

# Supplementary Materials

## Enumerating consistent subgraphs of directed acyclic graphs: an insight into biomedical ontologies

Yisu Peng\*, Yuxiang Jiang\*, and Predrag Radivojac

Department of Computer Science, Indiana University, Bloomington, Indiana, U.S.A.

This document contains supplementary text, algorithms, and detailed results related to the main manuscript.

## S1 Implementation Details

### S1.1 Handling Large Integers

The GNU Multiple Precision Arithmetic Library, available at <https://gmplib.org/>, was used to operate with large integers.

### S1.2 Source Code

The `cdag` algorithm is implemented in C++ and is freely available from GitHub, at the following location: <https://github.com/shawn-peng/counting-consistent-sub-DAG>.

## S2 Simulated directed acyclic graphs

We used Algorithm S1 to generate random directed acyclic graphs in Section 5.1.

## S3 Empirical Verification of Correctness

We generated over 100,000 random graphs with a single root and the number of vertices less than or equal to 25 (Algorithm S1). The in-degrees of the vertices were generated according to a Poisson distribution whose parameter  $\lambda$  was in turn sampled from a Gamma prior  $\Gamma(2.0, 1.0)$ . The parameter  $\lambda$  was kept constant in each individual graph and then generated anew for the next graph.

A brute-force program was used to verify the correctness of our algorithm. This program simply generates all vertex-induced subgraphs of the input graph and then checks for their consistency.

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\*Contributed equally to this work.

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**Algorithm S1:** Generating a random directed acyclic graph.

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**Input :** Desired number of vertices  $N_v$  of the DAG, Gamma shape parameters  $(\alpha, \beta)$

**Output:** A directed acyclic graph  $\mathcal{O}$

```

1 Function randdag( $\mathcal{O}$ )
2    $\lambda \sim \Gamma(\alpha, \beta)$ 
3   add vertex(1) to  $\mathcal{O}$  as root
4   for  $i = 2$  to  $N_v$  do
5      $X \sim \text{Pois}(\lambda)$ 
6      $n_p \leftarrow \min(X + 1, |\mathcal{O}.V|)$ 
7     add vertex( $i$ ) to  $\mathcal{O}$ 
8     select random  $n_p$  vertices from  $\mathcal{O}.V$ 
9     set these  $n_p$  vertices to be parents of vertex( $i$ )
10    end
11 end

```

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## S4 Proof for the Minimum Bound Pivot Selection

The following proposition is used for the “minimum bound” pivot selection strategy in the acceleration of the algorithm.

**Proposition S4.1.** *Given an ontology  $\mathcal{O}$ , let  $n$ ,  $e$ ,  $r$  be the number of its vertices, edges and roots, respectively. The quantity  $U = e - n + r$  is an upper-bound of the number of multi-parent vertices ( $m$ ) in  $\mathcal{O}$ .*

*Proof.* We prove this proposition by induction on the number of vertices. Since an ontology is a DAG, there exists a topological order of all vertices, say  $v_1, v_2, \dots, v_n$ . We consider a graph reconstruction by adding back vertices along with its incoming edges in this order.

Starting from an empty graph, we have both  $m$  and  $U$  being equal to 0, so the upper-bound holds. Suppose now that this bound holds for all graphs of size  $n = k$ ; i.e.,  $U^{(k)} \geq m^{(k)}$ . Add the next vertex  $v_{k+1}$  whose incoming degree is  $d$ .

If  $d \geq 1$ , we have  $e^{(k+1)} = e^{(k)} + d$ ,  $n^{(k+1)} = n^{(k)} + 1$  and  $r^{(k+1)} = r^{(k)}$ , hence

$$\begin{aligned}
U^{(k+1)} &= U^{(k)} + (d - 1) \\
&\geq m^{(k)} + (d - 1) \\
&\geq m^{(k)} + \min(d - 1, 1) \\
&= m^{(k+1)}.
\end{aligned}$$

Otherwise,  $d = 0$ ; i.e., the next vertex being added is a root, in which case  $e^{(k+1)} = e^{(k)}$ ,  $n^{(k+1)} = n^{(k)} + 1$  and  $r^{(k+1)} = r^{(k+1)} + 1$ ; hence,

$$U^{(k+1)} = U^{(k)} \geq m^{(k)} = m^{(k+1)}.$$

Therefore, this upper-bound holds for both cases, and by induction, the proposition holds.  $\square$

## S5 Additional Algorithms

Here we show the additional algorithms that are called in the advanced algorithm in the main paper: (i) the extended algorithm on trees (ii) the algorithm that identifies all branching vertices

and (iii) the flow algorithm which is used in both branching vertex detection and one of the pivot strategies.

