CS3000: Algorithms & Data
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Lecture 20:

- Network Flow: flows, cuts, duality
- Ford-Fulkerson

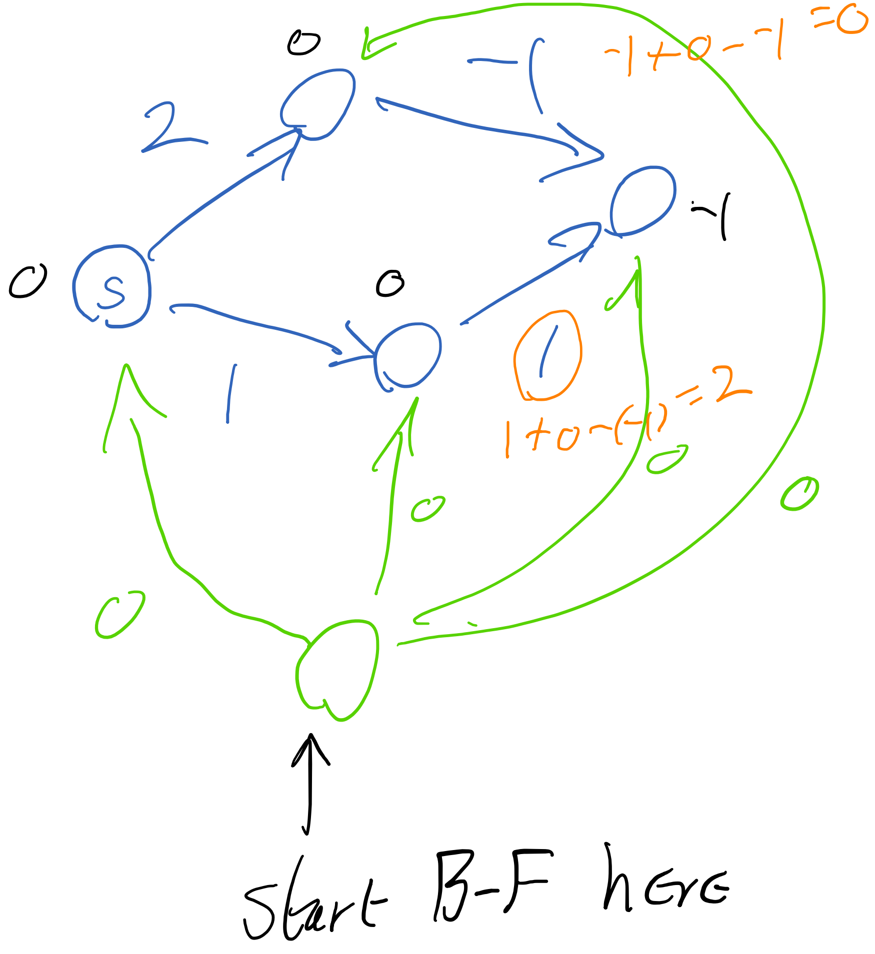
Apr 8, 2019

HW 7 # 1

Idea:

Use Bellman Ford to get a problem w/ nonneg weights.

- Run Dijkstra's Alg Starting from every Node



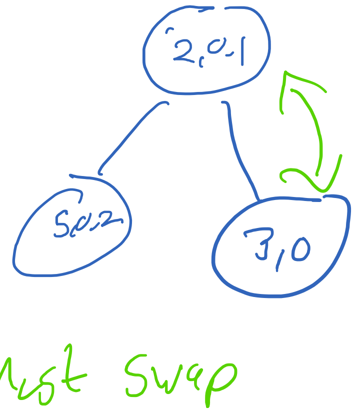
— Values of f

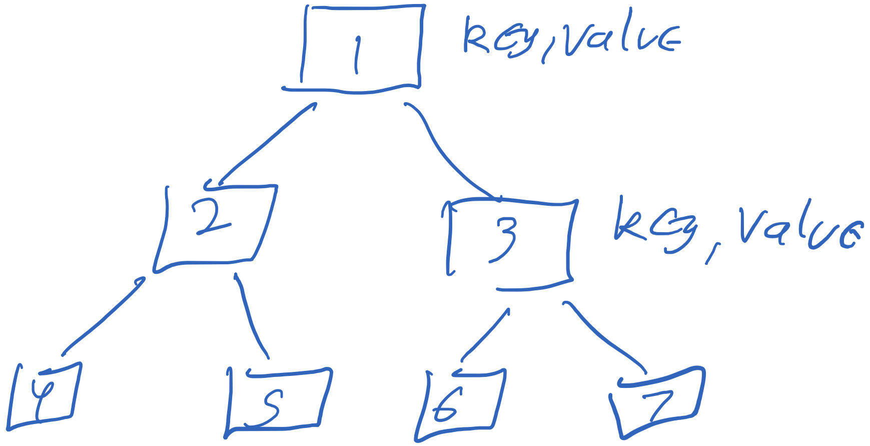
$l'_{u,v} = l_{u,v} + f(u) - f(v)$

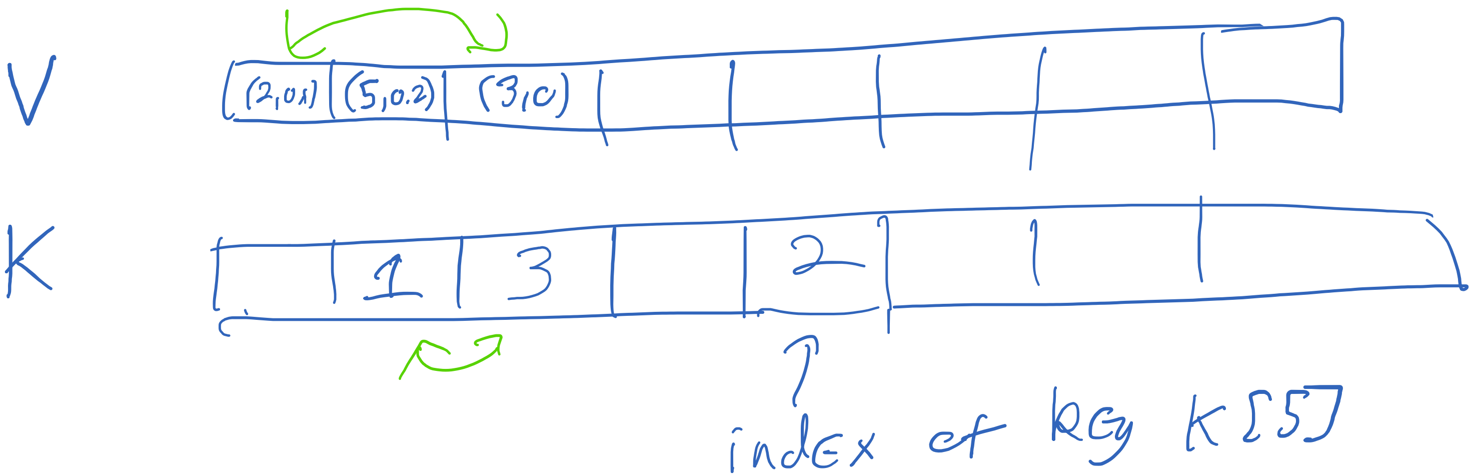
HW 7 # 3

Implement a Priority Queue using Binary Heap

Add (key, value) of (2, 0.1)
Add (5, 0.2)
Add (3, 0)







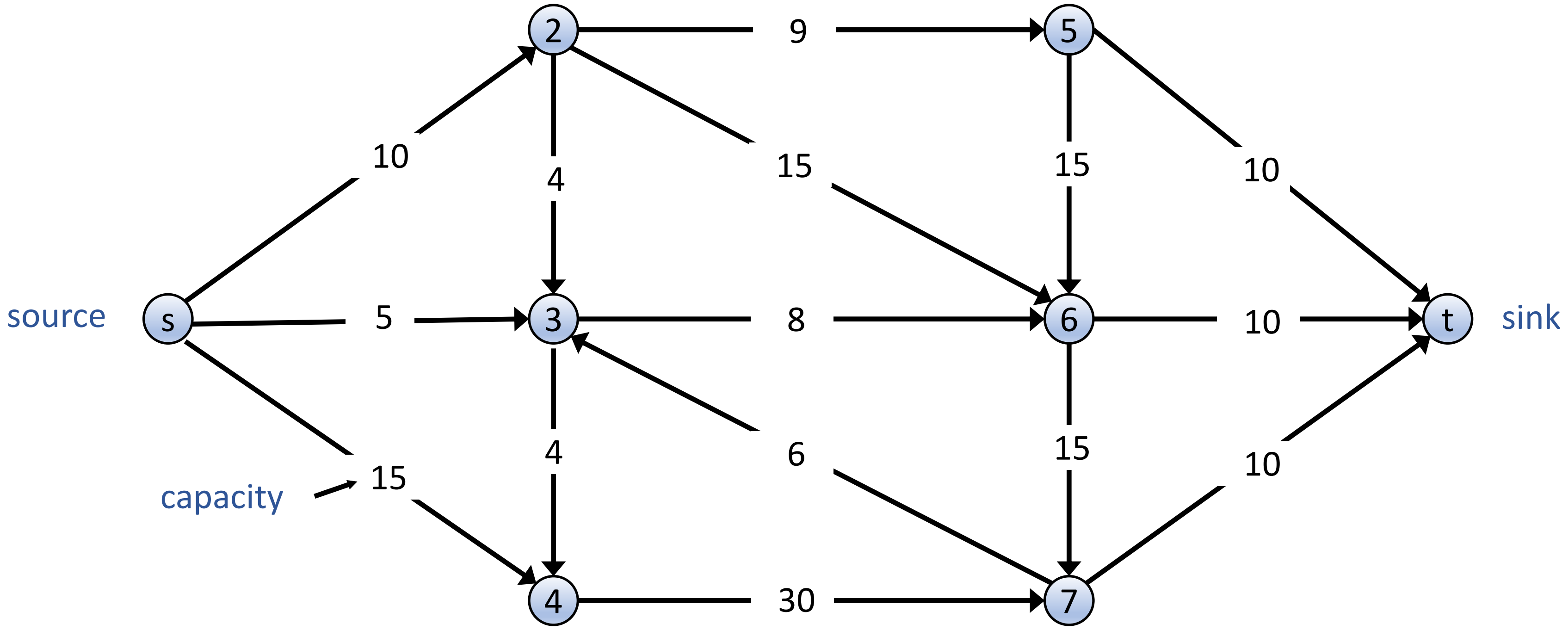
key is #
from 1 ... n

Flow Networks

Flow Networks

- Directed graph $G = (V, E)$
- Two special nodes: source s and sink t
- Edge capacities $c(e)$

Idea: Equilibrium
what flows in
must flow out



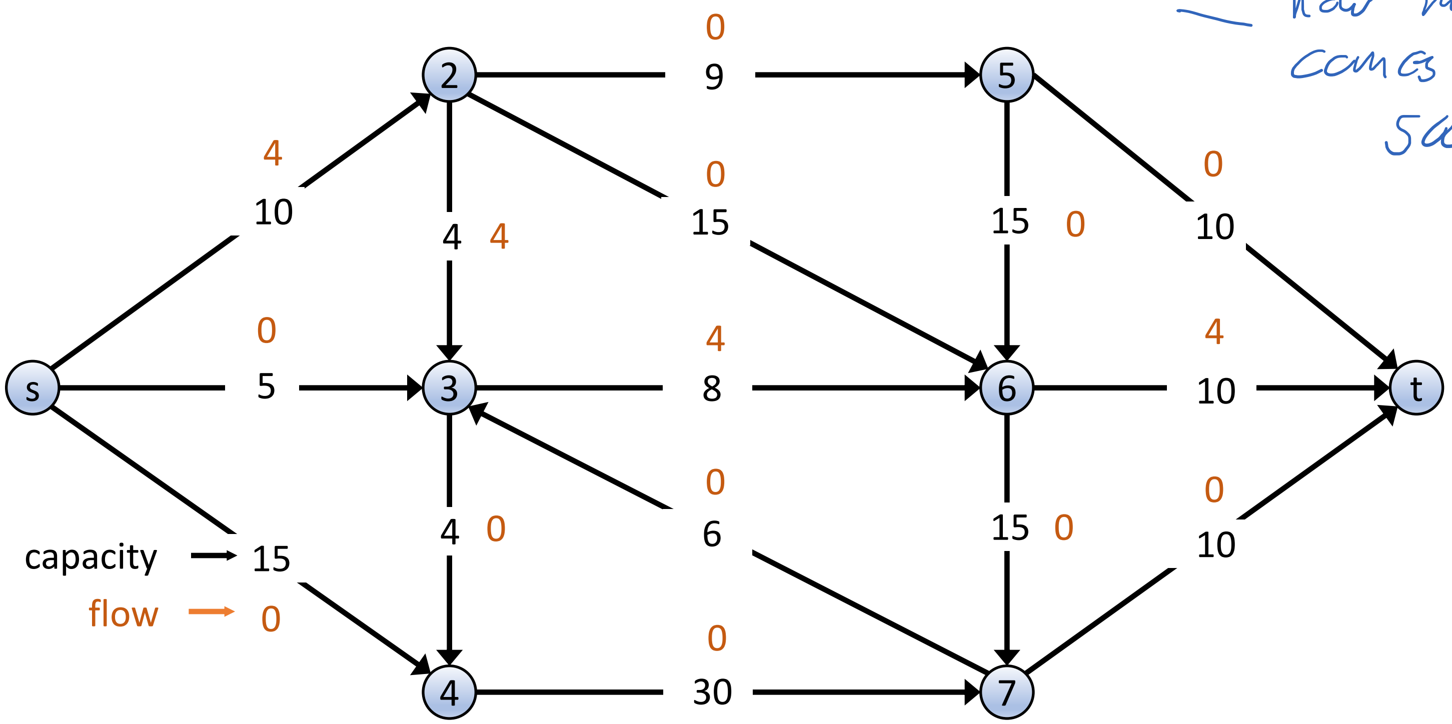
Flows

function defined on edges NOT VERTICES

- An s-t flow is a function $f(e)$ such that
  - For every $e \in E$, $0 \leq f(e) \leq c(e)$ (capacity)
  - For every $v \in V, v \neq s, v \neq t$, $\sum_{e \text{ in to } v} f(e) = \sum_{e \text{ out of } v} f(e)$ (conservation)

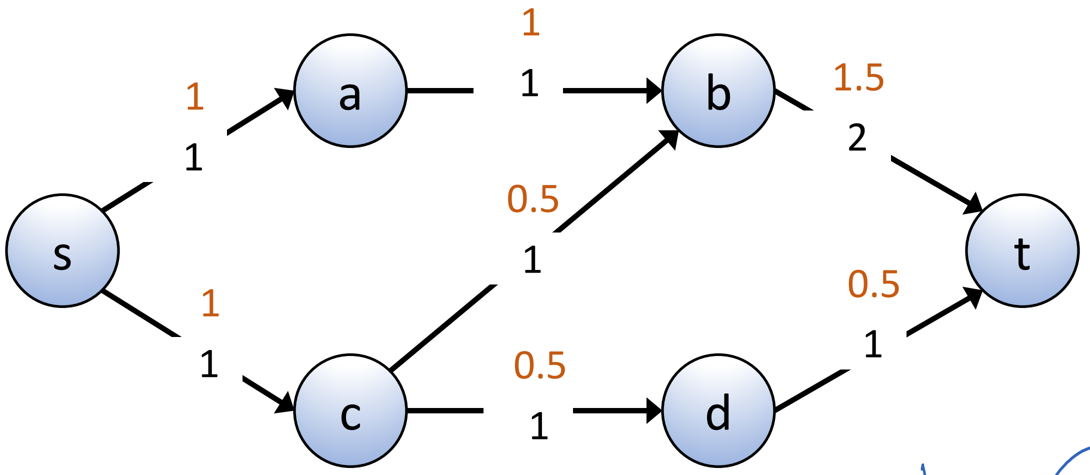
    what goes in what goes out
- The value of a flow is $val(f) = \sum_{e \text{ out of } s} f(e)$

how much comes out of source

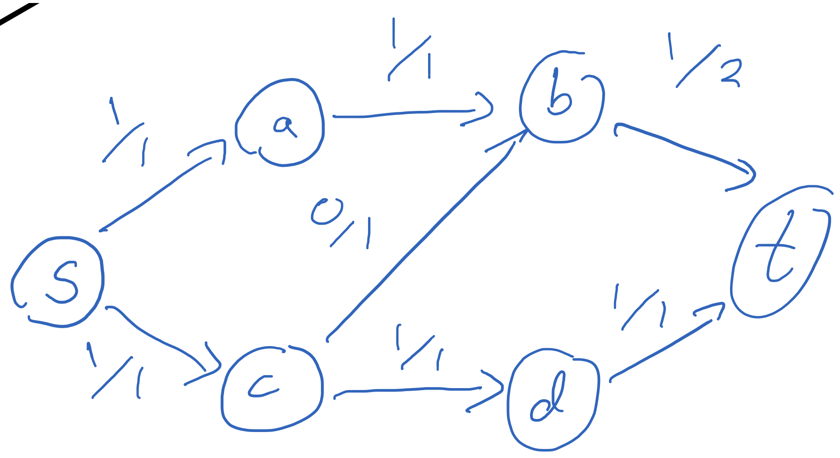


Maximum Flow Problem

- Given G = (V,E,s,t,{c(e)}), find an s-t flow of maximum value
- Is this a maximum flow?



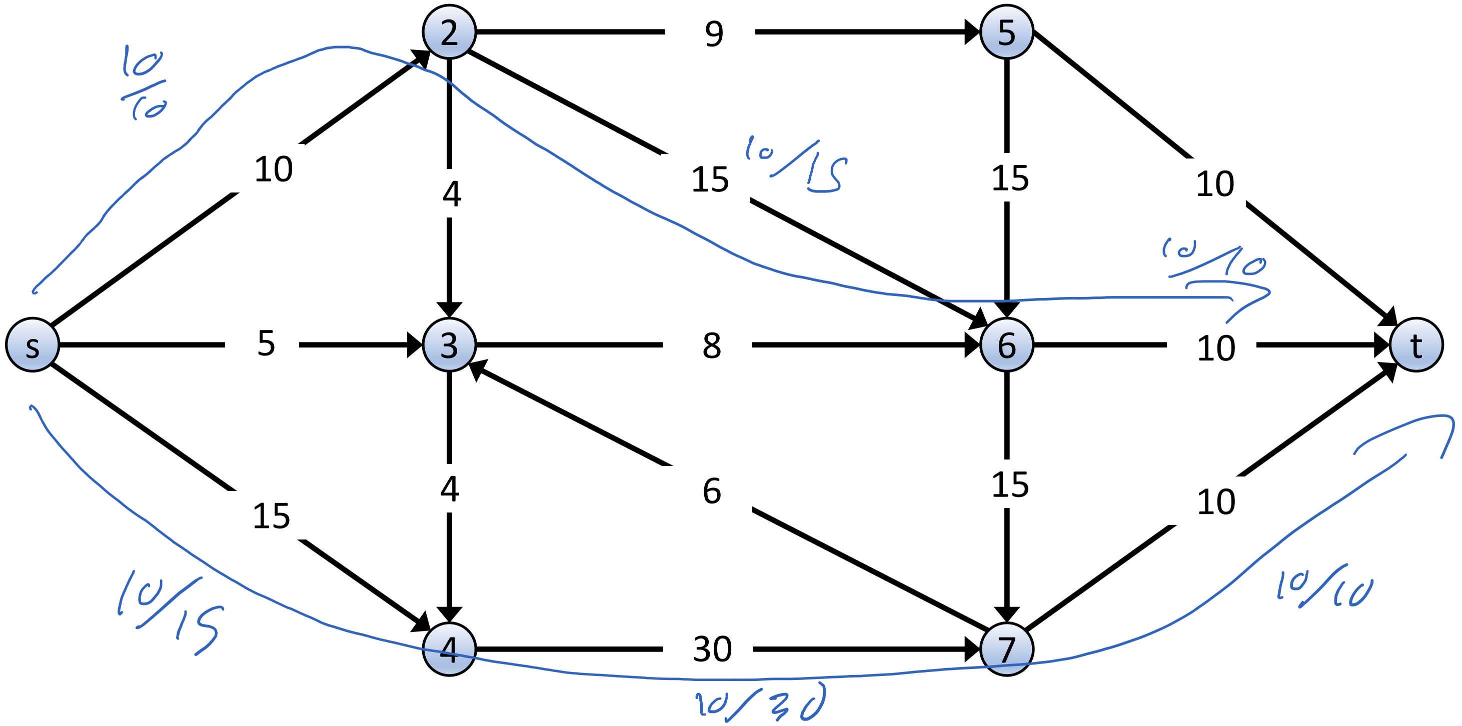
- Is there an integer maximum flow?



Maximum Flow Problem

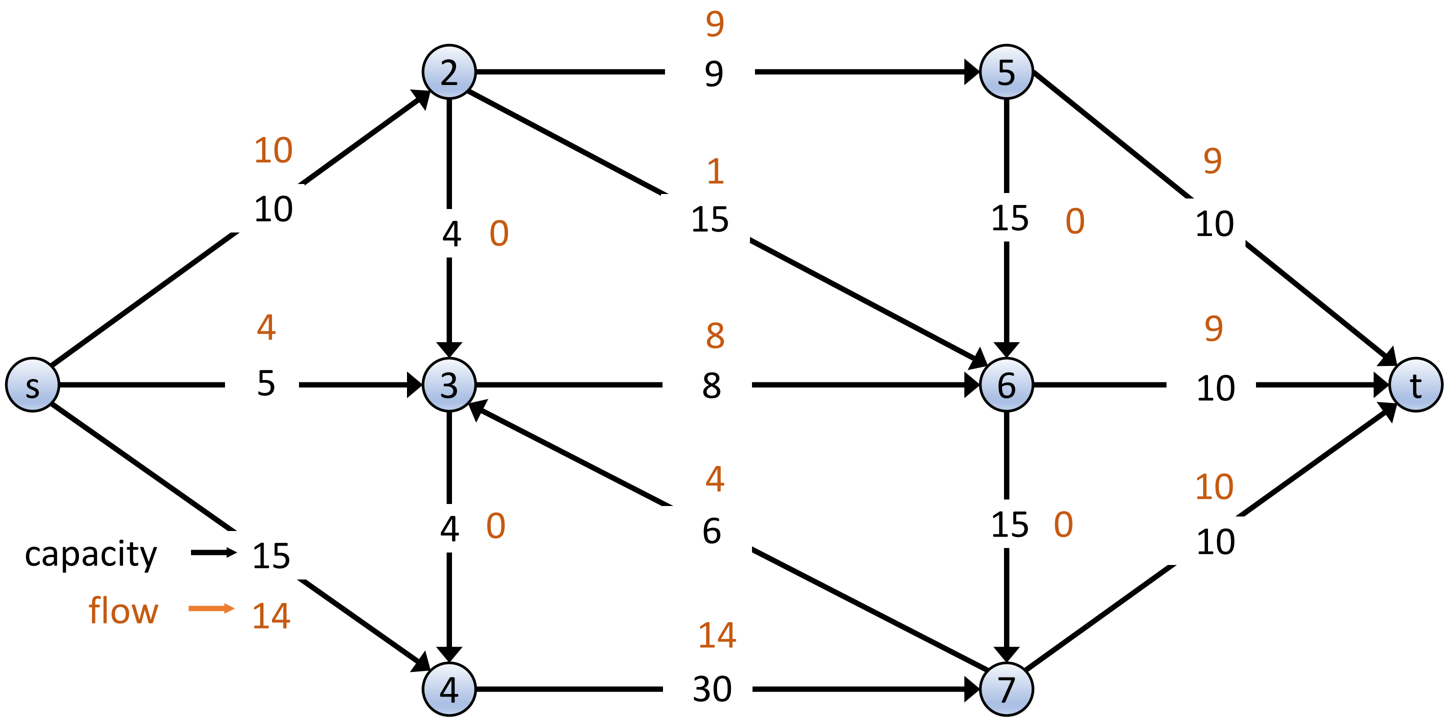
- Given G = (V,E,s,t,{c(e)}), find an s-t flow of maximum value

