

6 Homework

Due: Thursday, February 28, 2013.

Instructions

- Please, review the homework grading policy outlined in the course information page.
- On the *first page* of your solution write-up you *must* make explicit which problems are to be graded for regular credit, which problems are to be graded for extra credit, and which problems you did not attempt. Use a table that looks like this:

Problem	1	2	3	4	5	6	7	8	9	...
Credit	RC	RC	RC	EC	RC	EC	NA	NA	EC	...

where “RC” denotes “regular credit”, “EC” denotes “extra credit”, and “NA” denotes “not attempted”. Failure to include such a table will result in an arbitrary set of problems being graded for regular credit, no problems being graded for extra credit, and a 5% penalty assessment.

- You must also write down with whom you worked on the assignment. If this varies from problem to problem, write down this information separately with each problem.

Problems

Required: 5 of the following 6 problems

Points: 20 points per problem

1. Do Problem 2.30 parts a and d

Use the pumping lemma to show that the following languages are not context free.

(a) (a) $\{0^n 1^n 0^n 1^n \mid n \geq 0\}$

(b) (d) $\{t_1 \# t_2 \# \dots \# t_k \mid k \geq 2, \text{ each } t_i \in \{a, b\}^*, \text{ and } t_i = t_j \text{ for some } i \neq j\}$

2. Do Problem 2.31 or the Problem 2.32

(2.31) Let B be the language of all palindromes over $\{0,1\}$ containing equal number of 0s and 1s. Show that B is not context-free.

(2.32) Let $\Sigma = \{1,2,3,4\}$ and $C = \{w \in \Sigma^* \mid \text{in } w, \text{ the number of 1s equals the number of 2s, and the number of 3s equals the number of 4s}\}$. Show that C is not context-free.

3. Do Problem 2.34:

Consider the language $B = L(G)$, where $G = (V, \Sigma, R, R)$ is the following grammar: $V = \{S, T, U\}$; $\Sigma = \{0, \#\}$; and R is the set of rules:

$$S \rightarrow TT|U$$

$$T \rightarrow 0T|T0|\#$$

$$U \rightarrow 0U00|\#$$

The pumping lemma for context-free languages, Theorem 2.34, states the existence of a pumping length p for B .

What is the minimum p that works in the pumping lemma for the language B . Justify your answer.

4. Do Problem 2.44

If A and B are languages, define $A \diamond B = \{xy \in A \text{ and } y \in B \text{ and } |x| = |y|\}$. Show that if A and B are regular languages, then $A \diamond B$ is a CFL.

5. Do Problem 2.54 (skipping part b)

2.54 Let G be the following grammar:

$$S \rightarrow T \dagger$$

$$T \rightarrow TaTb|TbTa|\epsilon$$

- (a) Show that $L(G) = \{w \dagger \mid w \text{ contains equal number of } a\text{'s and } b\text{'s}\}$. Use proof by induction on the length of w .
- (c) Describe the DPDA that recognizes $L(G)$.

6. Do Problem 3.2 parts (b) and (d) and Problem 3.5:

- (a) 3.2 This exercise concerns TM M_1 , whose description and state diagram appear in Example 3.9. In each of the parts, give the

sequence of configurations that M_1 enters when started on the indicated input string:

- (b) 1#1
 - (d) 10#11
- (b) 3.5 Examine the formal definition of a Turing Machine to answer the following questions, and explain your reasoning:
- (a) Can a Turing machine ever write the blank symbol \sqcup on its tape?
 - (b) Can the tape alphabet Γ be the same as the input alphabet Σ ?
 - (c) Can a Turing machine's head *ever* be in the same location in two successive steps?
 - (d) Can a Turing machine contain just a single state?