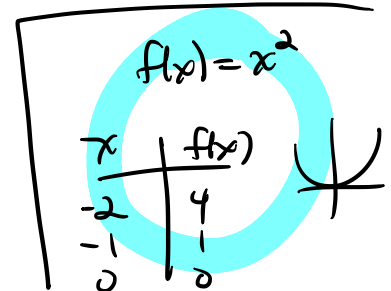
$$\begin{array}{r}
 1 \ 1 \ 1 \\
 + 1 \ 0 \ 1 \\
 \hline
 \end{array}$$

carry bit (1)  
 $a_0$  (1)  
 $b_0$  (1)  
 Sum bit  $S_0$  (0)

Sum bit

| $a_0$ | $b_0$ | $S_0$ |
|-------|-------|-------|
| 0     | 0     | 0     |
| 0     | 1     | 1     |
| 1     | 0     | 1     |
| 1     | 1     | 0     |

| $a_0$ | $b_0$ | $c$ |
|-------|-------|-----|
| 0     | 0     | 0   |
| 0     | 1     | 0   |
| 1     | 0     | 0   |
| 1     | 1     | 1   |



# Recap 2's comp.

4-bit 2's comp. <sup>sign</sup>  
↓ 4 2 1

math

① pos #'s → rep. in straight binary es. +6 0 1 1 0 x

② neg #'s → (a) rep mag. in binary es -6 0 1 1 0 x

(b) flip all bits 1 0 0 1  $(2^4 - 1) - x$

(c) add 1 1 0 1 0  $2^4 - x$

16 - 6 = 10

$$\begin{array}{r} 1111 \leftarrow 2^4 - 1 \\ - 0110 \leftarrow 6 \\ \hline 1001 \quad (2^4 - 1) - 6 \end{array}$$

E.g.

3: 0011

4: 0100

-3:

0011  
1100

1101

-4:

0100  
1011

1100

①  $4 - 3 \Rightarrow 1$

$$\begin{array}{rcl}
 4: & 0100 & \rightarrow 4 \\
 + -3: & 1101 & \rightarrow 2^4 - 3 \\
 \hline
 & 10001 & 2^4 + 1
 \end{array}$$

②  $3 - 4 \Rightarrow -1$

$$\begin{array}{rcl}
 3: & 0011 & \rightarrow 3 \\
 + -4: & 1100 & \rightarrow 2^4 - 4 \\
 \hline
 & 1111 & 2^4 - 1
 \end{array}$$

# Converting back from 2's comp:

one way

1 0 1 0

- 0 0 0 1

-----

1 0 1 1

0 1 0 0

→ 4

⇒ -4

$2^4 - x$

-1

$(2^4 - x) - 1$

flip

$(2^4 - 1) - [2^4 - x - 1]$

$x$

1 1 0 0

$2^4 - x$

flip

0 0 1 1

$(2^4 - 1) - [2^4 - x]$

add 1

0 1 0 0 ✓

$(2^4 - 1) - [2^4 - x] + 1$

$= x$

to convert a neg  
2's comp. # to its  
magnitude:

- ① flip all bits
- ② add 1
- ③ result is neg. mag.

## Overflow Rules

$$\begin{array}{r} 4: \quad 0100 \\ + 4: \quad 0100 \\ \hline 1000 \end{array}$$

$$\begin{array}{r} -6 \quad 0110 \\ \quad 1001 \\ \hline \boxed{1010} \\ -6 \quad 11010 \\ + -6 \quad 11010 \\ \hline 1 \boxed{0100} \end{array}$$

① pos + pos must be pos

$\Rightarrow$  if you have leading zeroes for both #s, but answer has a leading one  
 $\Rightarrow$  overflow

② neg + neg must be neg

$\Rightarrow$  if you have leading ones for both #s, but answer has leading zero  
 $\Rightarrow$  underflow