Idea:
Find a path from s-t
Max it out
Repeat



Maximum Flow Problem

- Given G = (V,E,s,t,{c(e)}), find an s-t flow of maximum value



Cuts

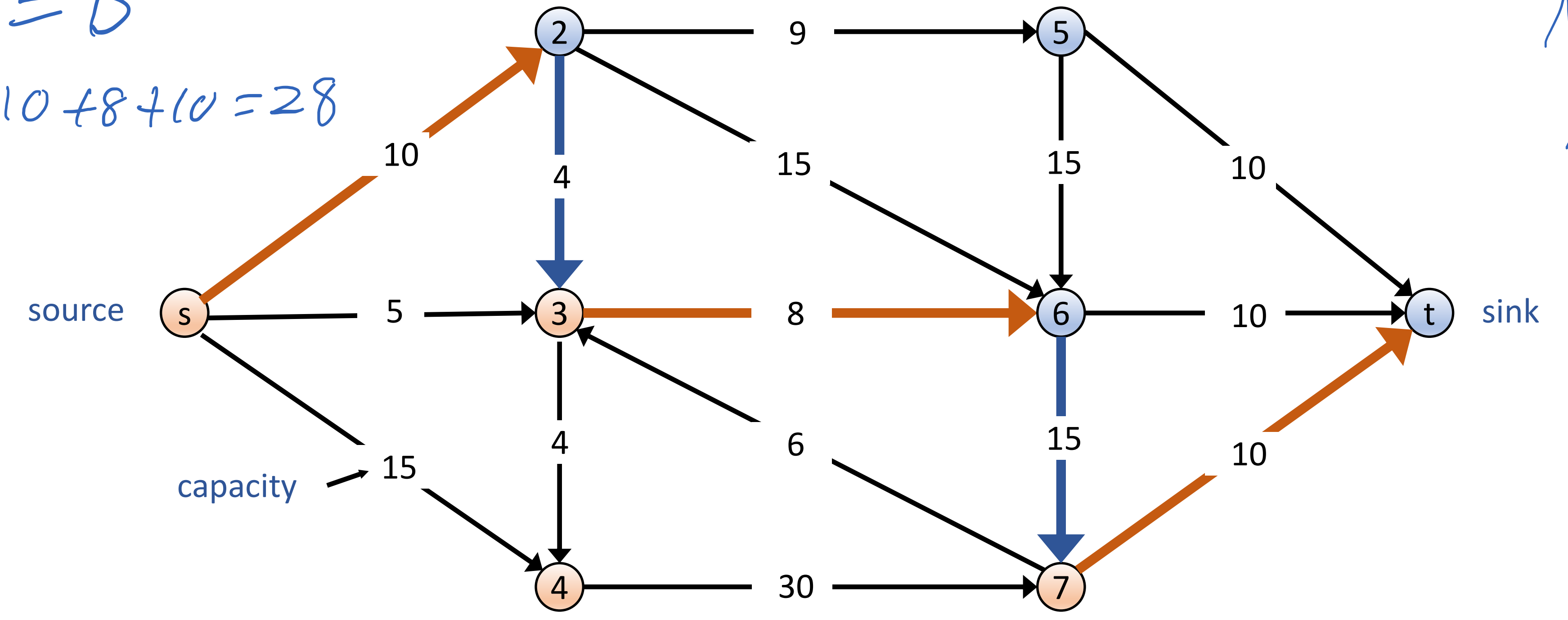
Every node is in A or B but not both

- An s-t cut is a partition (A, B) of V with $s \in A$ and $t \in B$

- The capacity of a cut (A,B) is $cap(A, B) = \sum_{e \text{ out of } A} c(e)$ (and into B)

Orange Nodes = A
Blue Nodes = B

Cap(A,B) = 10+8+10 = 28

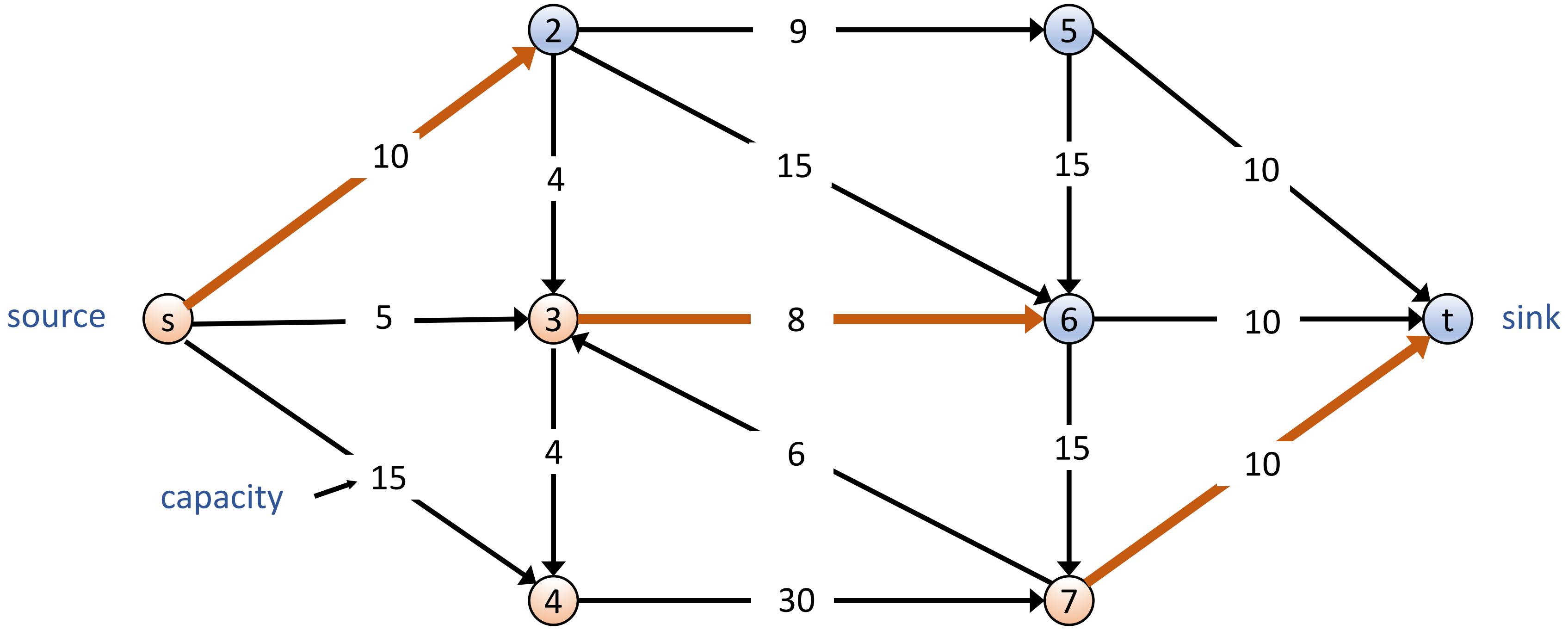


Note:

Max flow can be no more than 28

Minimum Cut problem

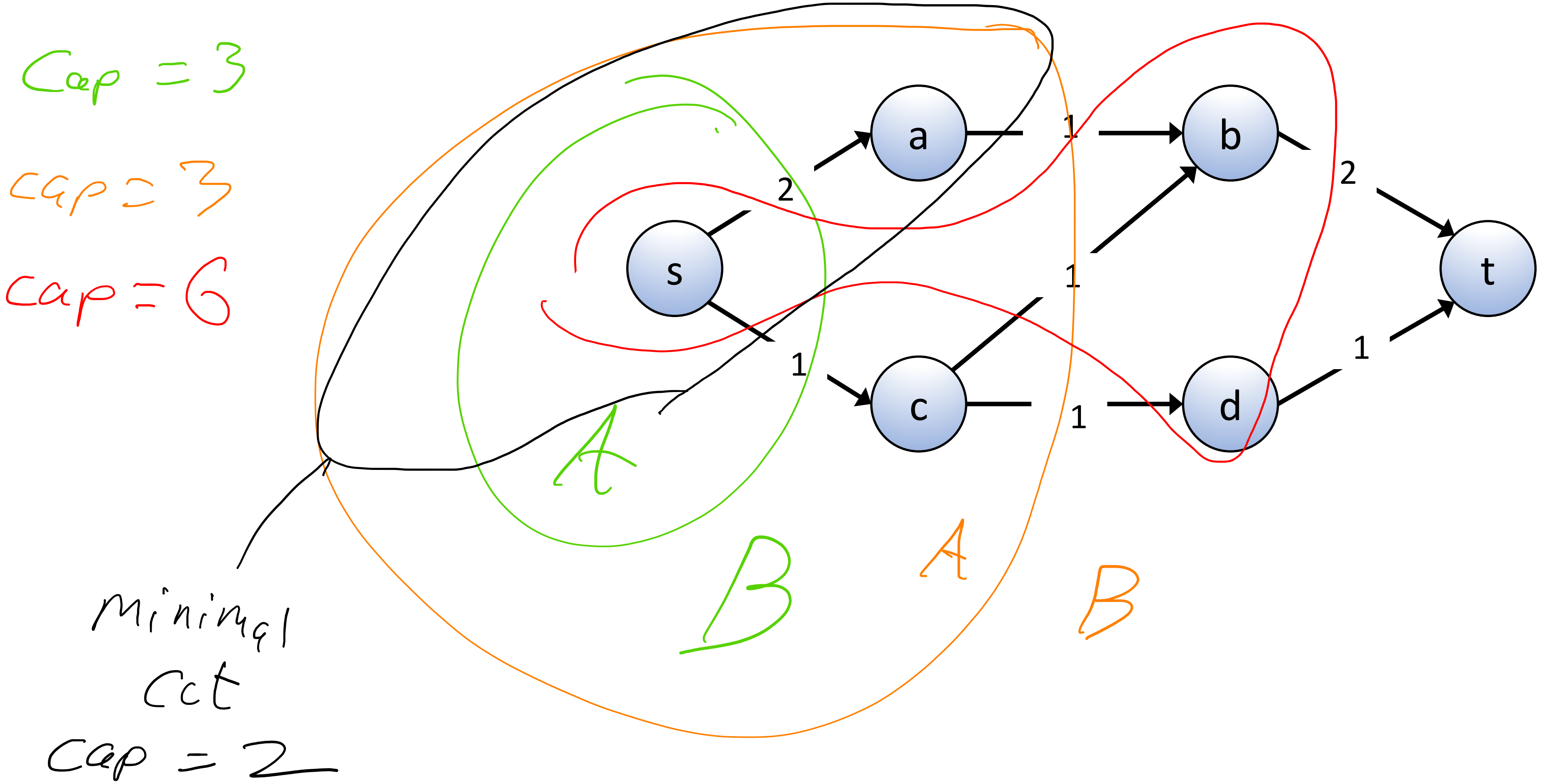
- Given G = (V,E,s,t,{c(e)}), find an s-t cut of minimum capacity



