Breadth first search
Uniform cost search

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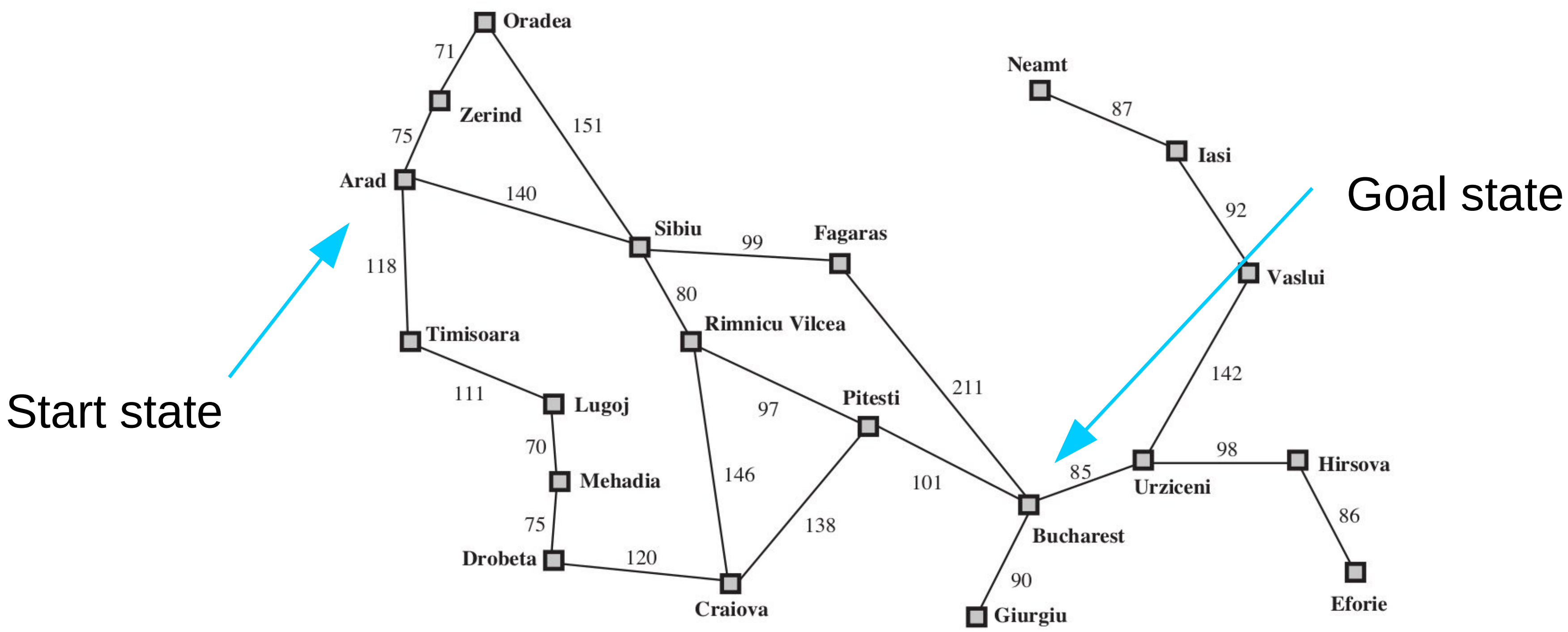
Northeastern University

Some images and slides are used from:
1. CS188 UC Berkeley
2. RN, AIMA

What is graph search?

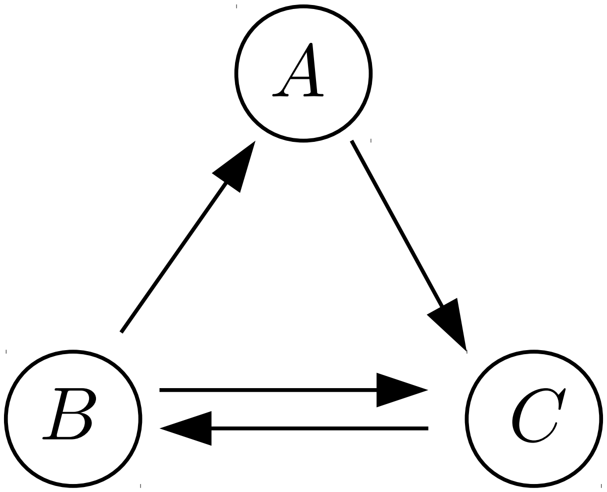


What is a graph?

Graph: $G = (V, E)$

Vertices: V

Edges: E



Directed graph

$$V = \{A, B, C\}$$
$$E = \{(B, A), (A, C), (B, C), (C, B)\}$$

What is a graph?

Graph: $G = (V, E)$

Vertices: V

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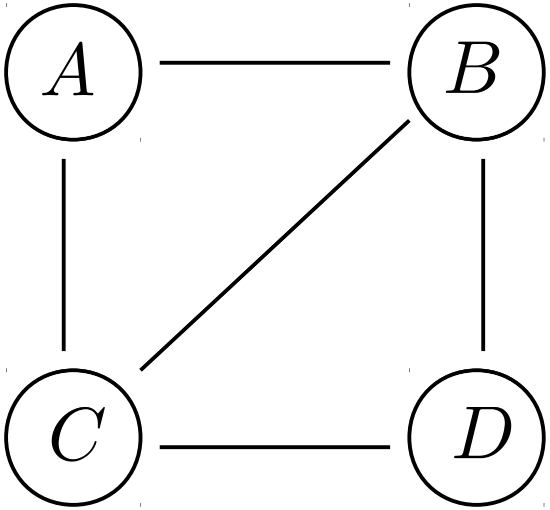
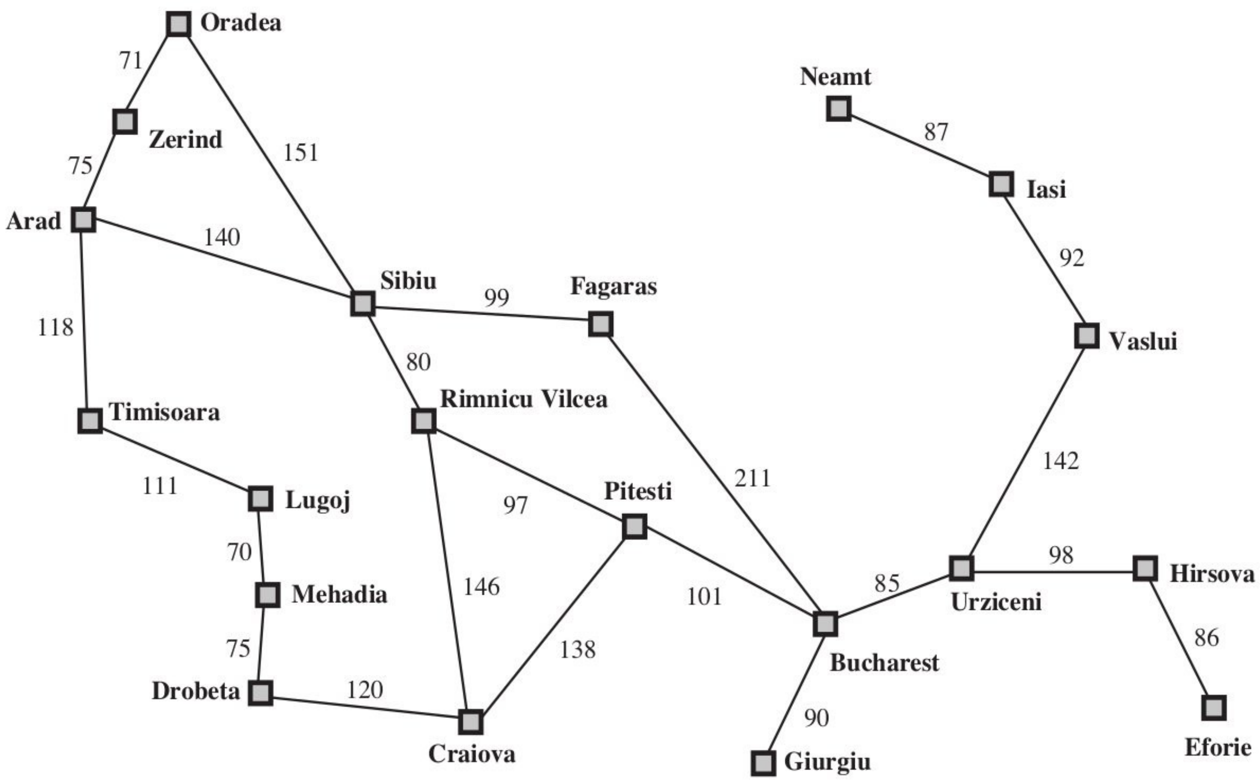


