

Problem Set 1 (due Friday, January 25)

1. (7 points) [Sipser 1.48] A regular language

Let $\Sigma = \{0, 1\}$ and let

$$D = \{w \mid w \text{ contains an equal number of occurrences of the substrings } 01 \text{ and } 10\}.$$

Thus, $101 \in D$ because 101 contains one 10 and one 01, but $1010 \notin D$ because 1010 contains two 10s and one 01. Show that D is a regular language.

2. ($4 \times 5 = 20$ points) A non-regular language

Consider the language $F = \{a^i b^j c^k \mid i, j, k \geq 0 \text{ and if } i = 1 \text{ then } j = k\}$.

- Show that F is not regular,
- Show that F satisfies the conditions of the pumping lemma.
- Explain why parts (a) and (b) do not contradict the pumping lemma.
- Show that F is context-free.

(Parts (a) through (c) are from [Sipser 1.54].)

3. (8 points) Designing a Turing machine

Design a Turing machine for the language of all strings with an equal number of 0s and 1s. The input alphabet is $\{0, 1, \}$.

4. ($10 + 3 + 2 = 15$ points) Input-read-only Turing machines

Define the *input-read-only* Turing machine as a single tape deterministic Turing machine that is not allowed to write on the portion of the tape containing the input string.

- Prove that input-read-only Turing machines can only recognize regular languages.
- Does the claim of part (a) hold if we allow the input-read-only Turing machine to have *two heads*? Briefly justify your answer. Note that we maintain the restrictions of a single tape and that the portion of the input cannot be written into.

The definition of a two-head Turing machine is fairly natural. The transition function of a two-head Turing machine takes as input a state and two symbols (one corresponding to each head) and returns a state, two symbols to be written (one corresponding to each head), and two moves (one corresponding to each head). We can assume an arbitrary tie-breaking mechanism for the scenario when two heads attempt to write on the same cell. Initially, both the heads are at the leftmost end of the tape.

- (c) Does the claim of part (a) hold if we allow the input-read-only Turing machine to have two tapes? Briefly justify your answer.

5. ($4 \times 5 = 20$ points) Classification of languages

Each of the following four parts gives a definition, description, or some properties of a language. In each case, tell whether the language is:

- A regular
- B context-free but not regular
- C Turing-decidable but not context-free
- D not Turing-decidable
- E there is insufficient information to tell.

Justify your answers with proofs, constructions, algorithms, or examples as needed.

If, for example, you decide that a language is context-free but not regular, you must give a context-free grammar that generates the language or a pushdown automaton that accepts the language and also prove that the language is not regular.

- (a) $L = \{0^i 1^j 2^k \mid \text{either } i + j = k \text{ or } i + k = j, \text{ and } i, j, k \geq 0\}$. The input alphabet is $\{0, 1, 2\}$.
- (b) $L = \{0^p \mid p \text{ is prime}\}$. The input alphabet is $\{0\}$.
- (c) $L = \{01^m 0 \mid m \text{ is a multiple of } 7\}$. The input alphabet is $\{0, 1\}$.
- (d) L is the set of balanced strings of three types of parentheses, $()$ and $[]$ and $\{, \}$. (For example, $\{[(())]\{\}\}$ is balanced, while $\{([\])\}$ is not balanced.)