CS4800: Algorithms — S'18 — Jonathan Ullman

Homework 7
Due Friday Mar 16 at 11:59pm via Gradescope

Name:
Collaborators:

- Make sure to put your name on the first page. If you are using the LaTeX template we provided, then you can make sure it appears by filling in the `yourname` command.
- This assignment is due Friday Mar 16 at 11:59pm via Gradescope. No late assignments will be accepted. Make sure to submit something before the deadline.
- Solutions must be typeset in LaTeX. If you need to draw any diagrams, you may draw them by hand as long as they are embedded in the PDF. I recommend using the source file for this assignment to get started.
- I encourage you to work with your classmates on the homework problems. *If you do collaborate, you must write all solutions by yourself, in your own words.* Do not submit anything you cannot explain. Please list all your collaborators in your solution for each problem by filling in the `yourcollaborators` command.
- Finding solutions to homework problems on the web, or by asking students not enrolled in the class is strictly prohibited.

Problem 1. *Minimum Spanning Trees*

In this problem we will see a new algorithm for finding an MST—the*anti-Kruskal* algorithm. Recall that *Kruskal's algorithm* starts wth $T = \emptyset$, considers edges e in *ascending* order of weight, and adds e to T as long as doing so would not create a cycle. The *anti-Kruskal algorithm* starts with $T = E$, considers edges e in *descending* order of weight, and removes e from T as long as doing so would not make T disconnected.

Explain why the anti-Kruskal algorithm outputs an MST T. You may assume that all edge-weights are distinct and you may use the cut and cycle properties of MSTs without proof.

Solution:

Problem 2. *All-Pairs Shortest Paths*

In the *all-pairs shortest paths* problem, you are given a directed, weighted graph with edge lengths $G = (V, E, \{\ell_e\})$, and have to find the length of the shortest path from s to t for *every pair* $s, t \in V$. For this HW we only want the *length* of the shortest path and not the path itself.

If all edge lengths are non-negative ($\ell_e \geq 0$), then we can solve this problem by running Dijkstra's algorithm from every source node $s \in V$, incurring running time $O(nm \log n)$. However, if lengths can be negative, then running Bellman-Ford from each source node $s \in V$ incurs running time $O(n^2 m)$. In this question we will study the following algorithm for solving all-pairs shortest paths in graphs with negative-length edges, but no negative-length cycles.

- Modify the input graph by adding an additional node z connected to every other node v by a zero-length edge (z, v).
- Run the Bellman-Ford algorithm on the modified graph with source z to find the length $f(v)$ of the shortest $z \to v$ path in the modified graph.
- Define new edge lengths $\ell'_{u,v} = \ell_{u,v} + f(u) - f(v)$ and let $G' = (V, E, \{\ell'_e\})$ be the input graph with these modified edge weights.
- For each source $s \in V$, run Dijkstra's algorithm on the graph G' with source s to find the length $d'(s, v)$ of the shortest $s \to v$ path in G' for every node v.
- For every $u, v \in V$, let $d(u, v) = d'(u, v) - f(u) + f(v)$. Output the values $\{d(u, v)\}$.

In this problem, we will show correctness and analyze the running time of this algorithm. The final three steps of the problem form the proof of correctness.

(a) What is the running time of this algorithm? Briefly explain your answer.

Solution:

(b) Prove that every edge in G' has non-negative length. That is, $\forall u, v \in V, \ell'_{u,v} \geq 0$. (Thus, Dijkstra's algorithm will correctly find the length $d'(u, v)$ of the shortest $u \to v$ path in G'.)

Solution:

(c) Prove that for every $u \to v$ path $P = u \to w_1 \to \cdots \to w_{k-1} \to v$, we have $\ell'_P = \ell_P + f(u) - f(v)$ where ℓ'_P, ℓ_P are the length of the path in G' and G, respectively.

Solution:

(d) Prove that for every $u, v \in V$, $d'(u, v) - f(u) + f(v)$ is the length of the shortest $u \to v$ path in the original graph G. Thus, the final lengths output by this algorithm are correct.

Solution: