

How do you plan in high dimensional state spaces?

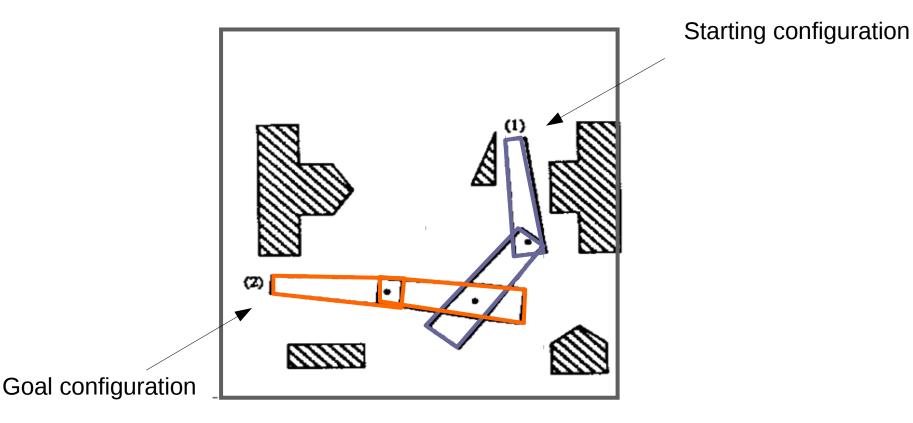
Problem we want to solve

<u>Given</u>:

- a point-robot (robot is a point in space)
- a start and goal configuration

<u>Find</u>:

 $-\ensuremath{\left. \mathsf{path} \right.}$ from start to goal that does not result in a collision



Problem we want to solve

<u>Given</u>:

- a point-robot (robot is a point in space)
- a start and goal configuration

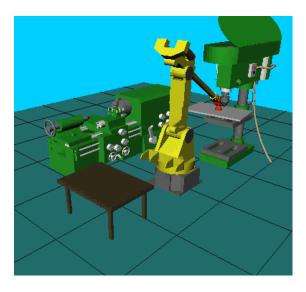
Find:

– path from start to goal that does not result in a collision

Assumptions:

- the position of the robot can always be measured perfectly
- the motion of the robot can always be controlled perfectly
- the robot can move in any directly instantaneously

<u>For example:</u> think about a robot workcell in a factory...



PRMs are specifically designed for high-dimensional configuration spaces – such as the c-space of a robot arm

<u>Problem:</u> robot arm configuration spaces are typically high dimensional

– for example, imagine using the wavefront planner to solve a problem w/ a 10-joint arm

– several variants of the path planning problem have been proven to be PSPACE-hard.

PRMs are specifically designed for high-dimensional configuration spaces – such as the c-space of a robot arm

General idea:

 create a randomized algorithm that will find a solution quickly in many cases

– eventually, the algorithm will be guaranteed to find a solution if one exists *with probability one*

PRMs are specifically designed for high-dimensional configuration spaces – such as the c-space of a robot arm

General idea:

 create a randomized algorithm that will find a solution quickly in many cases

– but, eventually, the algorithm will be guaranteed to find a solution if one exists *with probability one*

With probability one --> "Almost surely"

 the probably of an event NOT happening approaches zero as the algorithm continues to run

Example: an infinite sequence of coin flips contains at least one tail almost surely.

Infinite monkey theorem:

PRN

- SU

<u>Gen</u>

- Cr

case

– bu

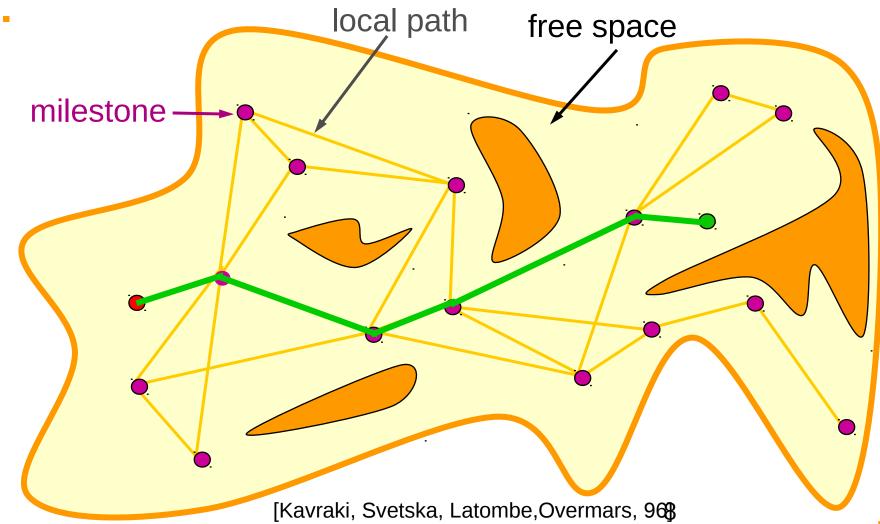
exis

A monkey typing keys randomly on a keyboard will produce any given text (the works of William Shakespeare) *almost surely*



Example: an infinite sequence of coin flips contains at least one tail almost surely. 7

Probabilistic Roadmap (PRM): multiple queries

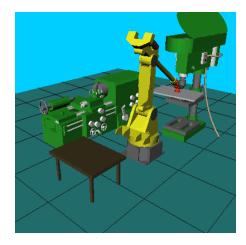


Classic multiple-query PRM

• Probabilistic Roadmaps for Path Planning in High-Dimensional Configuration Spaces, L. Kavraki et al., 1996.

Assumptions

- Static obstacles
- Many queries to be processed in the same environment
- Examples
 - Navigation in static virtual environments
 - Robot manipulator arm in a workcell



Overview

- Precomputation: roadmap construction
 - Uniform sampling
 - Resampling (expansion)
- Query processing

Uniform sampling

Input: geometry of the robot & obstacles
Output: roadmap G = (V, E)

1:
$$V \leftarrow \emptyset$$
 and $E \leftarrow \emptyset$.

2: repeat

- 3: $q \leftarrow a$ configuration sampled uniformly at random from C.
- 4: **if CLEAR**(q)**then**
- 5: Add q to V.
- 6: $N_{\alpha} \leftarrow a$ set of nodes in V that are close to q.
- 6: for each $q' \in N_q$, in order of increasing d(q,q')
- 7: **if** LINK(q',q)**then**
- 8: Add an edge between q and q' to E.

Some terminology

- The graph G is called a **probabilistic roadmap**.
- The nodes in G are called **milestones**.

Query processing

- Connect q_{init} and q_{goal} to the roadmap
- Start at q_{init} and q_{goal} , and try to connect with one of the milestones nearby
- Try multiple times

Error

- If a path is returned, the answer is always correct.
- If no path is found, the answer may or may not be correct. We hope it is correct with high probability.

Probabilistic completeness of PRM

Theorem (Kavraki et al 1998):

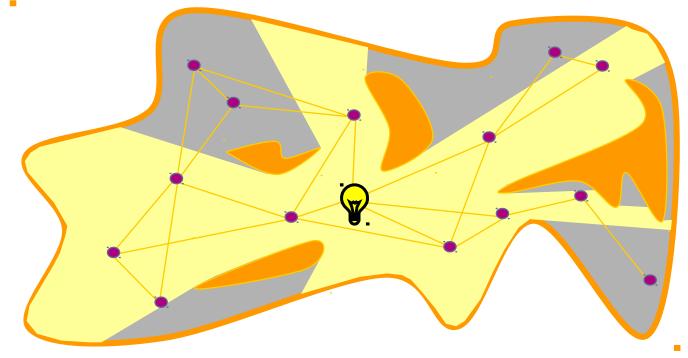
If a path planning problem is feasible, then there exist constants n_0 and a>0, such that:

$$P(\text{a path is found}) \ge 1 - e^{-an}$$

where *n*>n_0 is the number of samples

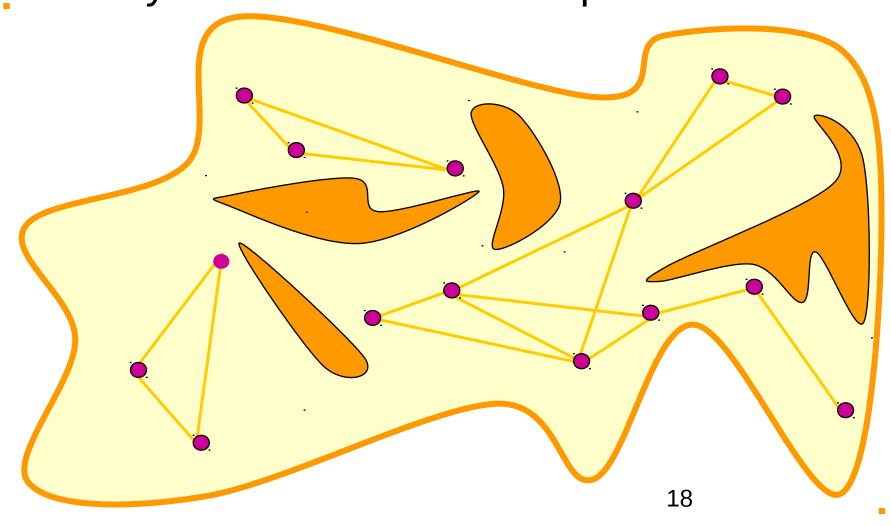
Why does it work? Intuition

• A small number of milestones **almost** "cover" the **entire** configuration space.



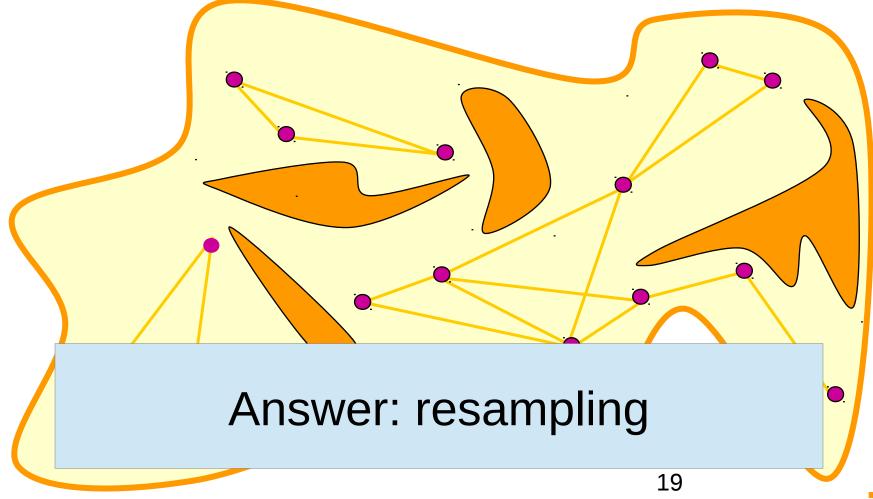
Difficulty

• Many small connected components



Difficulty

Many small connected components



Resampling (expansion)

• Failure rate

$$(q) = \frac{\text{num failed links}}{\text{num link attempts}}$$

• Weight

$$w(q) = \frac{r(q)}{\sum_{p} r(p)}$$

r

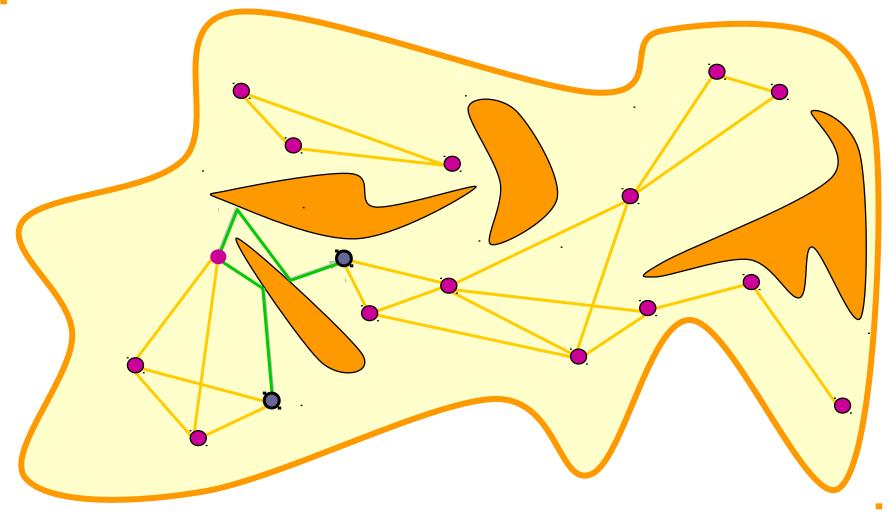
• Resampling probability Pr(q) = w(q)

Resampling (expansion)

Once a node is selected to be expanded:

- 1. Pick a random motion direction in c-space and move in this direction until an obstacle is hit.
- 2. When a collision occurs, choose a new random direction and proceed for some distance.
- 3. Add the resulting nodes and edges to the tree. Re-run tree connection step.

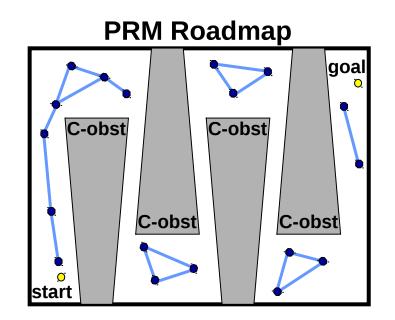
Resampling (expansion)



Gaussian sampler

So far, we have only discussed uniform sampling...

<u>Problem</u>: uniform sampling is not a great way to find paths through narrow passageways.



Gaussian sampler

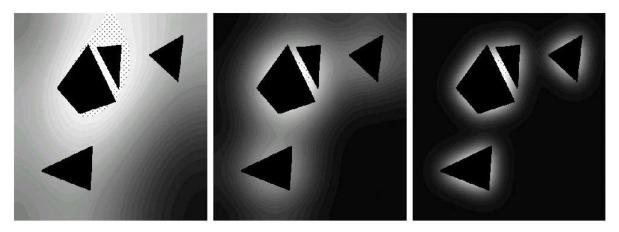
Gaussian sampler:

Sample points uniformly at random (as before)
For each sampled point, sample a second point
from a Gaussian distribution centered at the first
sampled point

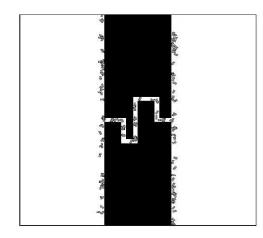
Discard the first sample if both samples are either free or in collision

– Keep the fist sample if the two samples are NOT both free or both in collision (that is, keep the sample if the free/collision status of the second sample is different from the first).

Gaussian sampler



Probability of sampling a point under the Gaussian sampler as a function of distance from a c-space obstacle



Example of samples drawn from Gaussian sampler

Single-Query PRM

NUS CS 5247 David Hsu

Lazy PRM

• *Path Planning Using Lazy PRM*, R. Bohlin & L. Kavraki, 2000.

Precomputation: roadmap construction

- Nodes
 - Randomly chosen configurations, which may or may **not** be collision-free
 - No call to clear
- Edges
 - an edge between two nodes if the corresponding configurations are close according to a suitable metric
 - no call to link

Query processing: overview

- 1. Find a shortest path in the roadmap
- 2. Check whether the nodes and edges in the path are collision.
- 3. If yes, then done. Otherwise, remove the nodes or edges in violation. Go to (1).

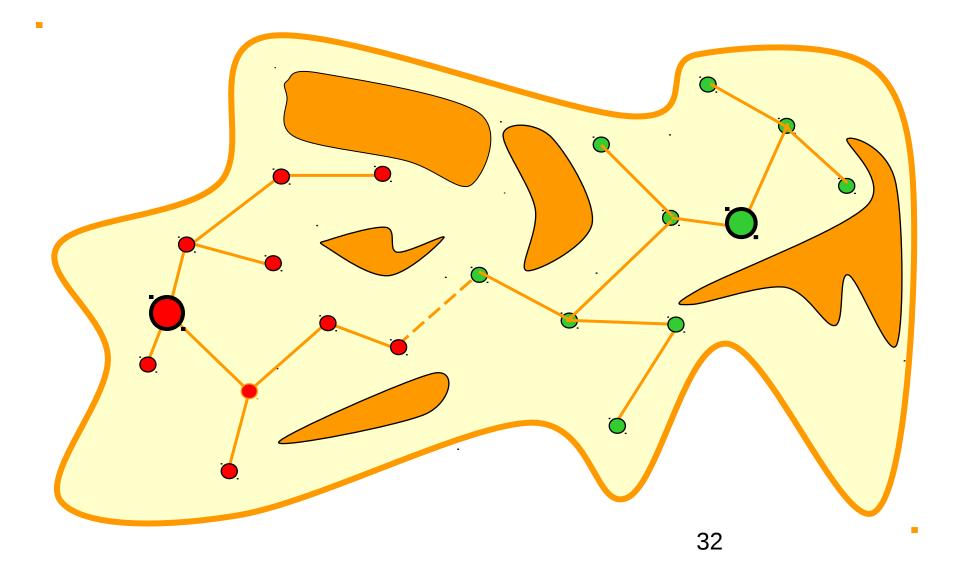
We either find a collision-free path, or exhaust all paths in the roadmap and declare failure.

Query processing: details

- Find the shortest path in the roadmap
 - A* algorithm
 - Dijkstra's algorithm (uniform cost search)
- Check whether nodes and edges are collisions free
 - -**CLEAR**(q)
 - -LINK (q_0, q_1)

Supplemental

Probabilistic Roadmap (PRM): single query



Multiple-Query PRM

NUS CS 5247 David Hsu