Minimum Cut Problem

- Given G = (V,E,s,t,{c(e)}), find an s-t cut of minimum capacity

- Find a minimum cut of this network

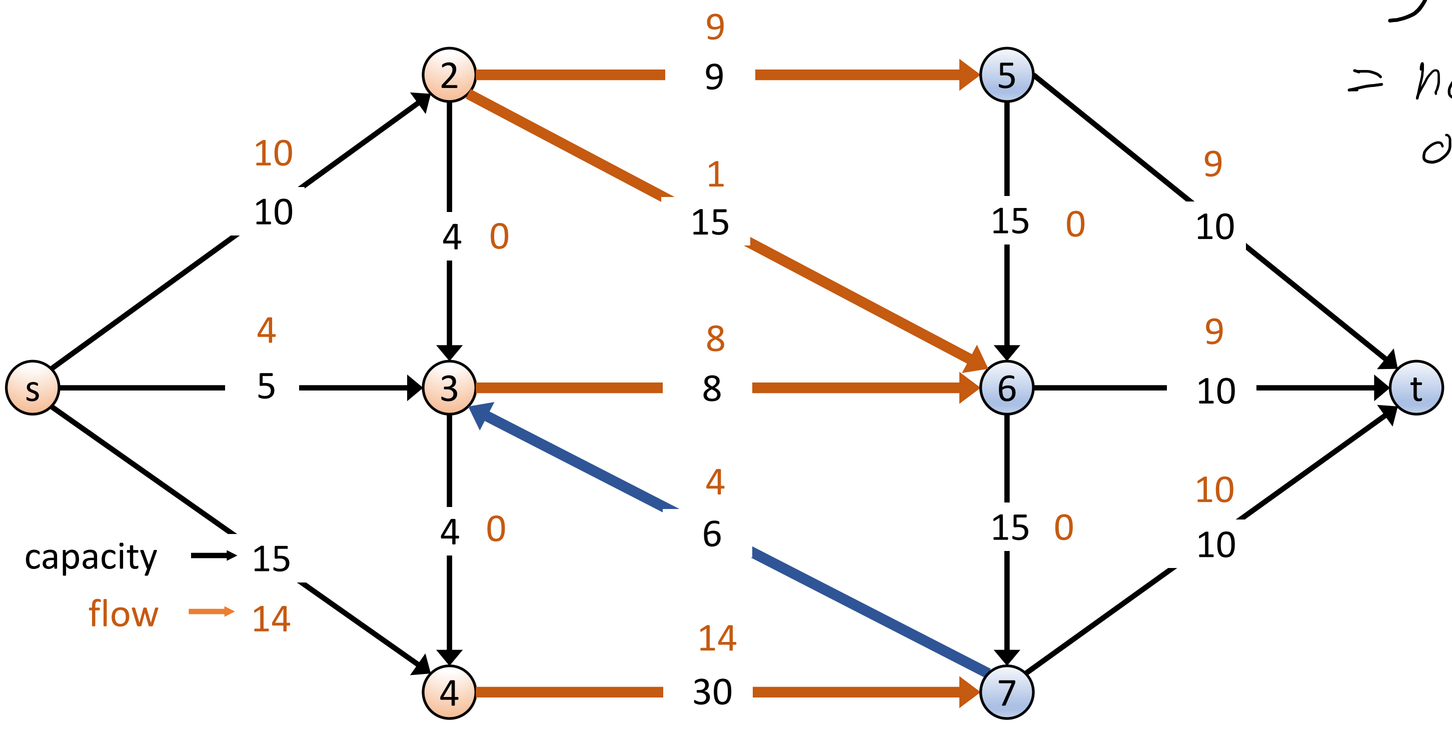


Flows vs. Cuts

- **Fact:** If f is any s-t flow and (A, B) is any s-t cut, then the net flow across (A, B) is equal to the amount leaving s

$$\sum_{e \text{ out of } A} f(e) - \sum_{e \text{ in to } A} f(e) = val(f)$$

~~net~~ amount of stuff flowing out of source = net amount of stuff leaving A



Max Flow Min Cut Duality

- **Weak Duality:** Let f be any s-t flow and (A, B) any s-t cut,

$$val(f) \leq cap(A, B)$$

- **Proof:**

$$Val(f) = \sum_{e \text{ out of } A} f(e) - \sum_{e \text{ into } A} f(e)$$

$$\leq \sum_{e \text{ out of } A} f(e) \quad \text{(non negativity)}$$

$$\leq \sum_{e \text{ out of } A} c(e) \quad \text{(capacity)}$$

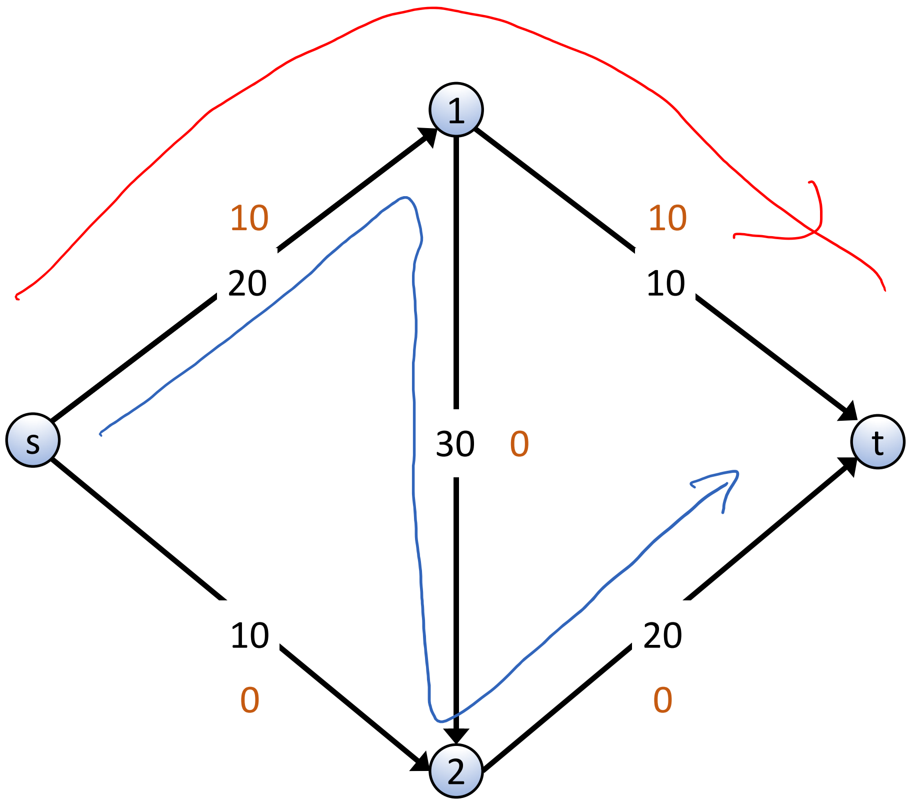
$$= cap(A, B) \quad \text{(defn)}$$

Augmenting Paths

path where all pipes are strictly below capacity.

- Given a network $G = (V, E, s, t, \{c(e)\})$ and a flow f, an **augmenting path** P is an $s \to t$ path such that $f(e) < c(e)$ for every edge $e \in P$

Could send a nonzero amount down this path

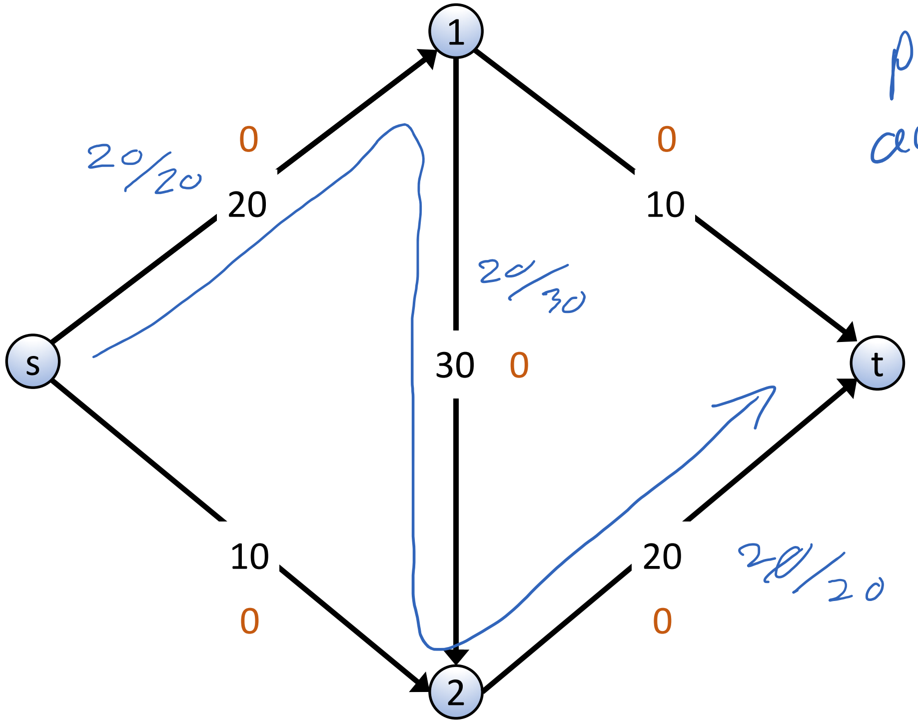


— Augmenting Path

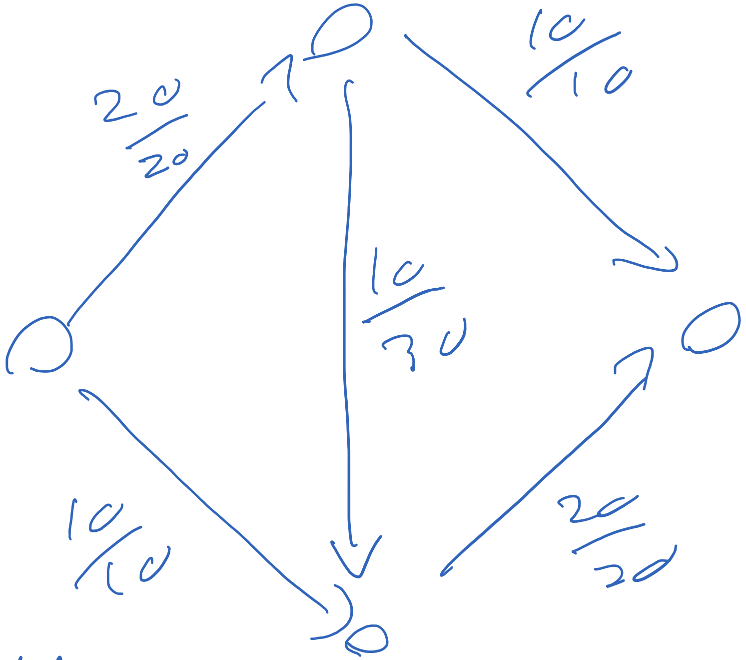
— Not Augmenting

Greedy Max Flow

- Start with $f(e) = 0$ for all edges $e \in E$
- Find an **augmenting path** P, max it out
- Repeat until you get stuck



