

### 3 Homework

**Due:** Thursday, January 31, 2013.

#### Instructions

- Please, review the homework grading policy outlined in the course information page.
- On the *first page* of your solution write-up you *must* make explicit which problems are to be graded for regular credit, which problems are to be graded for extra credit, and which problems you did not attempt. Use a table that looks like this:

Problem	1	2	3	4	5	6	7	8	9	...
Credit	RC	RC	RC	EC	RC	EC	NA	NA	EC	...

where “RC” denotes “regular credit”, “EC” denotes “extra credit”, and “NA” denotes “not attempted”. Failure to include such a table will result in an arbitrary set of problems being graded for regular credit, no problems being graded for extra credit, and a 5% penalty assessment.

- You must also write down with whom you worked on the assignment. If this varies from problem to problem, write down this information separately with each problem.

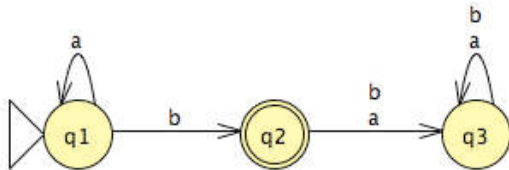
#### Problems

**Required:** 5 of the following 7 problems

**Points:** 20 points per problem

1. Give the regular expressions generating the languages from Problem 1 in Homework 2, i.e.:
  - (a)  $\{w \mid w \text{ the length of } w \text{ is at most } 5\}$
  - (b)  $\{w \mid w \text{ is any string except } aa \text{ and } aaa\}$
  - (c)  $\{w \mid w \text{ every odd position of } w \text{ is } a\}$

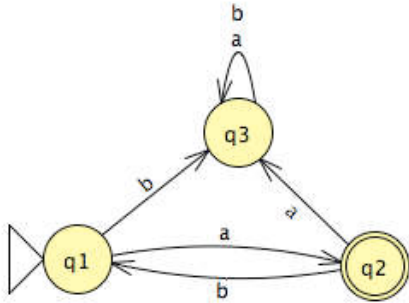
- (d)  $\{w \mid w \text{ contains at least two } a\text{s and at most one } b\}$
- (e)  $\{w \mid w \text{ begins with } a \text{ and ends with } b, \text{ or begins with } b \text{ and ends with } a\}$
2. Give the regular expressions generating the languages from Problem 2 in Homework 2, i.e.:
- (a)  $\{w \mid w \text{ begins with } 1, \text{ and ends with } 0\}$
- (b)  $\{w \mid w \text{ contains at least three } 1\text{s}\}$
- (c)  $\{w \mid w \text{ contains substring } 0101\}$  (i.e.,  $w = x0101y$ )
- (d)  $\{w \mid w \text{ has length at least } 3 \text{ and its third symbol is a } 0\}$
- (e)  $\{w \mid w \text{ does not contain the substring } 011\}$
3. Use the procedure described in Lemma 1.60 in the text to convert the following DFA to a regular expression in two different ways:
- eliminating first state 3, then state 2, then state 1
  - eliminating first state 1, then state 2, then state 3



Show the resulting GNFA after each step, and do not try to simplify your answer (except for eliminating all instances of  $\emptyset$  in unions and all instances of  $\epsilon$  in concatenations).

4. Use the procedure described in Lemma 1.60 in the text to convert the following DFA to a regular expression in two different ways:

- eliminating first state 3, then state 2, then state 1
- eliminating first state 1, then state 2, then state 3



Show the resulting GNFA after each step, and do not try to simplify your answer (except for eliminating all instances of  $\emptyset$  in unions and all instances of  $\epsilon$  in concatenations).

- For any string  $w = w_1w_2\dots w_n$ , the reverse of  $w$ , written  $w^R$ , is the string  $w$  in reverse order,  $w_n\dots w_2w_1$ . For any language  $A$ , let  $A^R = \{w^R \mid w \in A\}$ . Show that if  $A$  is regular, then so is  $A^R$ .
- Problem 1.42* For languages  $A$  and  $B$ , let the *perfect shuffle* of  $A$  and  $B$  be the language  $\{w \mid w = a_1b_1\dots a_kb_k, \text{ where } a_1\dots a_k \in A \text{ and } b_1\dots b_k \in B, \text{ for each } a_i, b_i \in \Sigma\}$ .  
Show that the class of regular languages is closed under *perfect shuffle*.
- Prove or disprove the following:
  - Every subset of a regular language is a regular language.
  - Every subset of a nonregular language is a nonregular language.
  - If  $A$  is a regular language and  $B$  is a language such that  $AB$  is regular, then  $B$  is regular.
  - If  $A$  is a regular language and  $B$  is a language such that  $A - B$  is regular, then  $B$  is regular.
  - For any language  $A$  and its complement  $A'$  the language  $A \cup A'$  is regular.