Undirected graph

$$V = \{A, B, C, D\}$$

$$E = \{\{A, C\}, \{A, B\}, \{C, D\}, \{B, D\}, \{C, B\}\}$$

Graph search



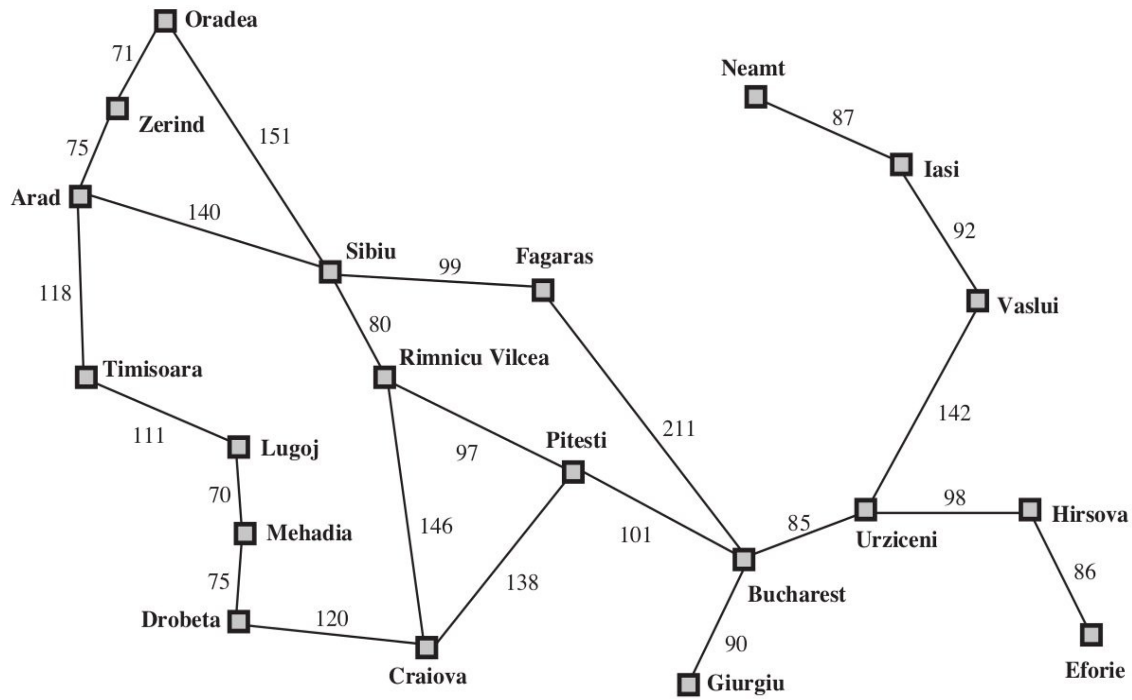
<u>Given</u>: a graph, *G*

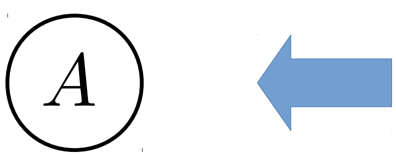
<u>Problem</u>: find a path from A to B

– A: start state

– B: goal state

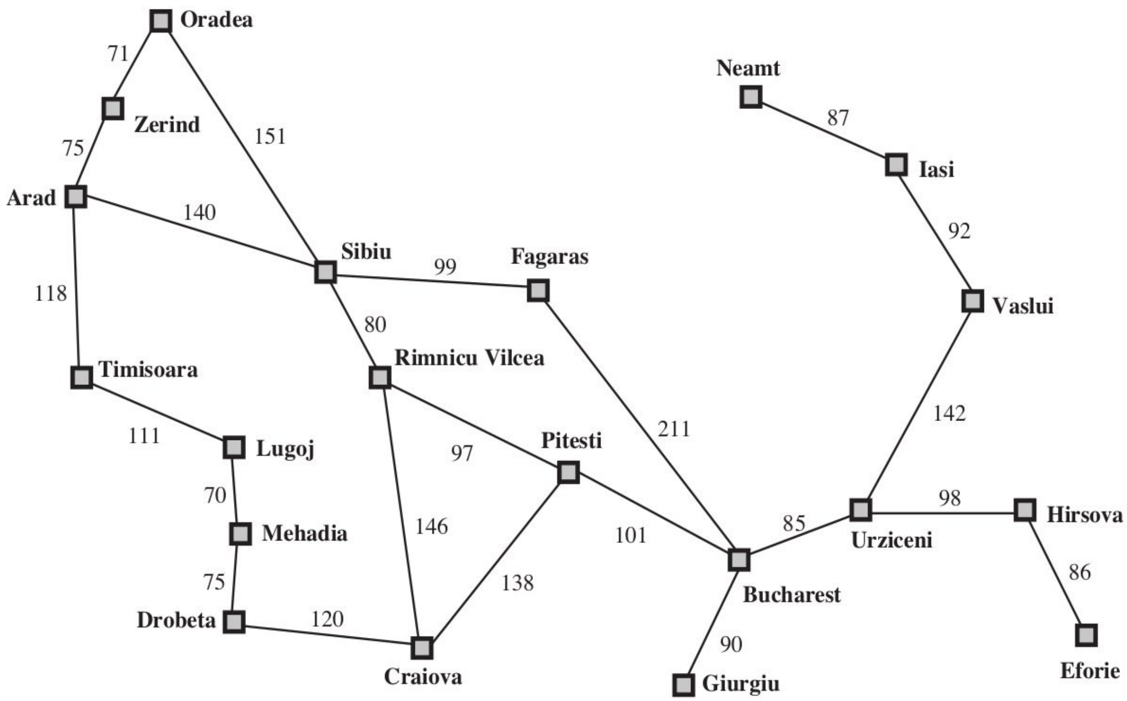
A search tree

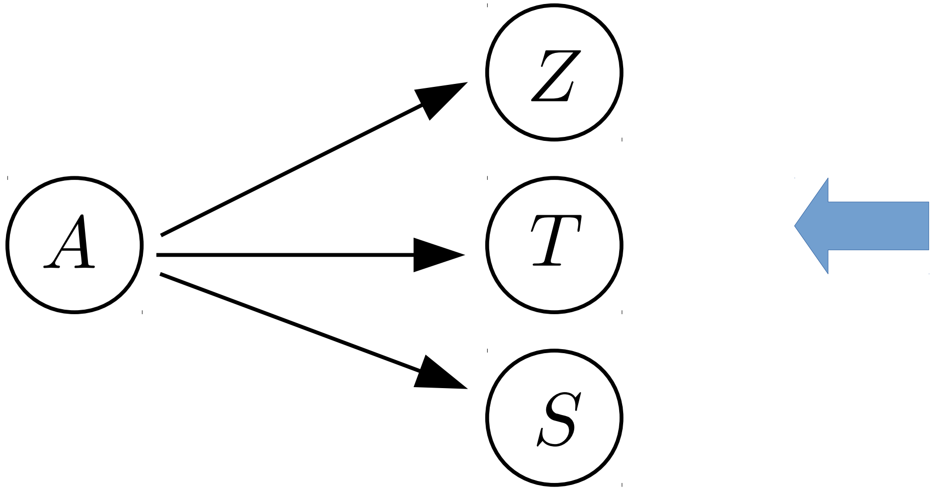




Start at *A*

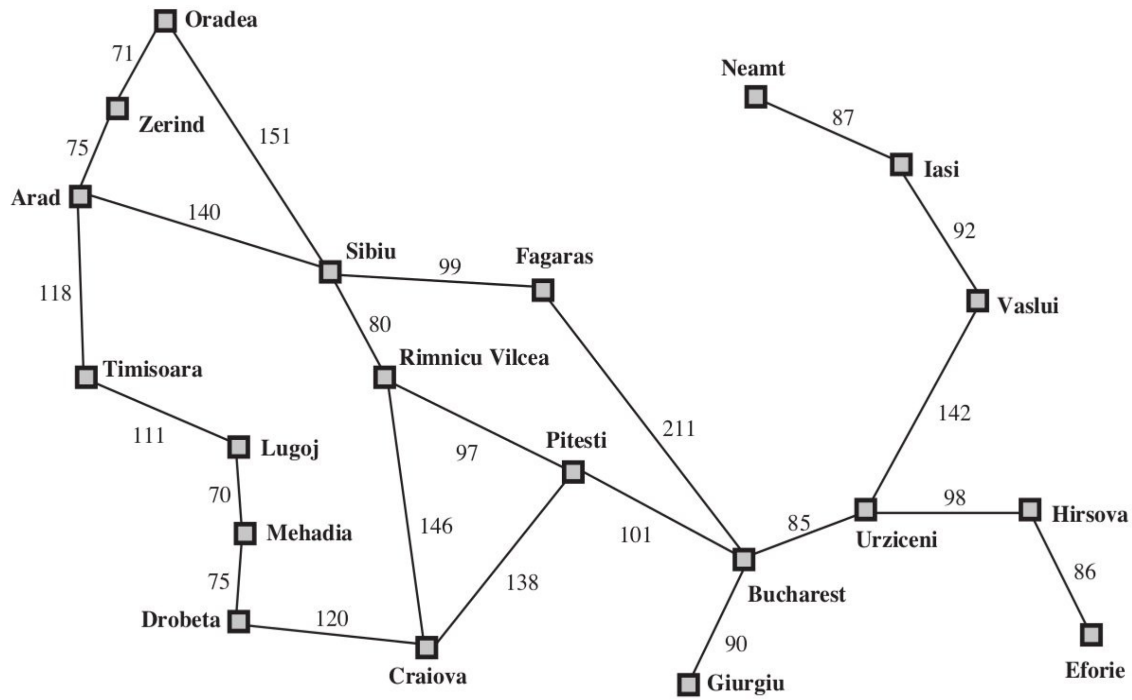
A search tree

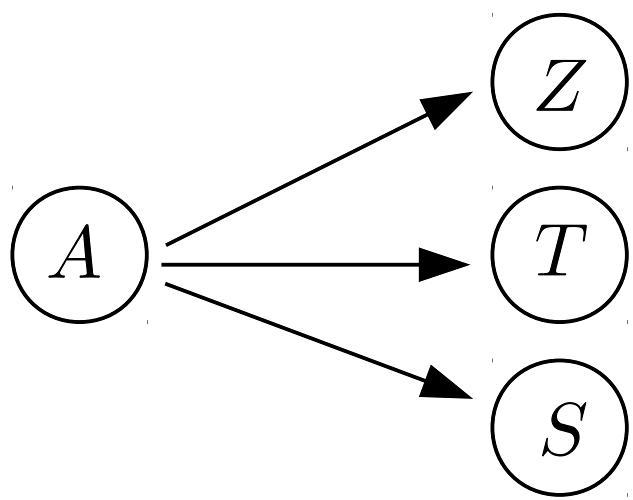




Successors of *A*

A search tree



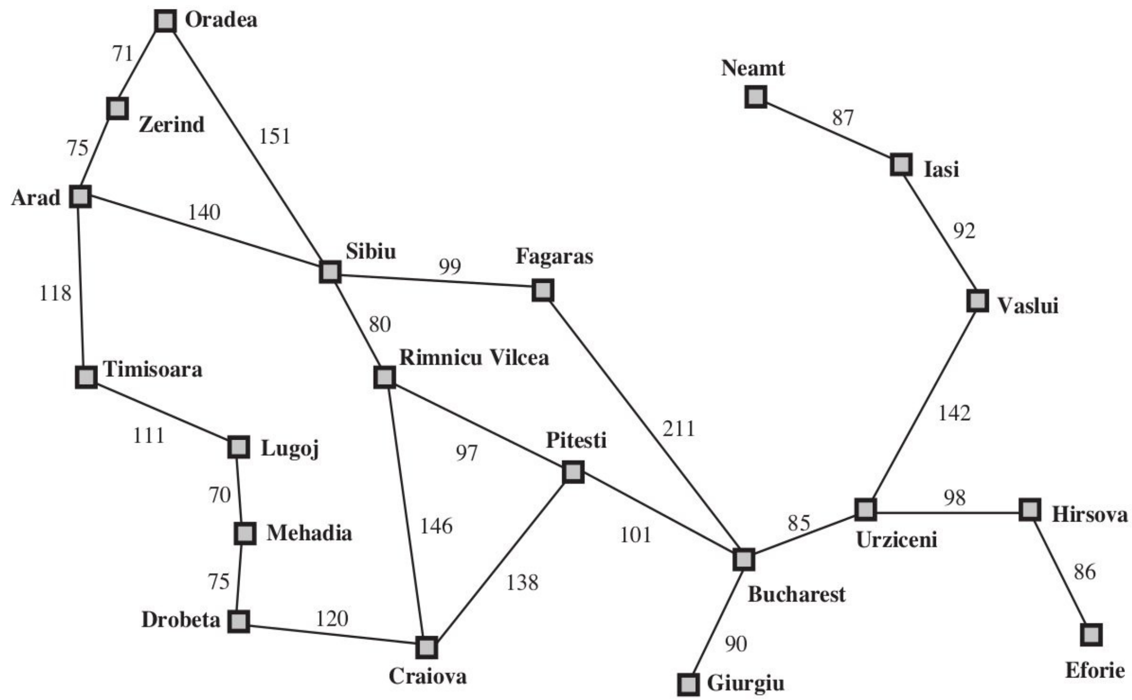


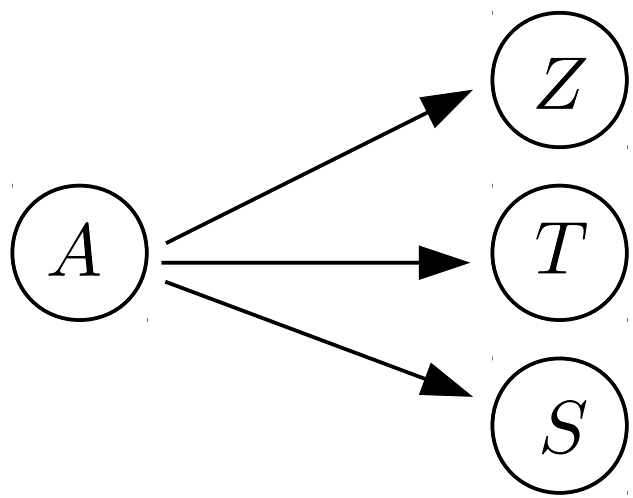
Successors of *A*

parent

children

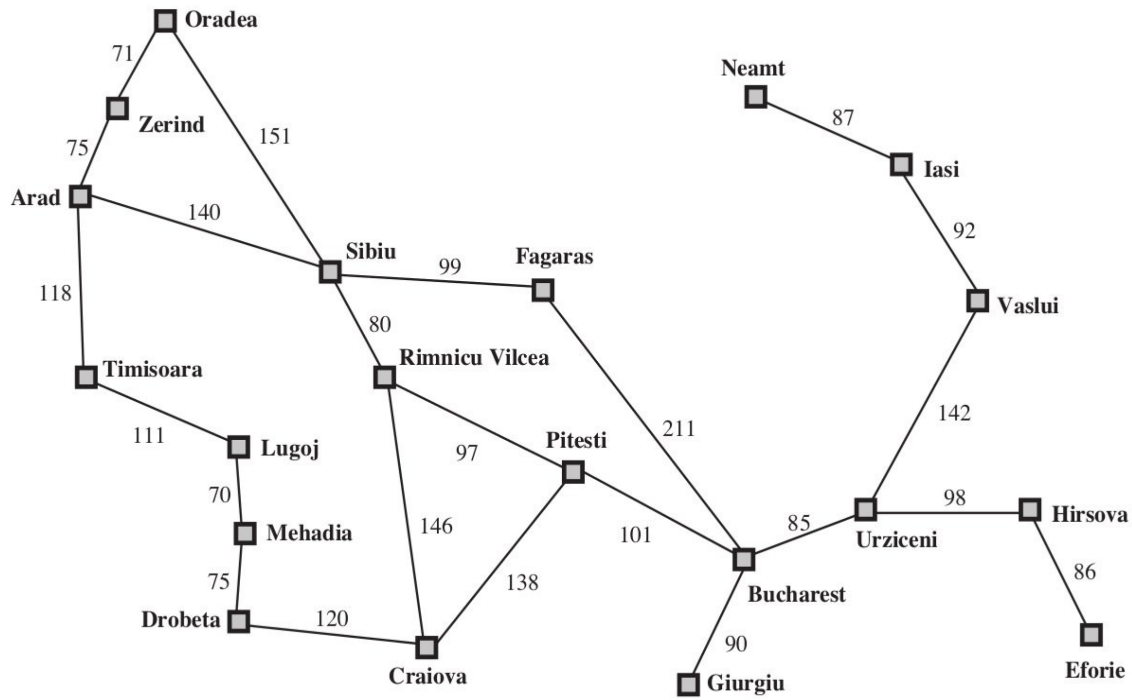
A search tree

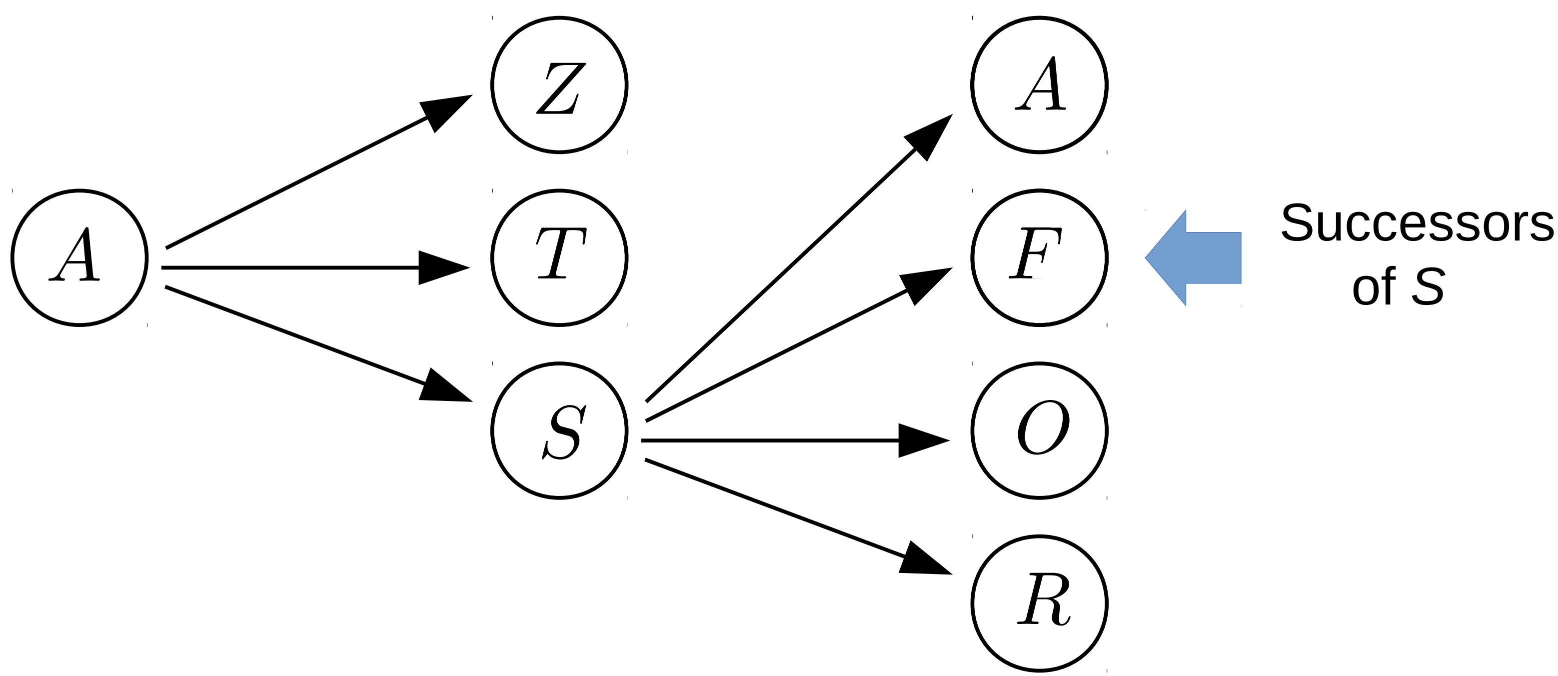




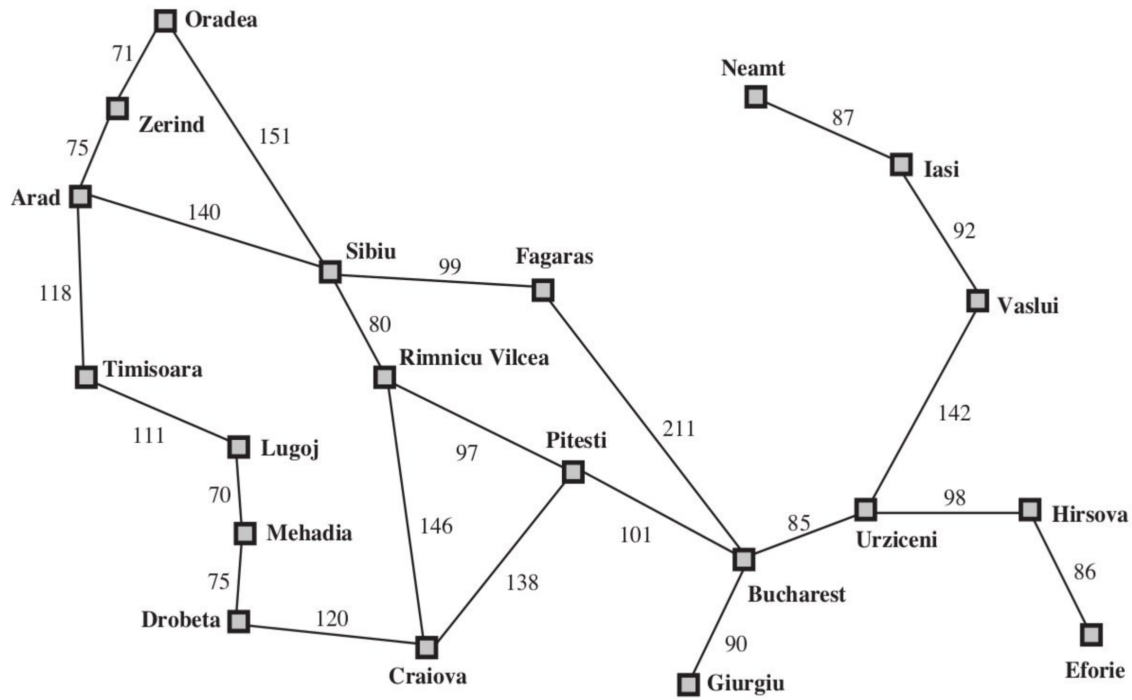
Let's expand *S* next

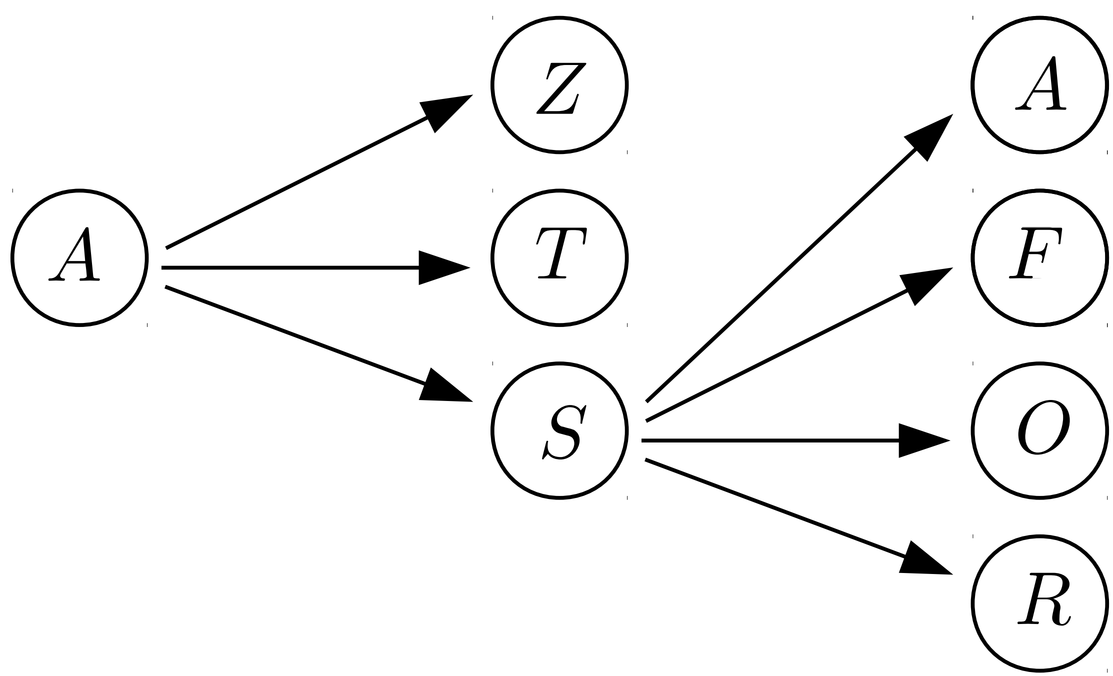
A search tree





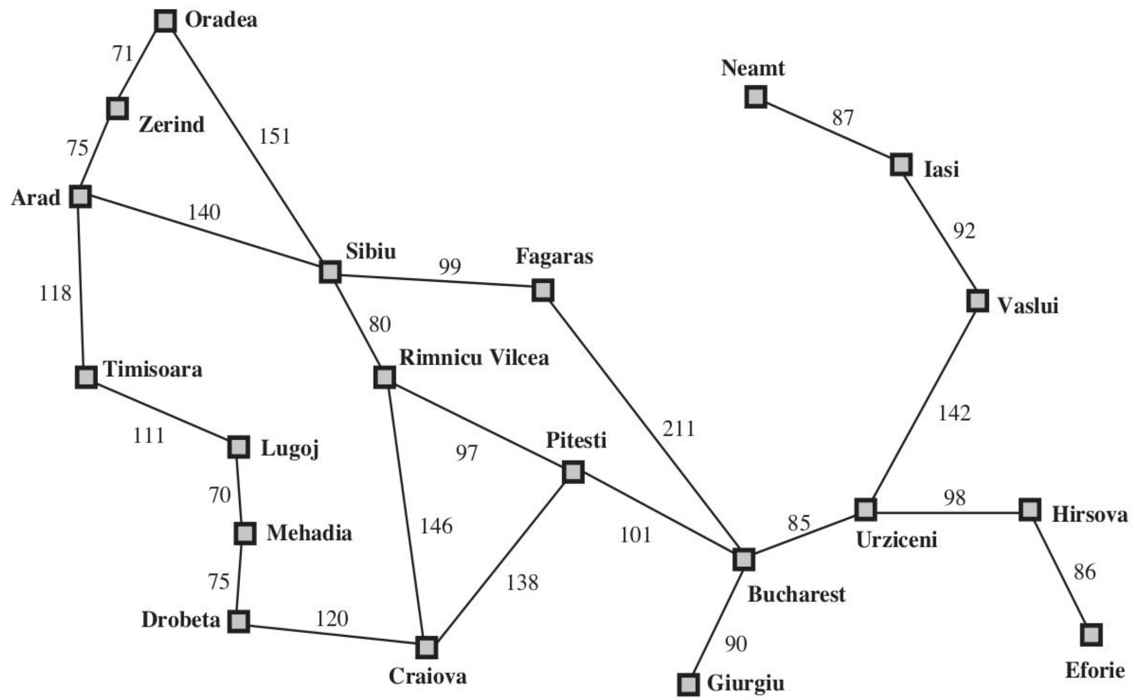
A search tree



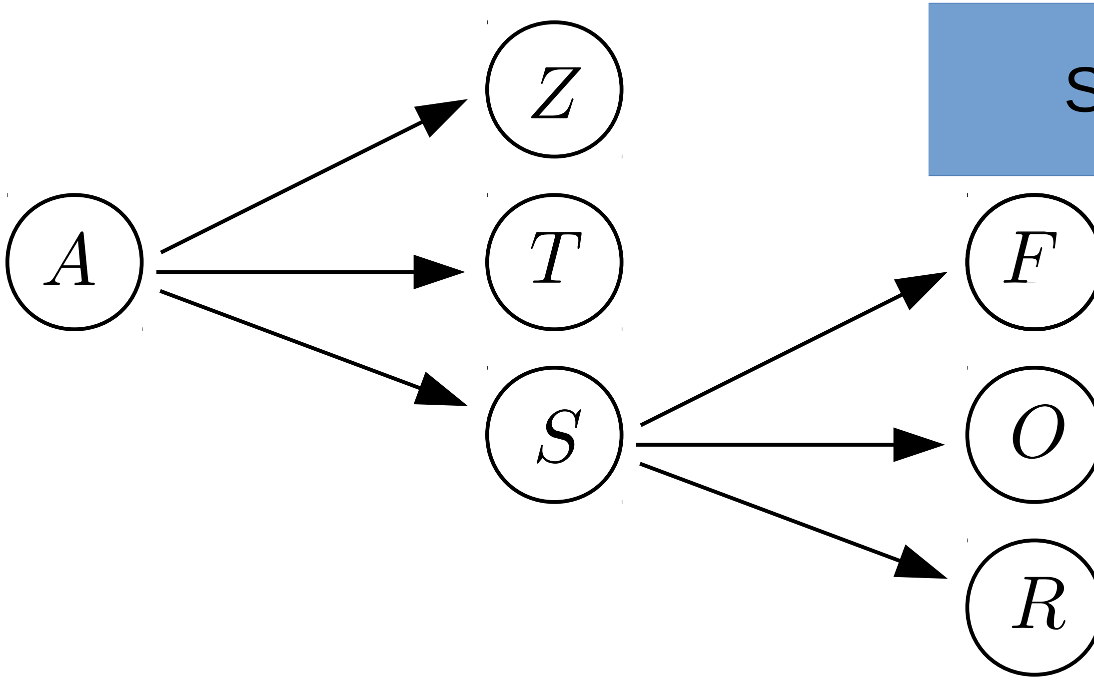


A was already visited!

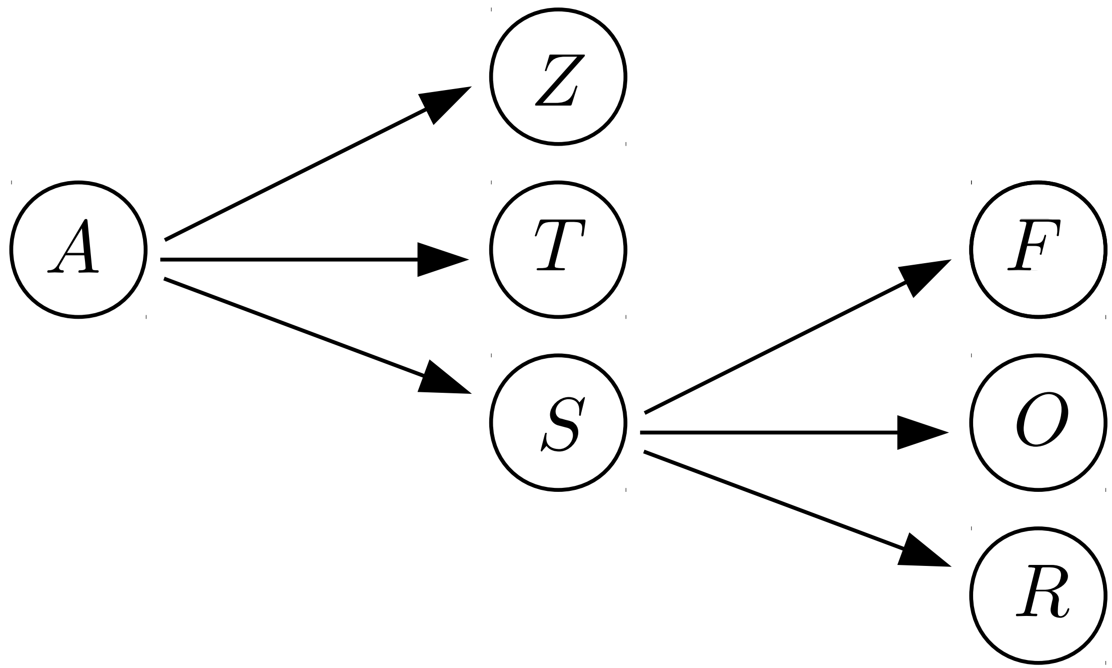
A search tree



So, prune it!