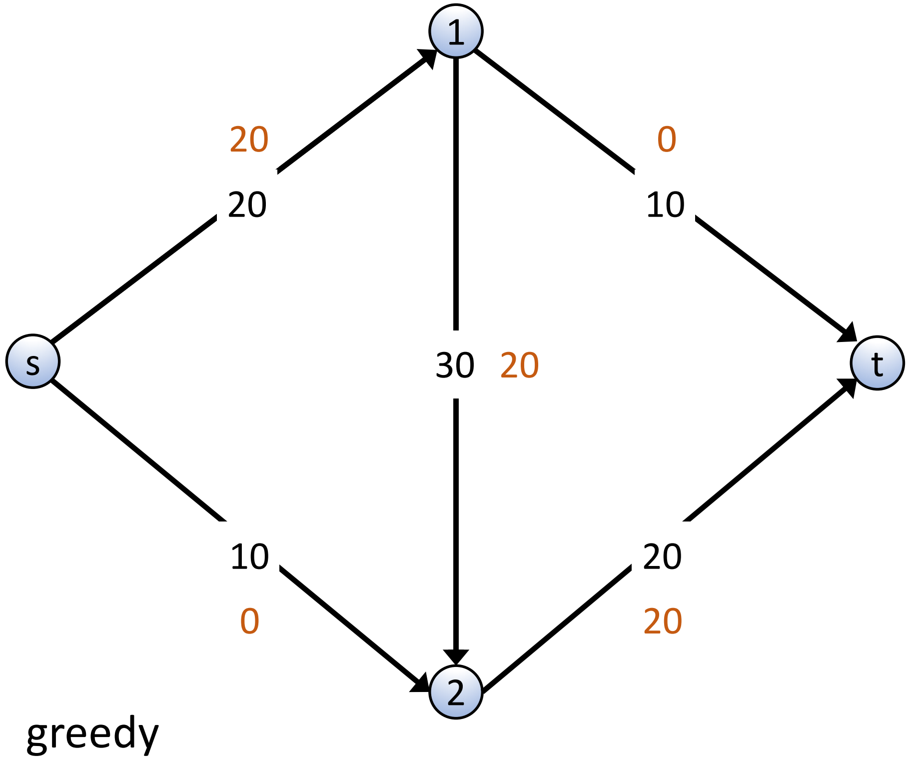
After this path, no more augmenting paths

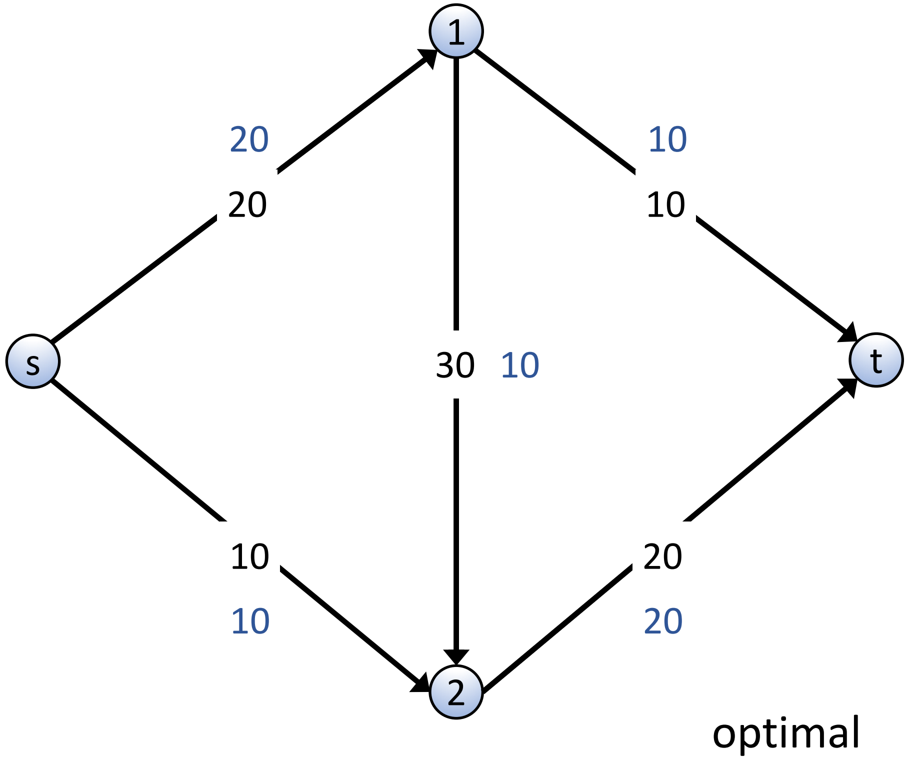


Better: 30 flow

Does Greedy Work?

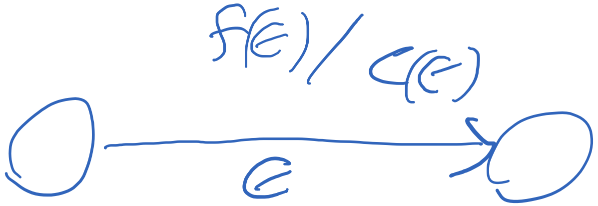
- Greedy gets stuck before finding a max flow
- How can we get from our solution to the max flow?



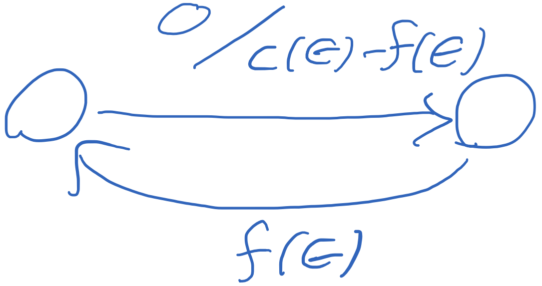


Residual Graphs

- Original edge: $e = (u, v) \in E$.
  - Flow $f(e)$, capacity $c(e)$



- Residual edge
  - Allows "undoing" flow
  - $e = (u, v)$ and $e^R = (v, u)$.
  - Residual capacity



- Residual graph $G_f = (V, E_f)$
  - Edges with positive residual capacity.
  - $E_f = \{e : f(e) < c(e)\} \cup \{e^R : f(e) > 0\}$.

Edges not maxed out

Edges which are from decisions I can take back

Augmenting Paths in Residual Graphs

- Let G_f be a **residual graph**
- Let P be an augmenting path in the **residual graph**
- **Fact:** $f' = \text{Augment}(G_f, P)$ is a valid flow

"max out this path"

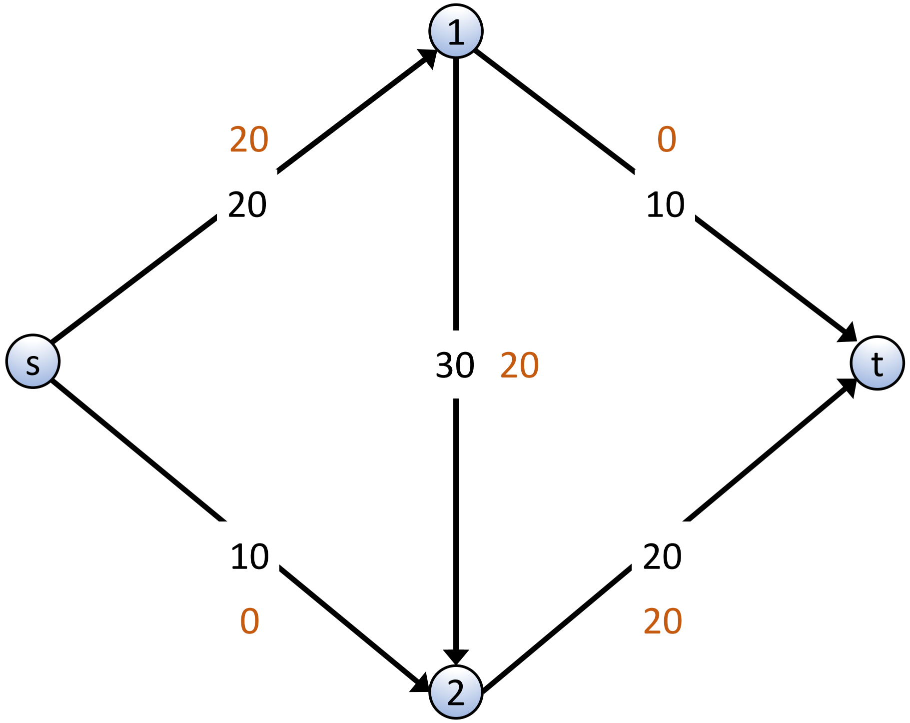
```
Augment(Gf, P)
    b ← the minimum (residual) capacity of an edge in P
    for e ∈ P
        if e ∈ E:    f(e) ← f(e) + b
        else:        f(e) ← f(e) - b
    return f
```

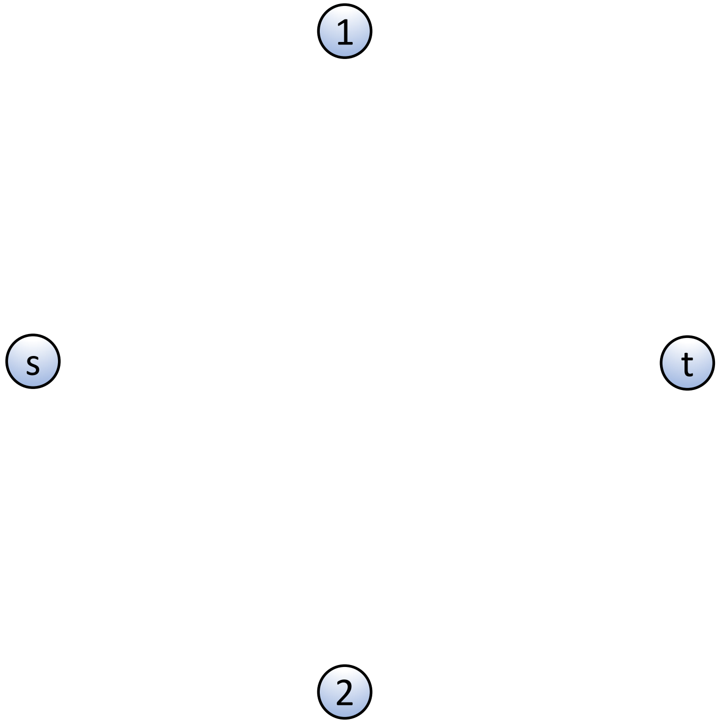
— was an original edge. So additional flow is added to what is flowing already

not an original edge, it is a residual edge
take back b units

Ford-Fulkerson Algorithm

- Start with $f(e) = 0$ for all edges $e \in E$
- Find an **augmenting path** P in the **residual graph**
- Max it out
- Repeat until you get stuck





Ford-Fulkerson Algorithm

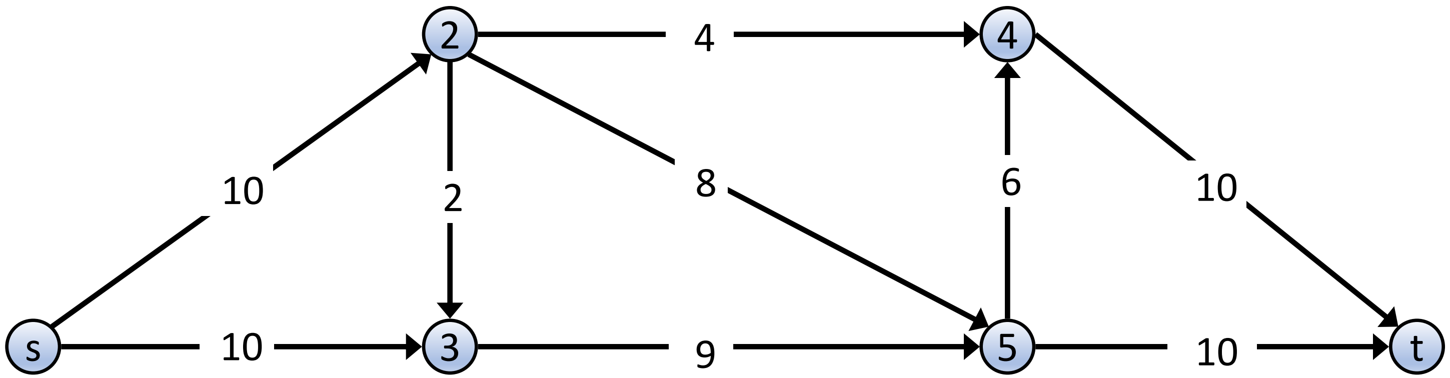
```
FordFulkerson(G,s,t,{c})
    for e ∈ E: f(e) ← 0
    Gf is the residual graph
                        augmenting
    while (there is an ^ s-t path P in Gf)
        f ← Augment(Gf,P)
        update Gf   (including residuals)

    return f
```

```
Augment(Gf, P)
    b ← the minimum capacity of an edge in P
    for e ∈ P
        if e ∈ E:   f(e) ← f(e) + b
        else:       f(e) ← f(e) - b
    return f
```