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**Algorithm S2:** The generalized version of `ctree` for trees.

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**Input :** A tree  $\mathcal{T}_r$ , rooted at  $r$ .  
**Output:** The number of consistent subgraphs in  $\mathcal{T}_r$ .

```

1 Function ctree*( $\mathcal{T}_r$ )
2   if  $\mathcal{T}_r$  is empty then
3     return 1;
4   else
5     return  $\varphi(r) + \prod_{u \in \mathcal{C}(r)} \text{ctree}^*(\mathcal{T}_u)$ ;    // compatible with pruning, reversing
6     ;
7   end
8 end
```

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**Algorithm S3:** Branching vertex detection.

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**Input :** A directed acyclic graph  $\mathcal{O}$   
**Output:** The branching vertices in  $\mathcal{O}$  in a reversed topological order.

```

1 Function branching_vertices( $\mathcal{O}$ )
2   flow( $\mathcal{O}$ );
3   foreach  $v$  in  $\mathcal{O}$  in a reversed topological order do
4     if  $v.\text{flow} = |\mathcal{D}^+(v)|$  then
5       append( $list, v$ );
6     end
7   end
8   return  $list$ ;
9 end
```

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## S6 Results on Biomedical Ontologies

This section provides more precise counts of consistent subgraphs in four biomedical ontologies, as discussed in Section 5.2 of the main paper and visualized in Figure 4. Table S1 provides counts that were completed before submitting this manuscript. The exact counts are available upon request.

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**Algorithm S4:** Propagate flows in a DAG.

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**Input :** A directed acyclic graph  $\mathcal{O}$

**Output:**  $\mathcal{O}$  with flow value saved to each vertex.

```

1 Function flow( $\mathcal{O}$ )
2   foreach  $v$  in  $\mathcal{O}$  in a reversed topological order do
3      $v.\text{flow} \leftarrow 1 + \sum_{u \in \mathcal{C}(v)} (u.\text{flow} / |\mathcal{P}(u)|)$ 
4   end
5   return  $\mathcal{O}$ ;
6 end

```

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Table S1: Number of consistent subgraphs in biomedical ontologies including vertices with up to a particular level.  $n$ : number of vertices,  $e$ : number of edges,  $m$ : number of multi-parent vertices,  $\ell$ : number of leaves, nrc: number of recursive calls.

Ontology	level	$n$	$e$	$m$	$\ell$	count	lower bound	upper bound	nrc
MFO	2	21	20	0	20	1.048e6	1.048e6	1.048e6	1
	3	156	155	0	143	5.465e43	1.115e43	5.465e43	1
	4	662	728	54	578	3.202e179	9.893e173	1.244e181	384
	5	2169	2460	247	1870	1.497e583	8.435e562	7.451e587	1.067e4
	6	6587	7403	722	5860	2.504e1812	1.086e1764	6.573e1827	5.857e4
	7	8394	9532	1014	7237	1.505e2256	3.582e2178	1.381e2279	1.521e5
	8	9415	10816	1226	7935	2.008e2491	4.710e2388	1.077e2524	3.686e6
	9	9996	11621	1407	8306	2.036e2616	2.265e2500	3.045e2659	1.075e7
	10	10300	12094	1536	8481	-	1.085e2553	9.609e2725	-
	11	10470	12430	1648	8573	-	5.372e2580	5.868e2763	-
	12	10641	12755	1755	8639	-	3.964e2600	2.209e2802	-
	13	10709	12861	1780	8672	-	3.405e2610	5.656e2816	-
	14	10762	12997	1828	8697	-	1.143e2618	2.554e2830	-
	15	10779	13025	1835	8703	-	7.312e2619	8.018e2833	-
	16	10788	13044	1843	8703	-	7.312e2619	4.021e2835	-
	17	10789	13046	1844	8703	-	7.312e2619	5.577e2835	-
MFO <sup>u</sup>	2	20	19	0	19	5.242e5	5.243e5	5.242e5	1
	3	143	142	0	130	6.724e39	1.361e39	6.724e39	1
	4	541	595	43	463	8.386e144	2.382e139	1.970e146	330
	5	1561	1789	188	1282	6.191e405	8.326e385	2.966e410	8449
	6	3688	4299	524	3033	5.892e959	1.057e913	8.019e973	5.136e4
	7	4952	5811	751	3924	6.943e1257	1.745e1181	4.438e1278	8.239e4
	8	5704	6780	926	4403	2.387e1424	2.723e1325	1.672e1454	1.348e6
	9	6108	7360	1069	4619	1.239e1503	2.868e1390	6.152e1541	5.957e6
	10	6342	7709	1161	4740	-	7.624e1426	4.269e1591	-
	11	6474	7977	1253	4804	-	1.406e1446	1.493e1620	-
	12	6591	8200	1327	4842	-	3.866e1457	5.313e1645	-
	13	6634	8268	1344	4860	-	1.013e1463	4.432e1654	-
	14	6654	8320	1363	4870	-	1.038e1466	3.623e1659	-
	15	6661	8333	1367	4872	-	4.151e1466	7.937e1660	-
	16	6662	8334	1367	4872	-	4.151e1466	1.160e1661	-

Ontology	level	$n$	$e$	$m$	$\ell$	count	lower bound	upper bound	nrc
BPO	2	22	21	0	21	2.097e6	2.097e6	2.097e6	1
	3	123	147	25	101	2.376e31	2.535e30	1.311e32	238
	4	573	734	149	471	1.430e147	6.097e141	4.118e150	2.183e4
	5	1777	2631	672	1365	7.835e436	8.053e410	1.177e453	5.308e6
	6	4288	7182	2049	3057	-	1.773e920	1.616e1067	-
	7	8109	14783	4593	5198	-	5.674e1564	8.161e1961	-
	8	12719	24596	7966	7449	-	2.357e2242	4.714e3004	-
	9	17243	34704	11414	9397	-	6.010e2828	1.309e4015	-
	10	21445	44378	14816	11036	-	1.469e3322	1.247e4938	-
	11	24683	51711	17443	12102	-	1.161e3643	3.029e5637	-
	12	26944	56834	19280	12780	-	1.457e3847	3.641e6118	-
	13	28424	60105	20489	13231	-	8.470e3982	1.518e6433	-
	14	29134	61702	21071	13401	-	1.268e4034	2.585e6577	-
	15	29402	62314	21294	13439	-	3.484e4045	2.794e6630	-
	16	29509	62551	21377	13465	-	2.338e4053	9.061e6652	-
	17	29551	62648	21410	13477	-	9.577e4056	1.750e6661	-
	18	29572	62697	21431	13485	-	2.452e4059	5.407e6665	-
	19	29575	62700	21431	13487	-	9.807e4059	2.433e6666	-
BPO <sup>u</sup>	2	22	21	0	21	2.097e6	2.097e6	2.097e6	1
	3	114	137	24	92	2.714e29	4.952e27	2.714e29	268
	4	489	626	128	398	7.082e127	6.456e119	7.082e127	2.470e4
	5	1471	2219	585	1106	1.458e357	8.693e332	3.142e372	4.579e6
	6	3419	5776	1667	2353	-	2.107e708	6.968e841	-
	7	6267	11369	3539	3851	-	1.847e1159	1.062e1503	-
	8	9649	18501	5980	5359	-	1.659e1613	1.410e2257	-
	9	12935	25726	8461	6622	-	2.634e1993	8.698e2980	-
	10	15775	32228	10762	7518	-	1.392e2263	1.322e3589	-
	11	17932	37166	12525	8039	-	9.553e2419	2.283e4043	-
	12	19481	40657	13745	8405	-	1.436e2530	3.209e4367	-
	13	20472	42852	14565	8588	-	1.760e2585	1.618e4571	-
	14	20918	43853	14923	8639	-	3.964e2600	2.116e4658	-
	15	21077	44218	15053	8641	-	1.586e2601	4.939e4688	-
	16	21128	44336	15094	8640	-	7.928e2600	1.059e4698	-
	17	21148	44385	15111	8639	-	3.964e2600	3.202e4701	-
	18	21154	44402	15117	8641	-	1.586e2601	6.313e4702	-

Ontology	level	$n$	$e$	$m$	$\ell$	count	lower bound	upper bound	nrc
CCO	2	16	15	0	15	32769	32768	32769	1
	3	78	84	7	67	8.155e20	1.476e20	8.155e20	26
	4	351	413	52	301	4.895e94	4.074e90	3.441e95	328
	5	623	840	166	476	2.249e156	1.951e143	3.113e162	2441
	6	953	1416	346	713	2.127e233	4.309e214	1.387e245	2.192e4
	7	1225	1881	470	880	3.543e291	8.061e264	9.098e310	1.669e5
	8	1737	2765	742	1301	1.492e423	4.365e391	1.110e451	9.573e6
	9	2185	3584	969	1600	-	4.446e481	4.844e565	
	10	2619	4534	1292	1899	-	4.529e571	2.657e680	
	11	3051	5482	1578	2164	-	2.685e651	5.198e786	
	12	3491	6425	1887	2400	-	2.965e722	2.320e894	
	13	3780	7105	2119	2540	-	4.132e764	3.286e962	
	14	3964	7533	2255	2630	-	5.116e791	7.860e1006	
	15	4061	7751	2321	2668	-	1.406e803	9.184e1029	
	16	4082	7798	2336	2672	-	2.250e804	3.675e1034	
	17	4084	7801	2337	2672	-	2.250e804	1.102e1035	
	18	4085	7802	2337	2672	-	2.250e804	1.469e1035	
CCO <sup>u</sup>	2	14	13	0	13	8193	8192	8193	1
	3	66	72	7	55	2.116e17	3.603e16	2.116e17	26
	4	269	316	38	228	5.486e71	4.314e68	2.721e72	246
	5	455	619	124	339	1.287e112	1.120e102	1.639e117	1505
	6	675	1023	255	481	5.816e160	6.243e144	6.137e170	1.058e4
	7	864	1361	352	585	4.194e198	1.266e176	2.439e215	4.806e4
	8	1159	1945	558	794	7.022e266	1.042e239	4.178e291	3.177e6
	9	1441	2479	711	961	-	1.949e289	5.107e360	
	10	1761	3176	954	1181	-	3.284e355	3.304e444	
	11	2080	3886	1162	1357	-	3.146e408	7.276e521	
	12	2376	4521	1372	1504	-	5.612e452	3.908e593	
	13	2565	4966	1519	1591	-	8.684e478	3.450e637	
	14	2695	5270	1613	1656	-	3.204e498	6.089e669	
	15	2765	5432	1662	1682	-	2.150e506	7.189e686	
	16	2775	5455	1668	1682	-	2.150e506	2.361e689	
	17	2776	5456	1668	1682	-	2.150e506	3.541e689	
	18	2777	5457	1668	1682	-	2.150e506	4.722e689	

Ontology	level	$n$	$e$	$m$	$\ell$	count	lower bound	upper bound	nrc
HPO	2	6	5	0	5	33	32	33	1
	3	59	58	0	53	1.374e16	9.007e15	1.374e16	1
	4	285	288	4	238	3.946e73	4.417e71	4.199e73	158
	5	867	925	59	698	5.195e218	1.315e210	3.386e219	1.329e6
	6	2135	2340	199	1635	2.088e523	1.528e492	7.485e526	2.523e6
	7	4162	4695	508	3057	-	1.773e920	1.126e1003	
	8	6375	7330	891	4514	-	7.079e1358	9.364e1506	
	9	8307	9679	1270	5777	-	1.123e1739	1.740e1946	
	10	9733	11489	1616	6653	-	5.657e2002	2.414e2264	
	11	10700	12841	1958	7244	-	4.584e2180	2.020e2481	
	12	11271	13892	2311	7508	-	1.359e2260	2.226e2605	
	13	11709	14841	2641	7712	-	3.494e2321	1.920e2705	
	14	12121	15862	3035	7872	-	5.107e2369	4.017e2806	
	15	12149	15920	3054	7888	-	3.347e2374	1.085e2813	
	16	12167	15974	3072	7891	-	2.677e2375	8.845e2817	
HPO <sup>u</sup>	2	5	4	0	4	17	16	17	1
	3	45	44	0	40	1.753e12	1.100e12	1.753e12	1
	4	206	209	4	166	1.005e52	9.354e49	1.134e52	134
	5	606	656	51	463	2.224e147	2.382e139	1.227e148	7.231e5
	6	1481	1649	164	1062	4.622e350	4.941e319	5.988e353	5.993e5
	7	2785	3203	402	1871	-	1.687e563	1.123e647	
	8	4097	4824	680	2625	-	1.599e790	3.328e930	
	9	5213	6239	953	3241	-	4.347e975	1.581e1169	
	10	6026	7328	1200	3653	-	4.598e1099	1.882e1338	
	11	6470	7956	1357	3839	-	4.510e1155	3.712e1427	
	12	6667	8289	1470	3884	-	1.587e1169	2.640e1466	
	13	6745	8441	1518	3892	-	4.062e1171	8.671e1480	
	14	6777	8524	1546	3874	-	1.550e1166	5.623e1486	
	15	6783	8538	1552	3875	-	3.099e1166	7.346e1487	
	16	6785	8544	1554	3873	-	7.748e1165	2.145e1488	