A search tree

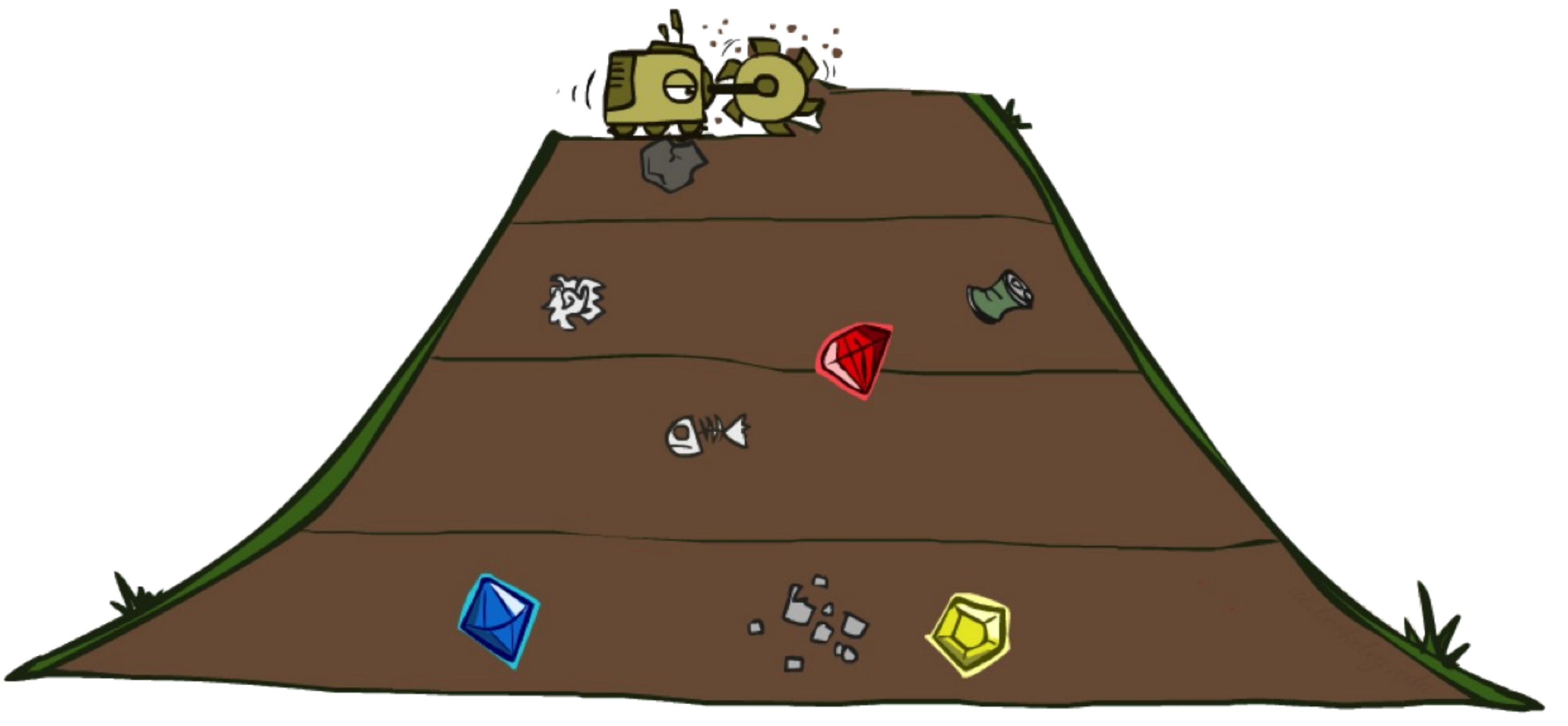


<u>In what order should we expand states?</u>

– here, we expanded *S*, but we could also have expanded *Z* or *T*

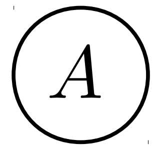
– different search algorithms expand in different orders

Breadth first search (BFS)



Slide: Adapted from Berkeley CS188 course notes (downloaded Summer 2015)

Breadth first search (BFS)



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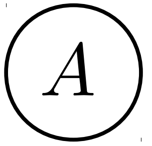
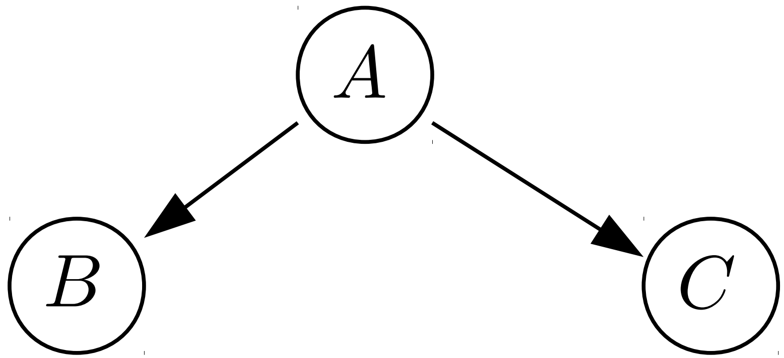


Start node

Breadth first search (BFS)



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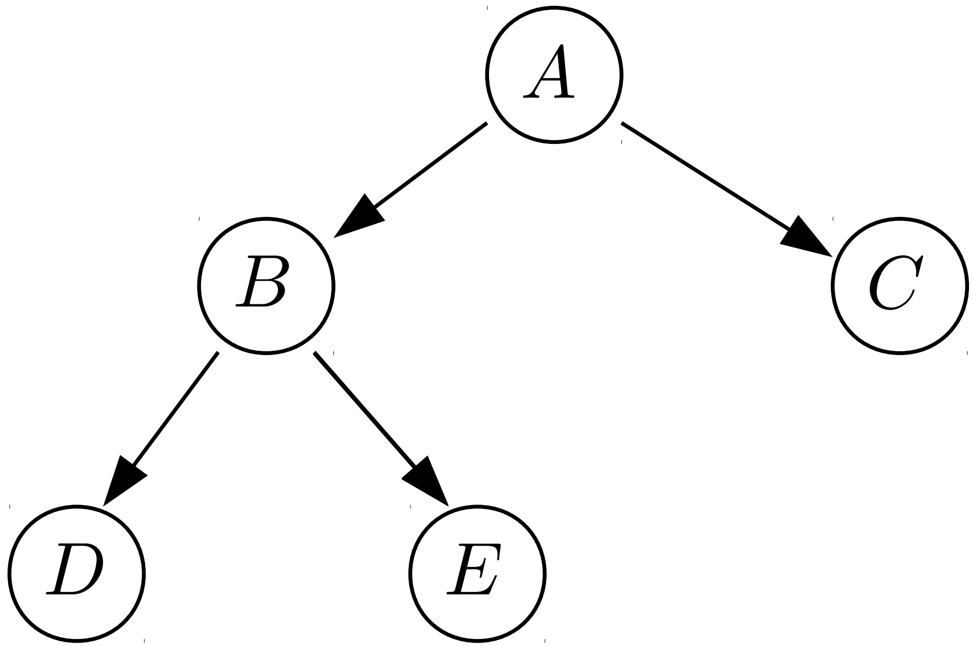

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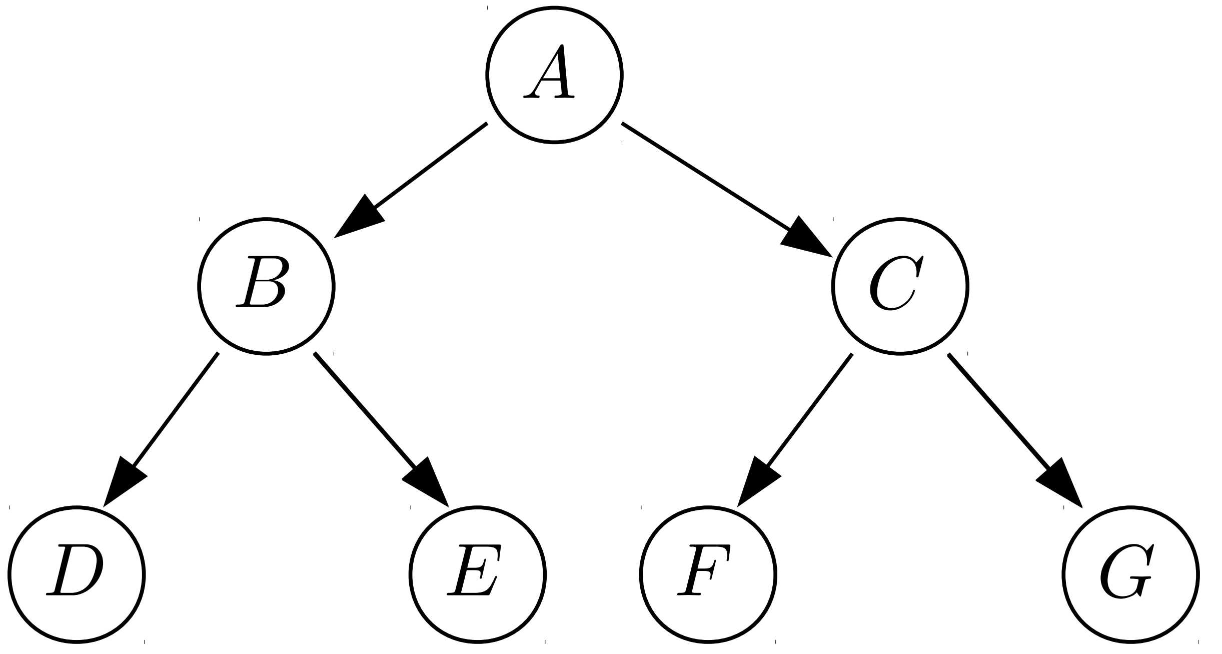


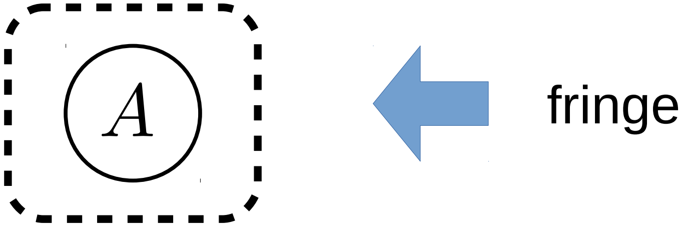
Breadth first search (BFS)

Fringe

We're going to maintain a queue called the fringe

– initialize the fringe as an empty queue

Breadth first search (BFS)

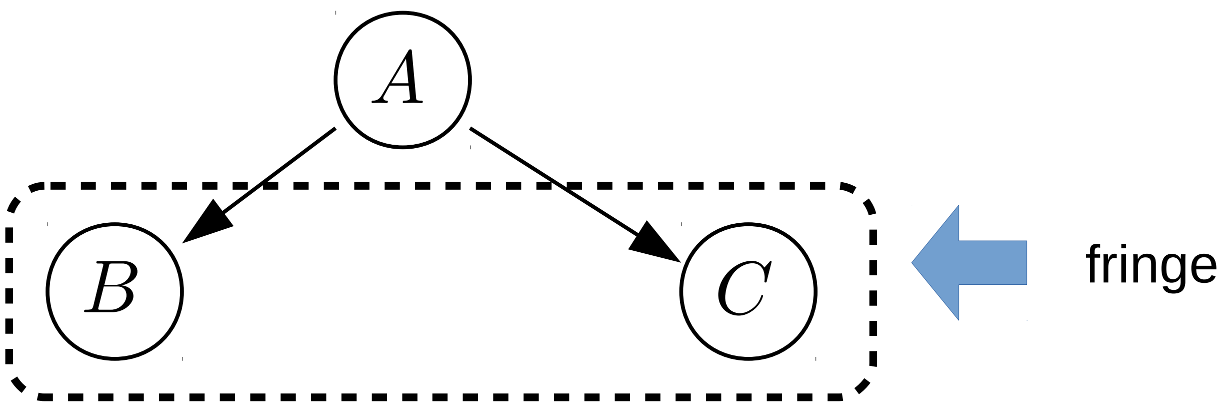


Fringe
A

– add *A* to the fringe

Breadth first search (BFS)

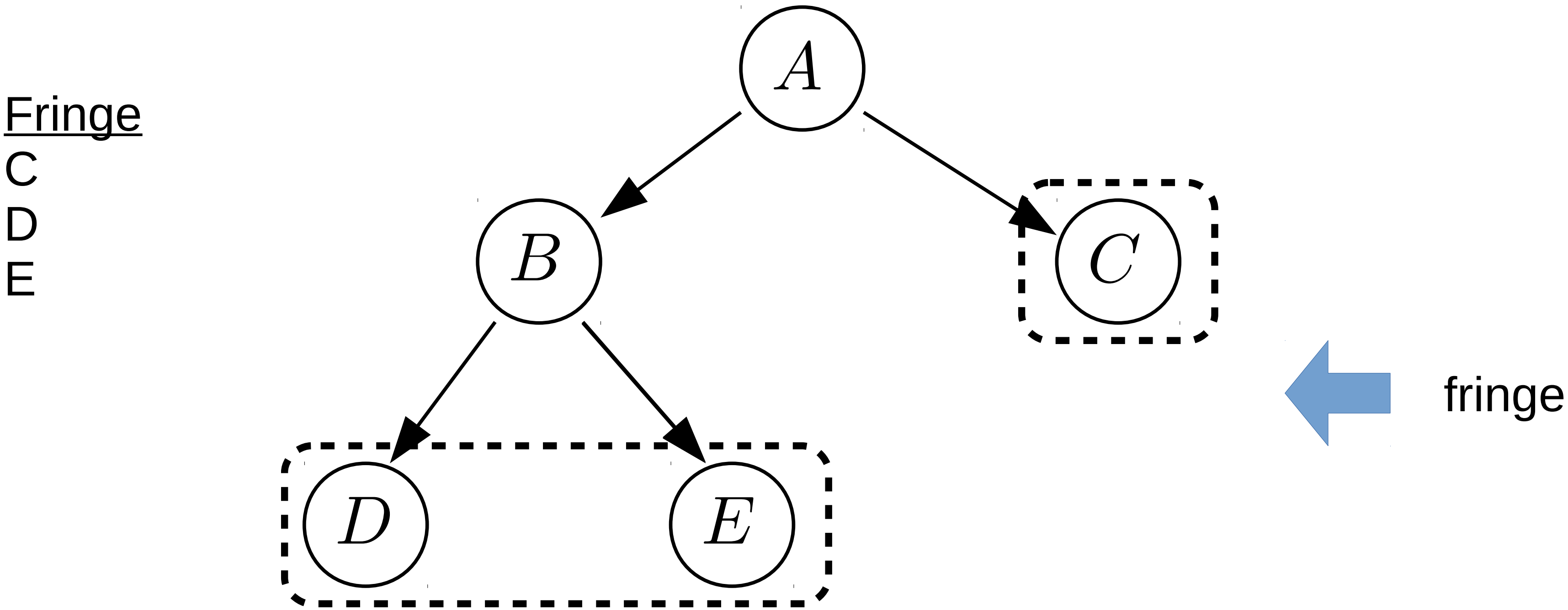
<u>Fringe</u>
B
C



-- remove *A* from the fringe

-- add successors of *A* to the fringe

Breadth first search (BFS)

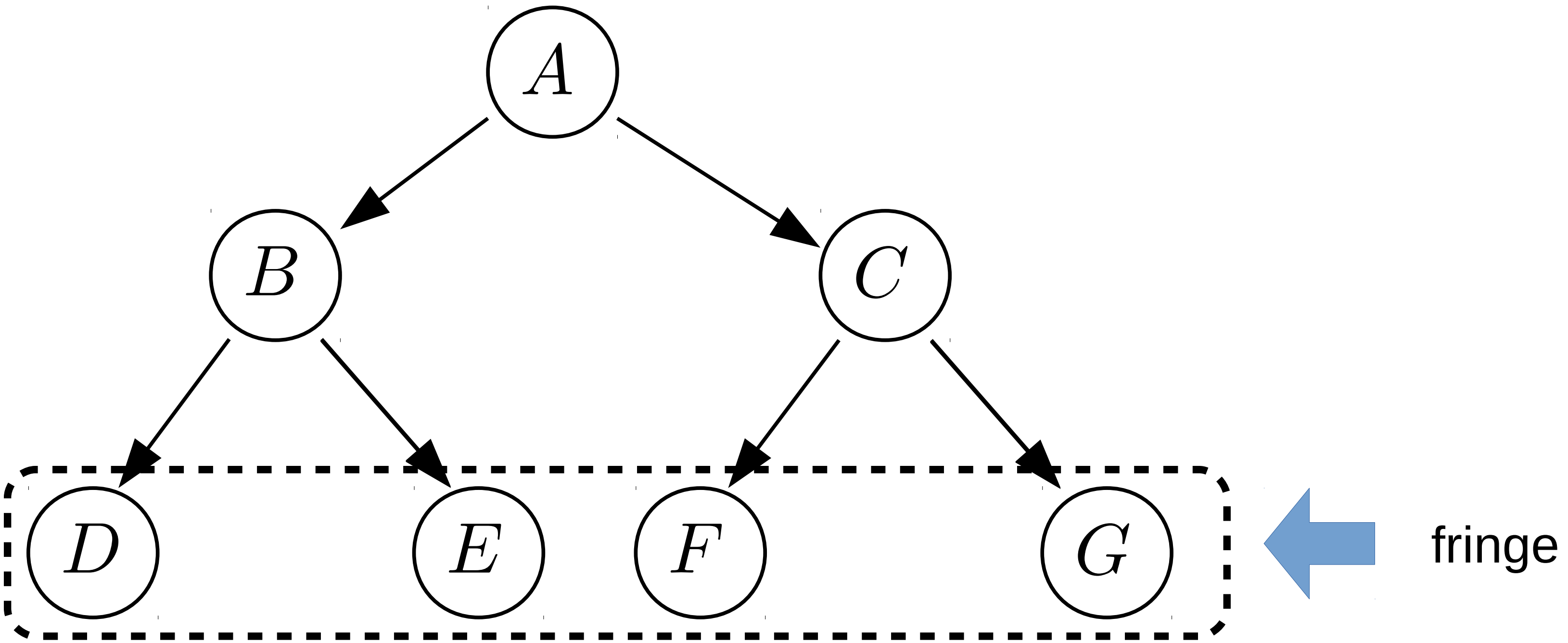


-- remove *B* from the fringe

-- add successors of *B* to the fringe

Breadth first search (BFS)

<u>Fringe</u>
D
E
F
G

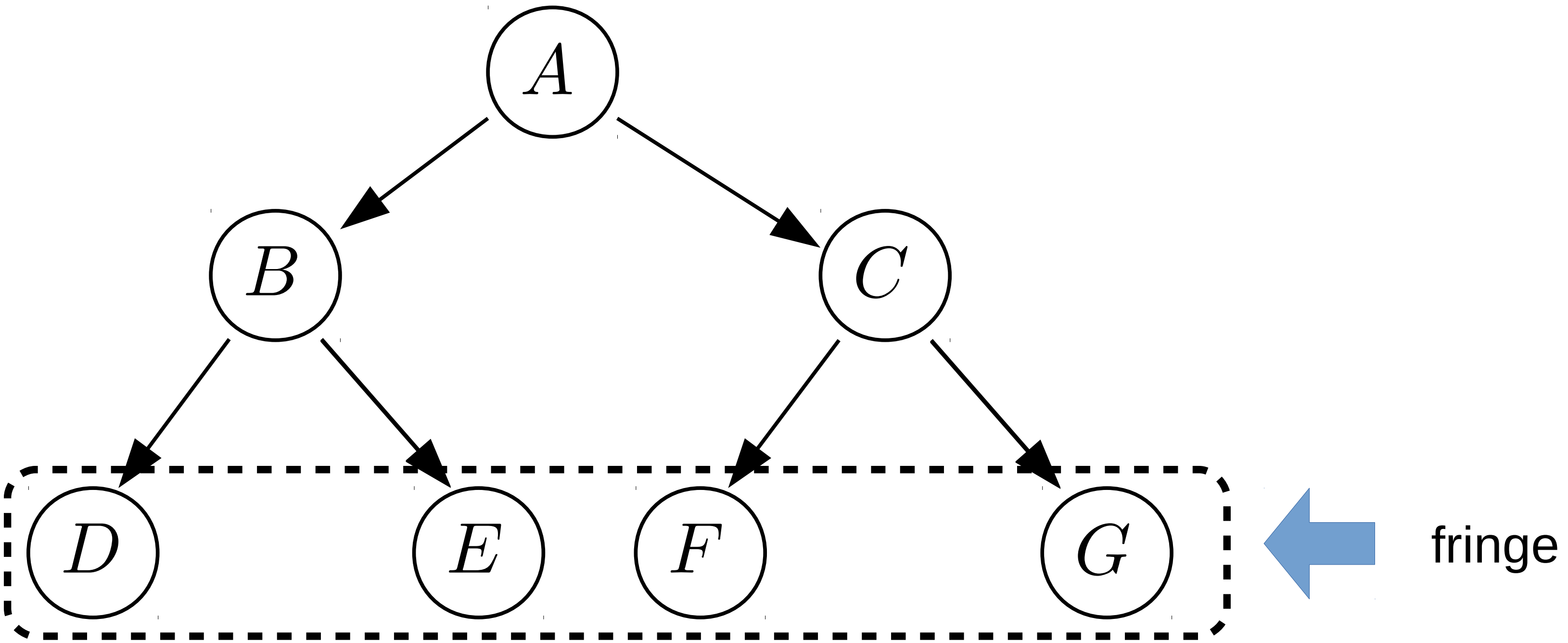


-- remove *C* from the fringe

-- add successors of *C* to the fringe

Breadth first search (BFS)

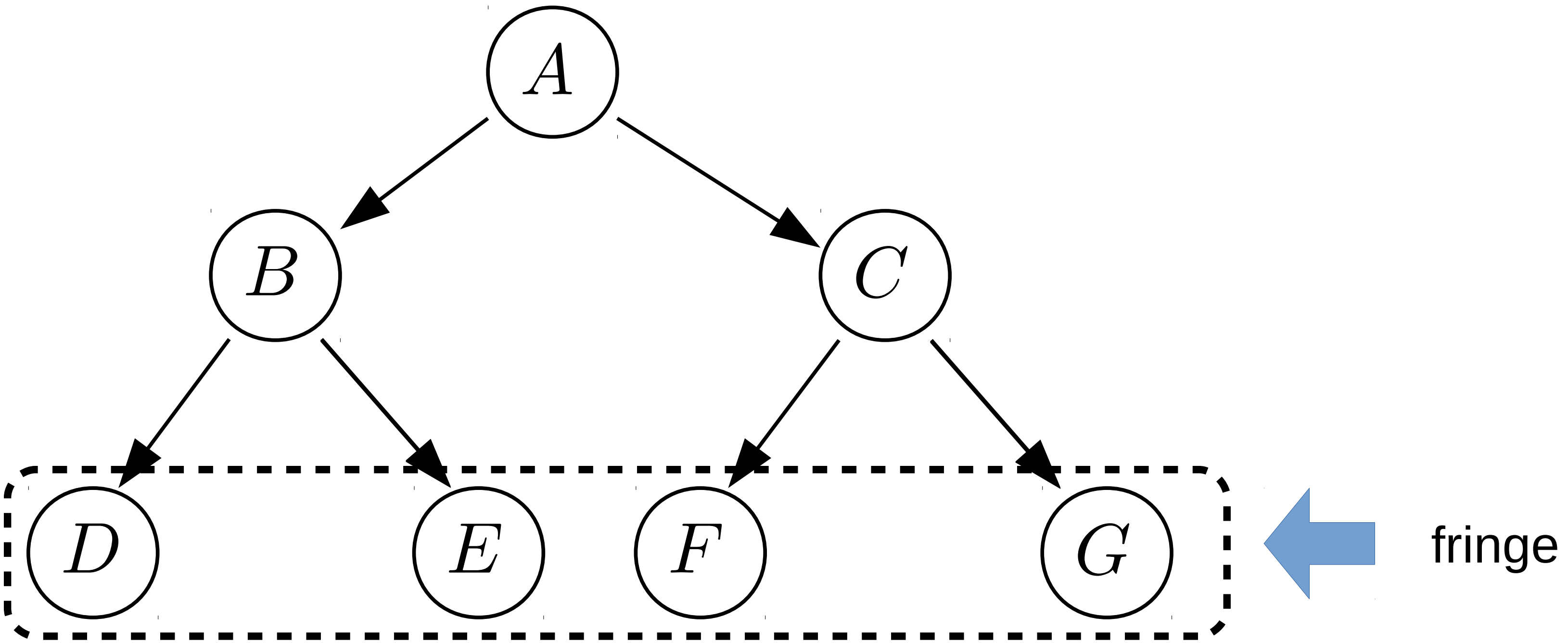
<u>Fringe</u>
D
E
F
G



Which state gets removed next from the fringe?

Breadth first search (BFS)

Fringe
D
E
F
G



Which state gets removed next from the fringe?

What kind of a queue is this?

Breadth first search (BFS)

<u>Fringe</u>
D
E
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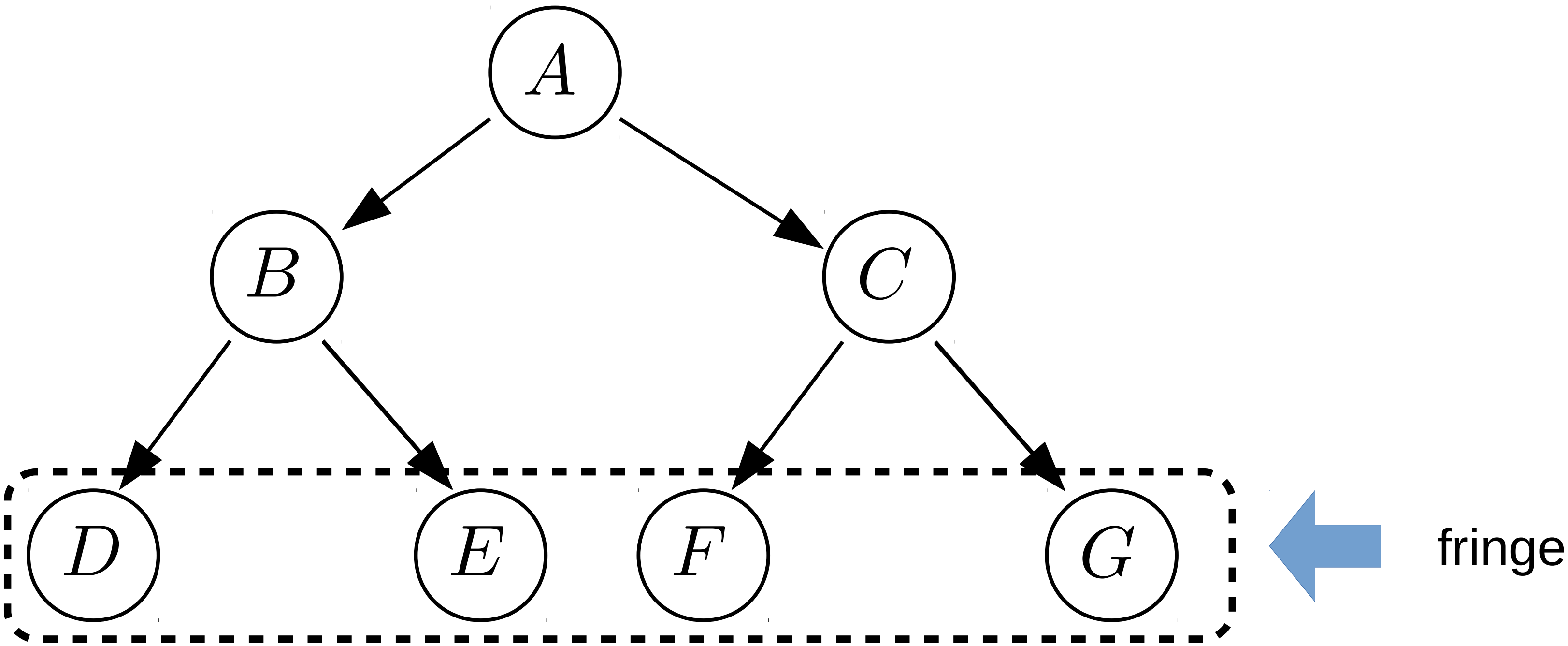
G



Which state gets removed next from the fringe?

What kind of a queue is this?

FIFO Queue!
(first in first out)

Breadth first search (BFS)

```
function BREADTH-FIRST-SEARCH(problem) returns a solution, or failure
  node ← a node with STATE = problem.INITIAL-STATE, PATH-COST = 0
  if problem.GOAL-TEST(node.STATE) then return SOLUTION(node)
  frontier ← a FIFO queue with node as the only element
  explored ← an empty set
  loop do
     if EMPTY?(frontier) then return failure
     node ← POP(frontier)   /* chooses the shallowest node in frontier */
     add node.STATE to explored
     for each action in problem.ACTIONS(node.STATE) do
        child ← CHILD-NODE(problem, node, action)
        if child.STATE is not in explored or frontier then
           if problem.GOAL-TEST(child.STATE) then return SOLUTION(child)
           frontier ← INSERT(child, frontier)
```

Figure 3.11 Breadth-first search on a graph.