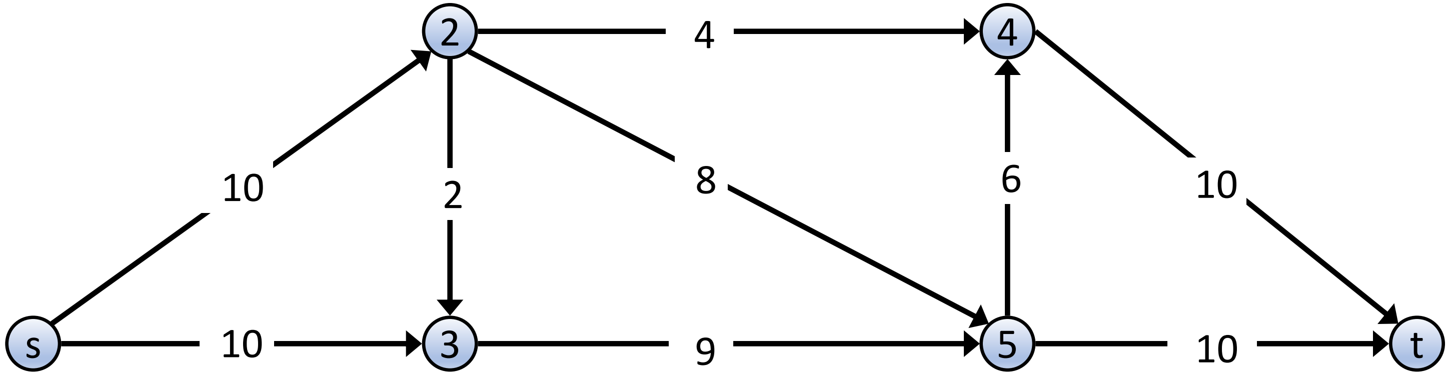
Ford-Fulkerson Demo

G:



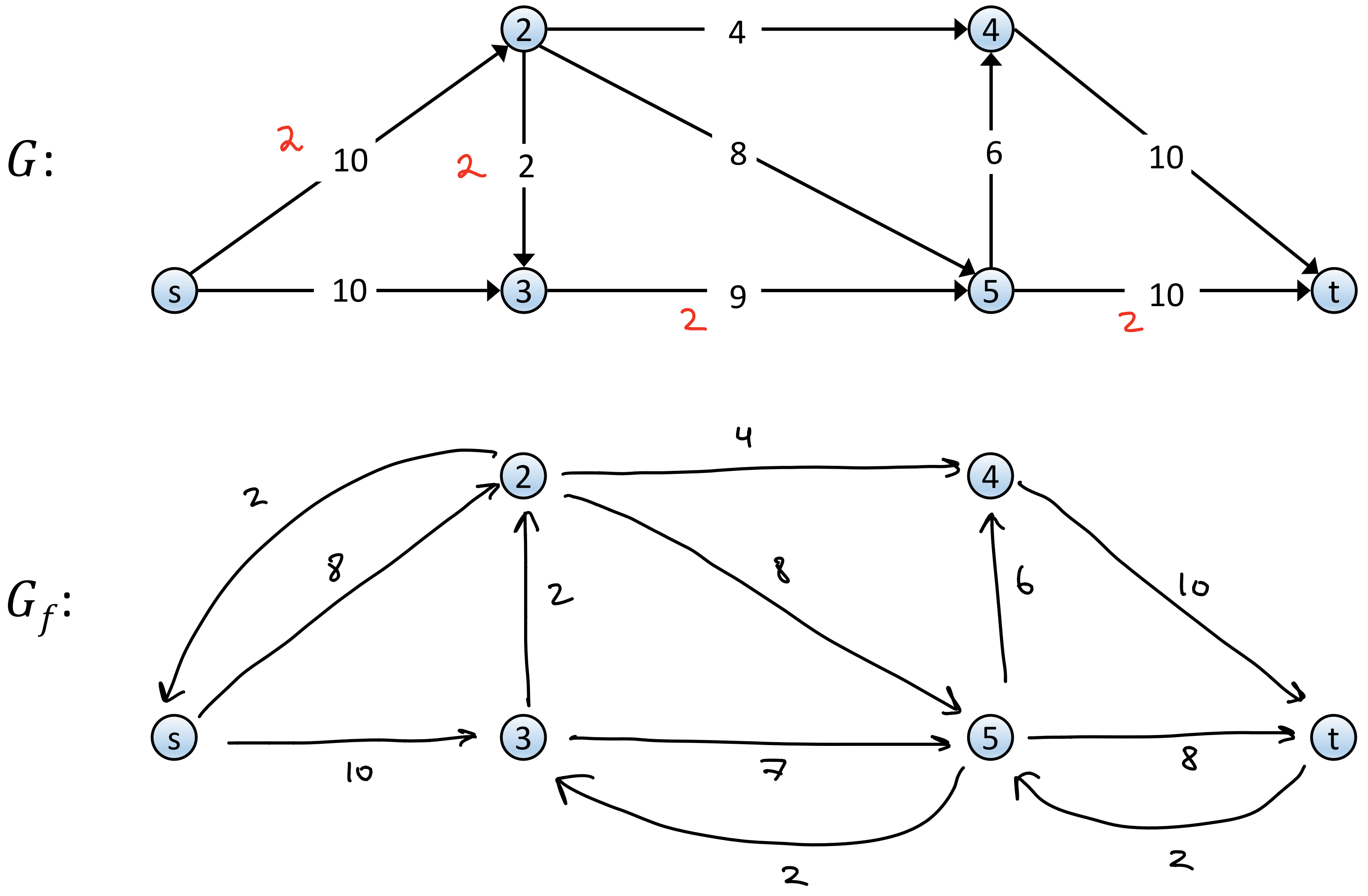
Ford-Fulkerson Demo

$G:$



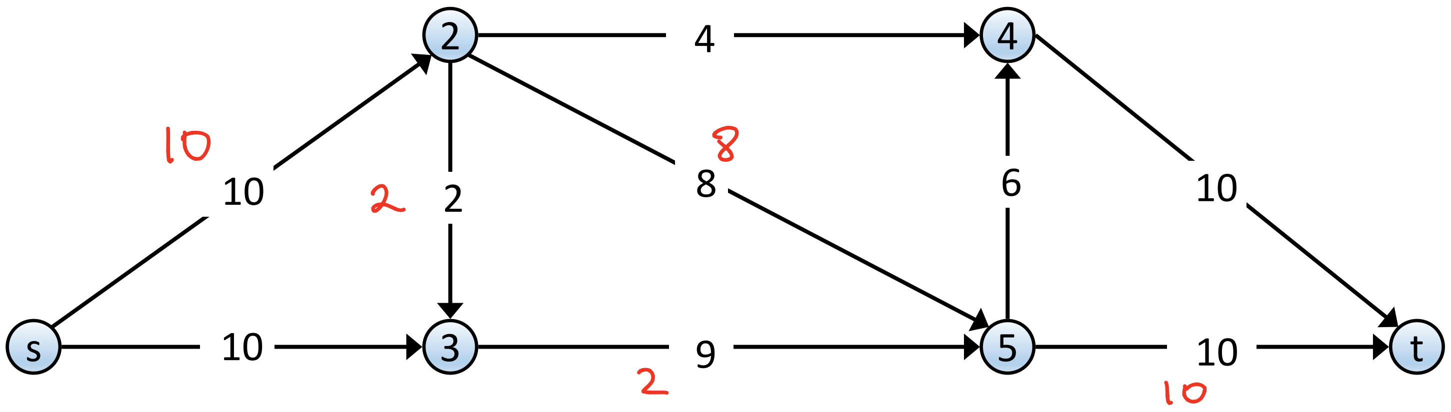
$G_f:$

Ford-Fulkerson Demo

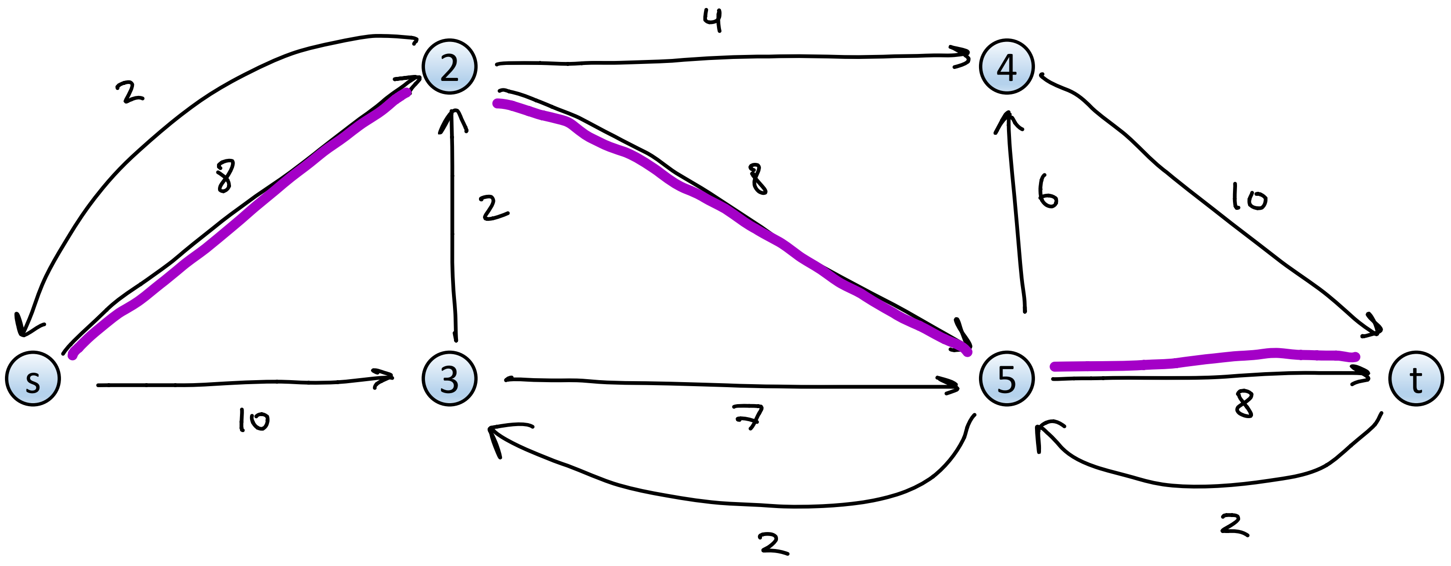


Ford-Fulkerson Demo

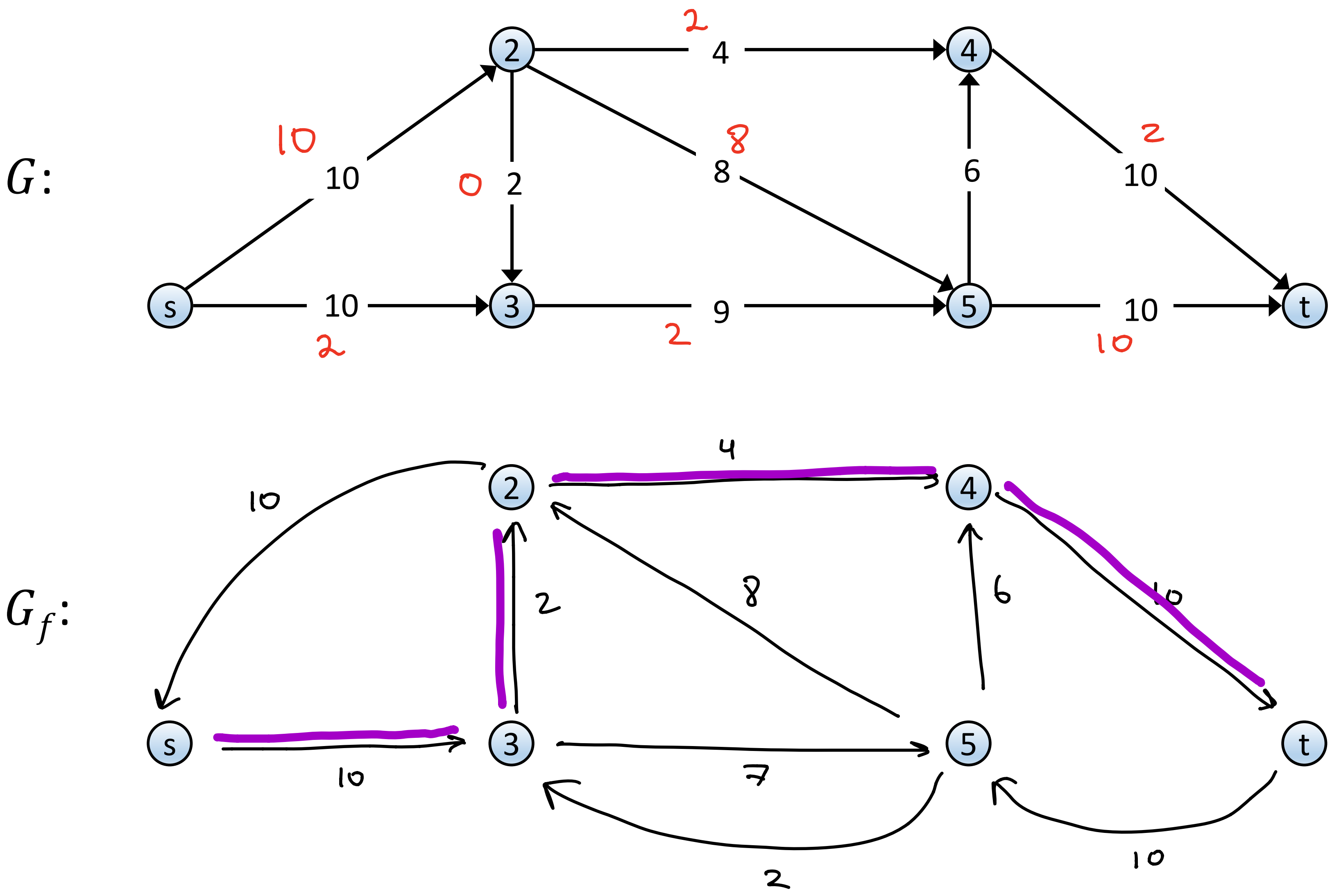
G:



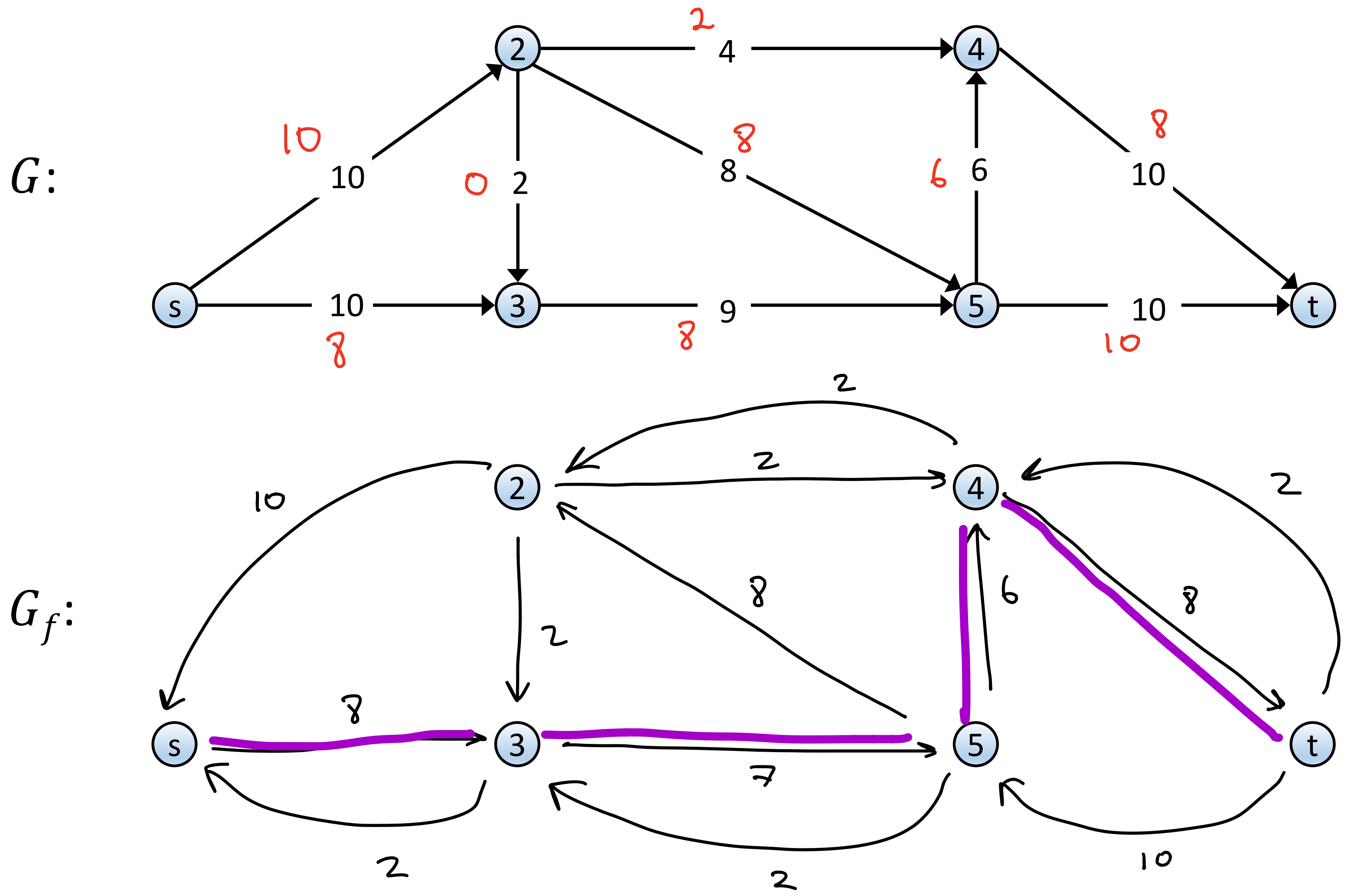
G_f:



Ford-Fulkerson Demo

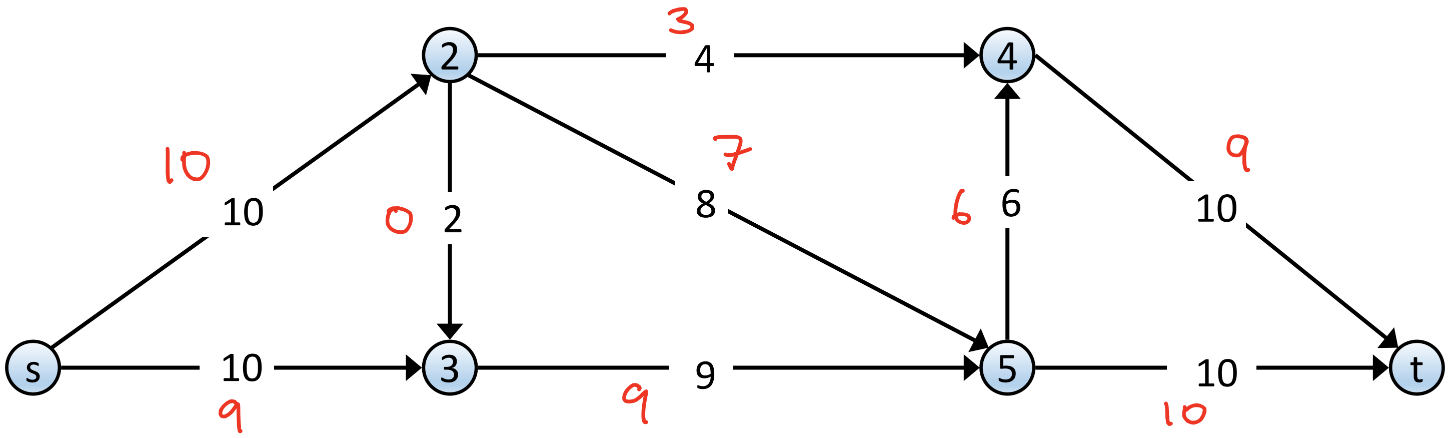


Ford-Fulkerson Demo

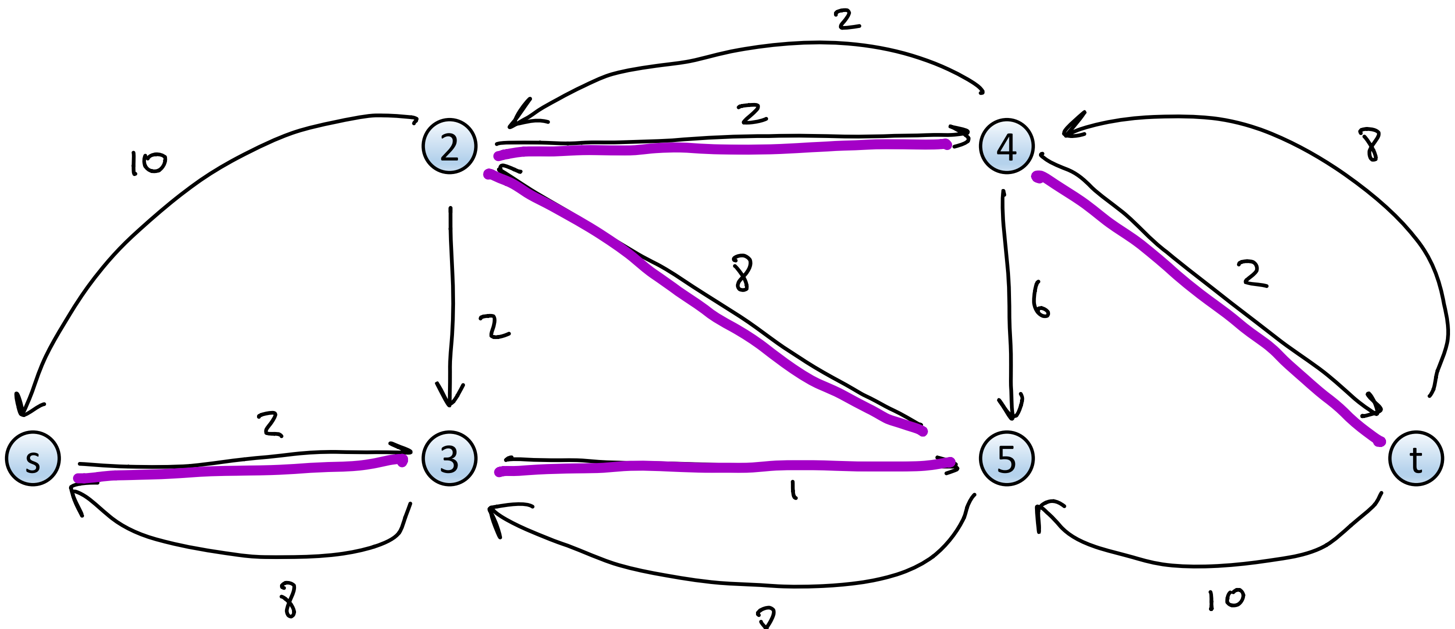


Ford-Fulkerson Demo

G:



G_f:

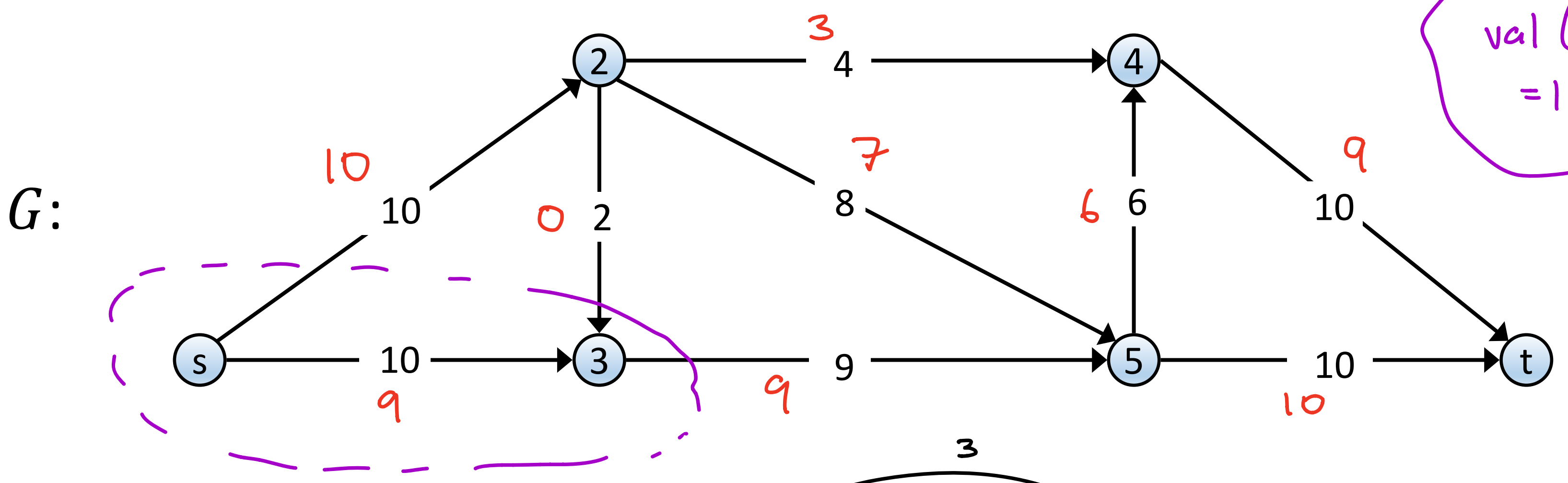


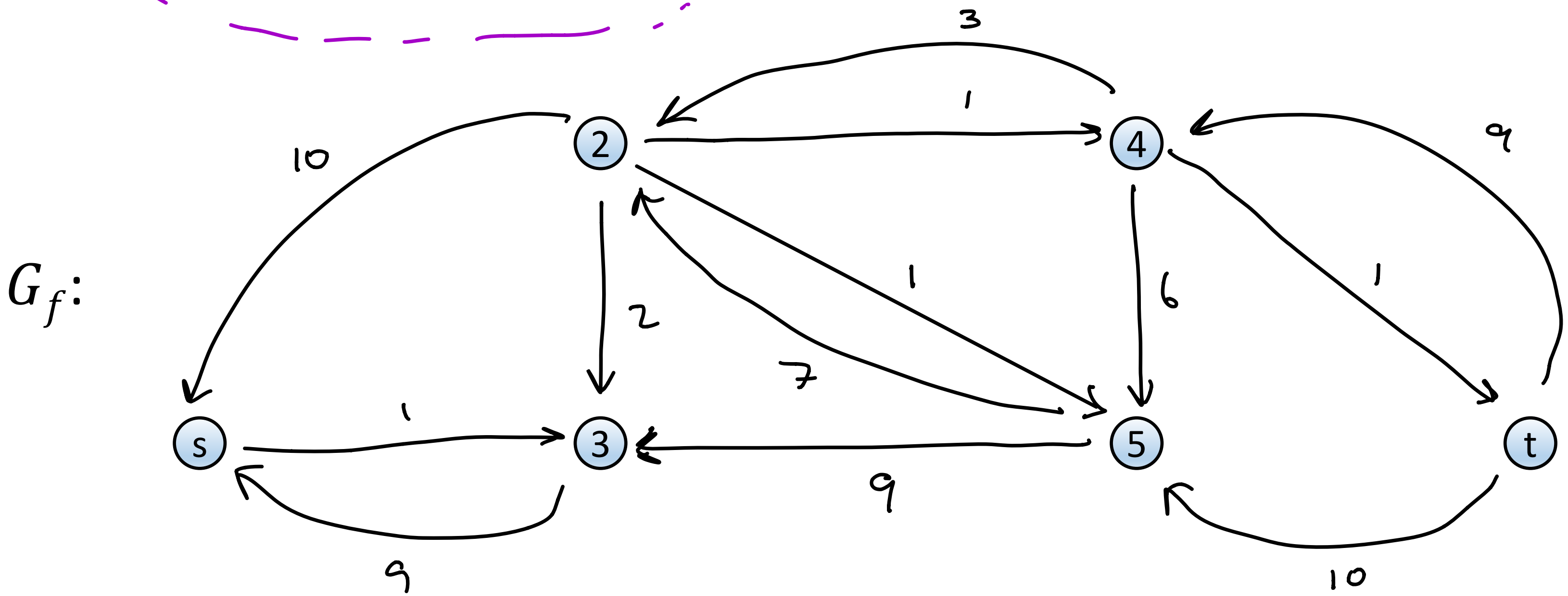
Ford-Fulkerson Demo

A = { s, 3 }
B = { 2, 4, 5, t }

cap(A,B) = 19

val(f) = 19





What do we want to prove?

- FF Terminates
- FF finds a maximum s-t flow
- There is always a cut (A,B) such that val(f) = cap(A,B)

Ford-Fulkerson Algorithm – Run Time

```
FordFulkerson(G,s,t,{c})
    for e ∈ E: f(e) ← 0
    Gf is the residual graph

    while (there is an s-t path P in Gf)
        f ← Augment(Gf,P)
        update Gf

    return f
```

```
Augment(Gf, P)
    b ← the minimum capacity of an edge in P
    for e ∈ P
        if e ∈ E:   f(e) ← f(e) + b
        else:       f(e) ← f(e) - b
    return f
```

Running Time of Ford-Fulkerson

- For **integer capacities**, $\leq val(f^*)$ augmentation steps

- Can perform each augmentation step in $O(m)$ time
  - find augmenting path in $O(m)$
  - augment the flow along path in $O(n)$
  - update the residual graph along the path in $O(n)$

- For integer capacities, FF runs in $O\big(m \cdot val(f^*)\big)$ time
  - $O(mn)$ time if all capacities are $c_e = 1$
  - $O(mnC_{\max})$ time for any integer capacities
  - Problematic when capacities are large

Correctness of Ford-Fulkerson

- **Theorem:** f is a maximum s-t flow if and only if there is no augmenting s-t path in G_f

- **(Strong) MaxFlow-MinCut Duality:** The value of the max s-t flow equals the capacity of the min s-t cut

- We'll prove that the following are equivalent for all f
  1. There exists a cut (A, B) such that $val(f) = cap(A, B)$
  2. Flow f is a maximum flow
  3. There is no augmenting path in G_f

Optimality of Ford-Fulkerson

- **Theorem:** the following are equivalent for all f
  1. There exists a cut (A, B) such that $val(f) = cap(A, B)$
  2. Flow f is a maximum flow
  3. There is no augmenting path in G_f

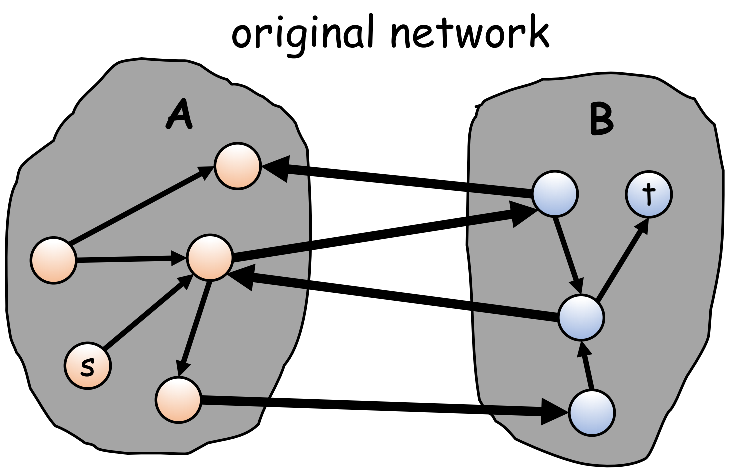
Optimality of Ford-Fulkerson

- **(3 → 1)** If there is no augmenting path in G_f, then there is a cut (A, B) such that $val(f) = cap(A, B)$
  - Let A be the set of nodes reachable from s in G_f
  - Let B be all other nodes

Optimality of Ford-Fulkerson

- **(3 → 1)** If there is no augmenting path in G_f, then there is a cut (A, B) such that $val(f) = cap(A, B)$
  - Let A be the set of nodes reachable from s in G_f
  - Let B be all other nodes
  - **Key observation:** no edges in G_f go from A to B

- If e is $A \rightarrow B$, then $f(e) = c(e)$
- If e is $B \rightarrow A$, then $f(e) = 0$



Summary

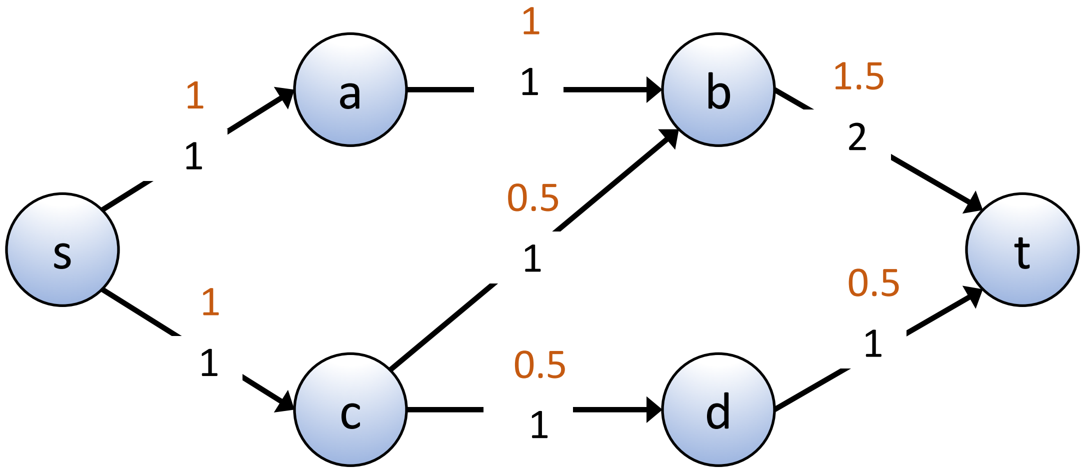
- **The Ford-Fulkerson Algorithm solves maximum s-t flow**
  - Running time $O\left(m \cdot val(f^*)\right)$ in networks with integer capacities
  - Space $O(n+m)$

- **MaxFlow-MinCut Duality:** The value of the maximum s-t flow equals the capacity of the minimum s-t cut
  - If f^* is a maximum s-t flow, then the set of nodes reachable from s in G_{f^*} gives a minimum cut
  - Given a max-flow, can find a min-cut in time $O(n+m)$

- **Every graph with integer capacities has an integer maximum flow**
  - Ford-Fulkerson will return an integral maximum flow

Ask the Audience

- Is this a maximum flow?



- Is there an **integer maximum flow**?
- Does every graph with integer capacities have an integer maximum flow?

Summary

- **The Ford-Fulkerson Algorithm solves maximum s-t flow**
  - Running time $O(m \cdot val(f^*))$ in networks with integer capacities
  - Space $O(n+m)$

- **MaxFlow-MinCut Duality:** The value of the maximum s-t flow equals the capacity of the minimum s-t cut
  - If f^* is a maximum s-t flow, then the set of nodes reachable from s in G_{f^*} gives a minimum cut
  - Given a max-flow, can find a min-cut in time $O(n+m)$

- **Every graph with integer capacities has an integer maximum flow**
  - Ford-Fulkerson will return an integer maximum flow