Breadth first search (BFS)

```
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```

Figure 3.11 Breadth-first search on a graph.

What is the purpose of the *explored* set?

BFS Properties

Is BFS complete?
– is it guaranteed to find a solution if one exists?

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What is the time complexity of BFS?
– how many states are expanded before finding a sol'n?
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  – d: depth of shallowest solution
  – complexity = ???

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What is the space complexity of BFS?
– how much memory is required?
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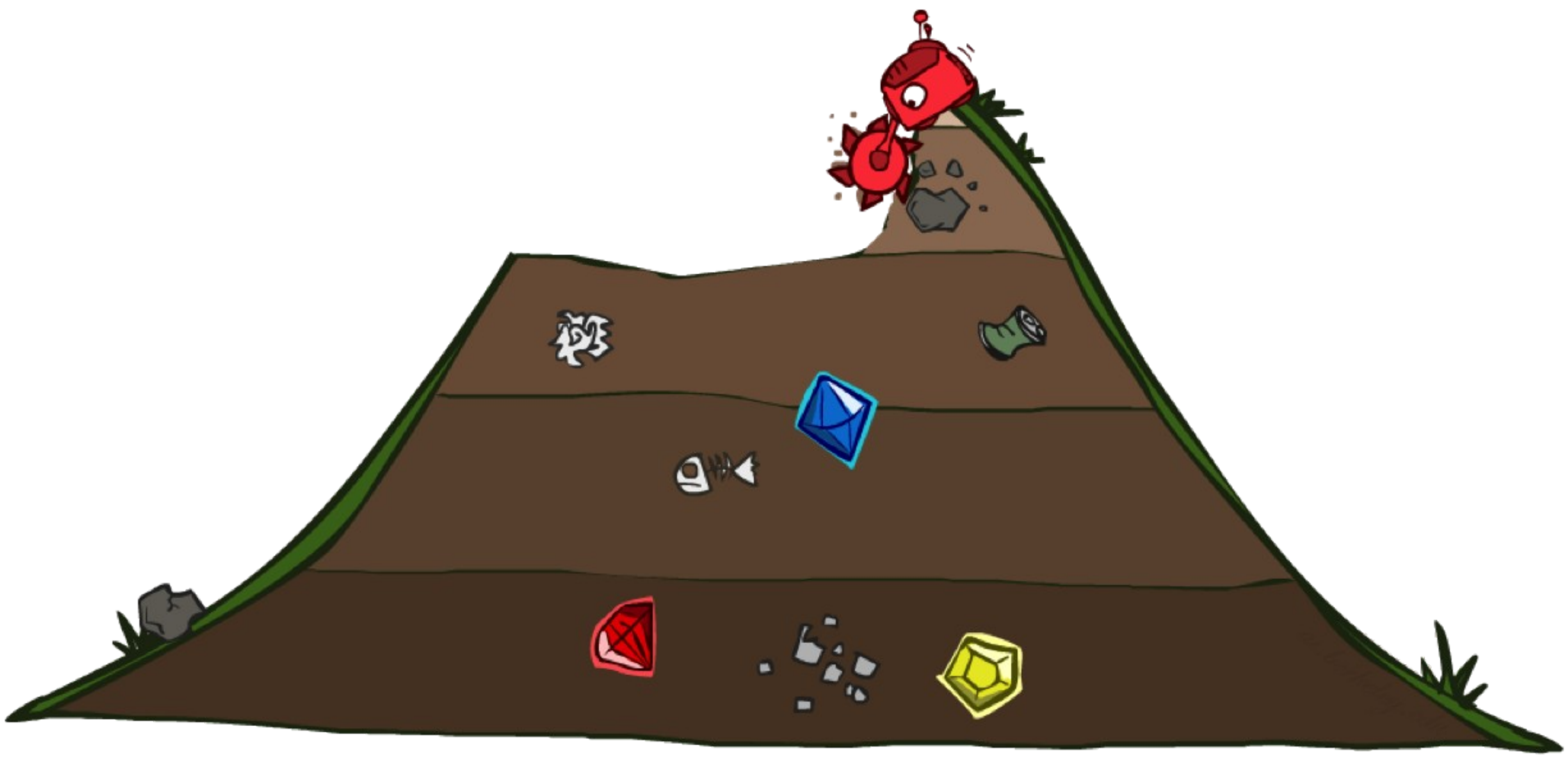
Is BFS optimal?
– is it guaranteed to find the best solution (shortest path)?

Another BFS example...

Uniform Cost Search (UCS)



Slide: Adapted from Berkeley CS188 course notes (downloaded Summer 2015)

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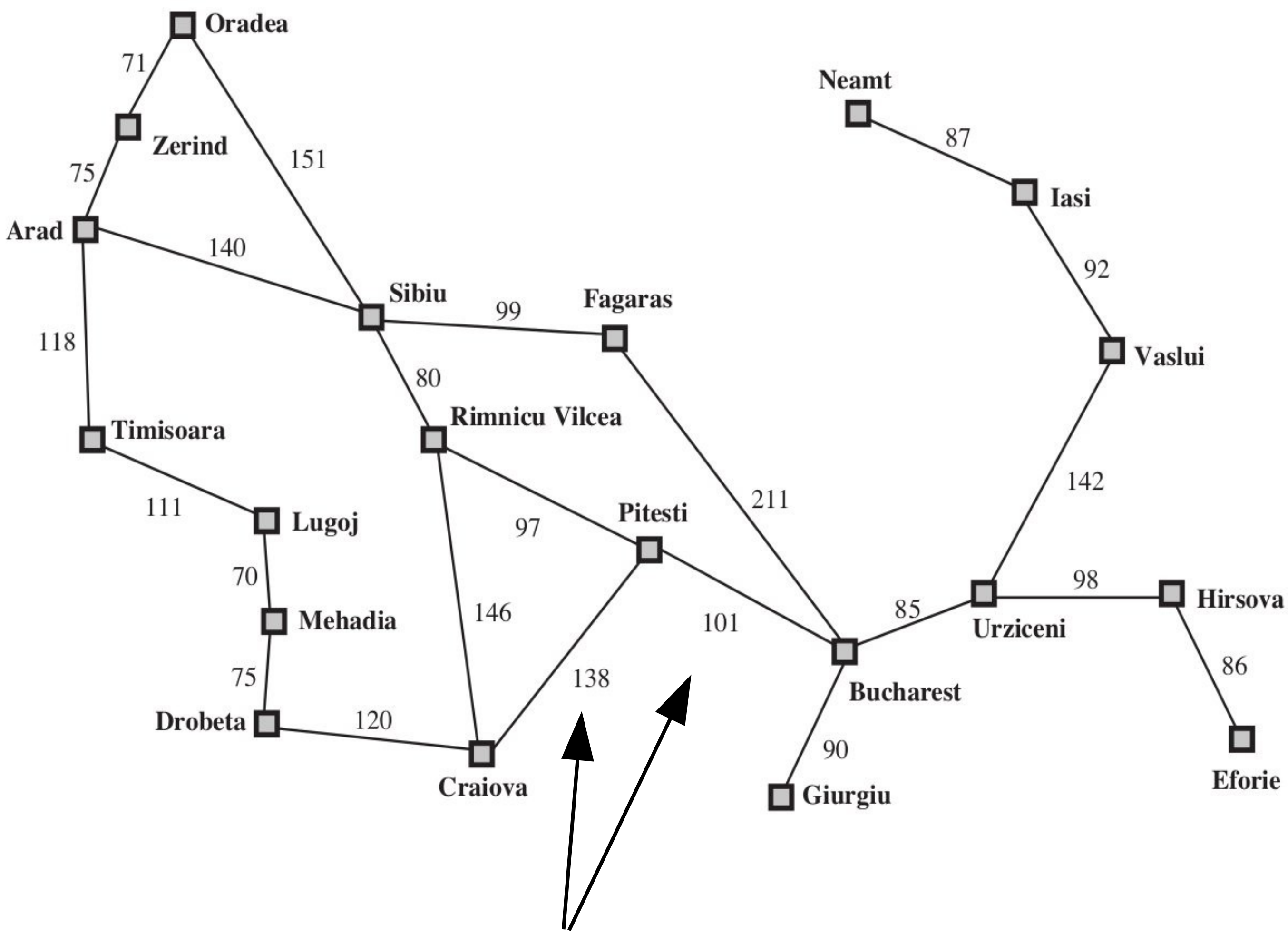
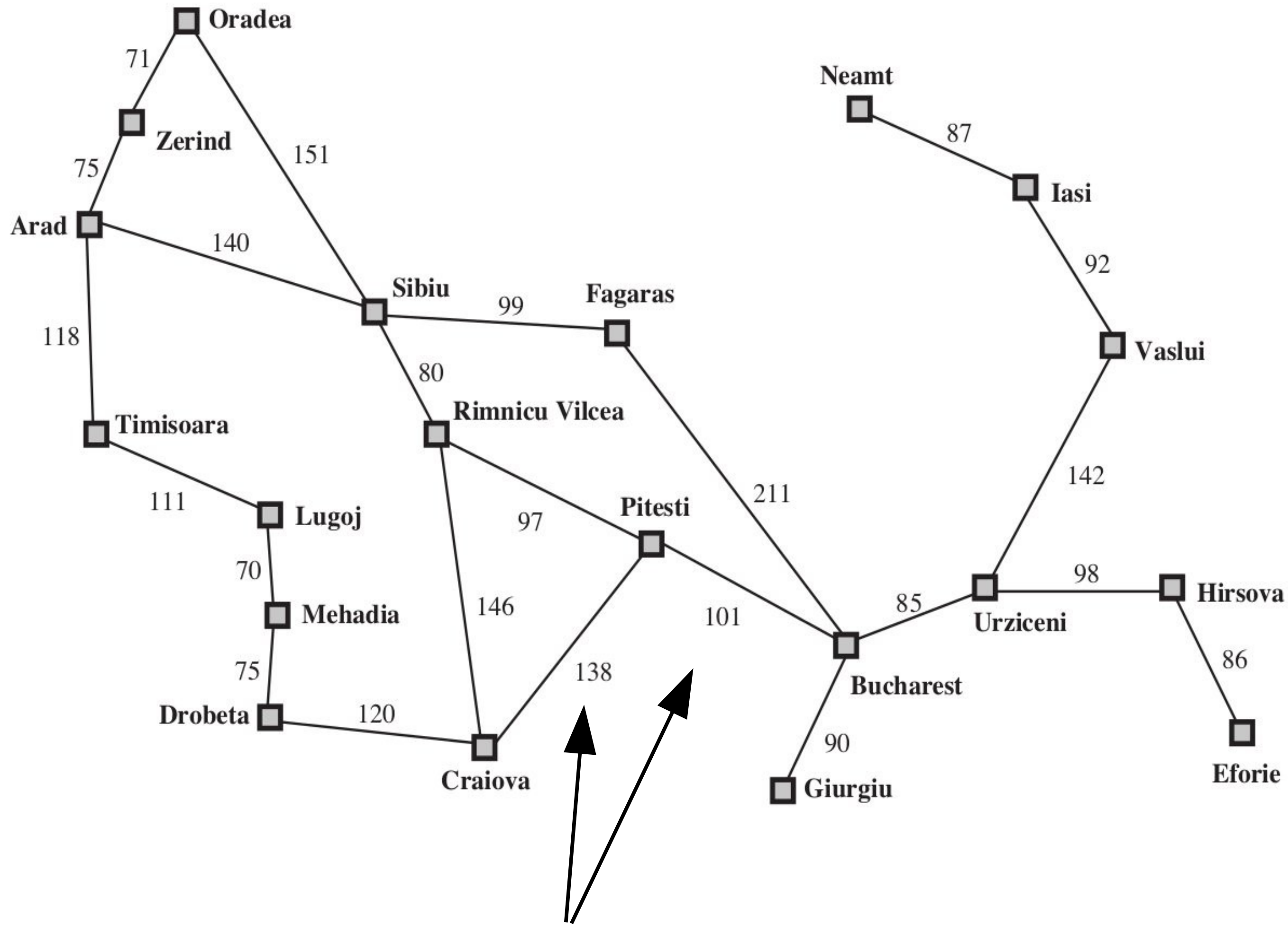


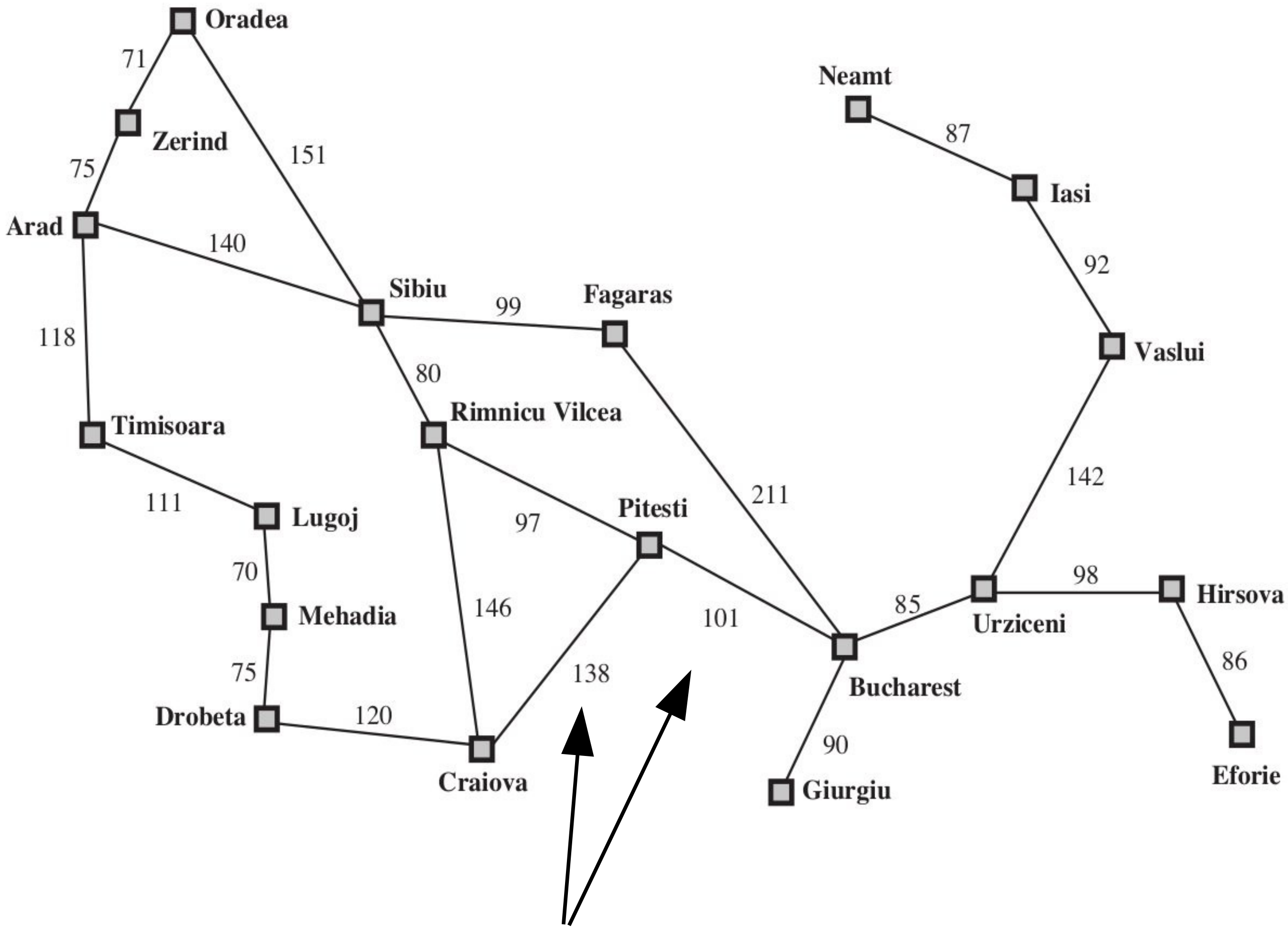
Notice the distances between cities

Uniform Cost Search (UCS)



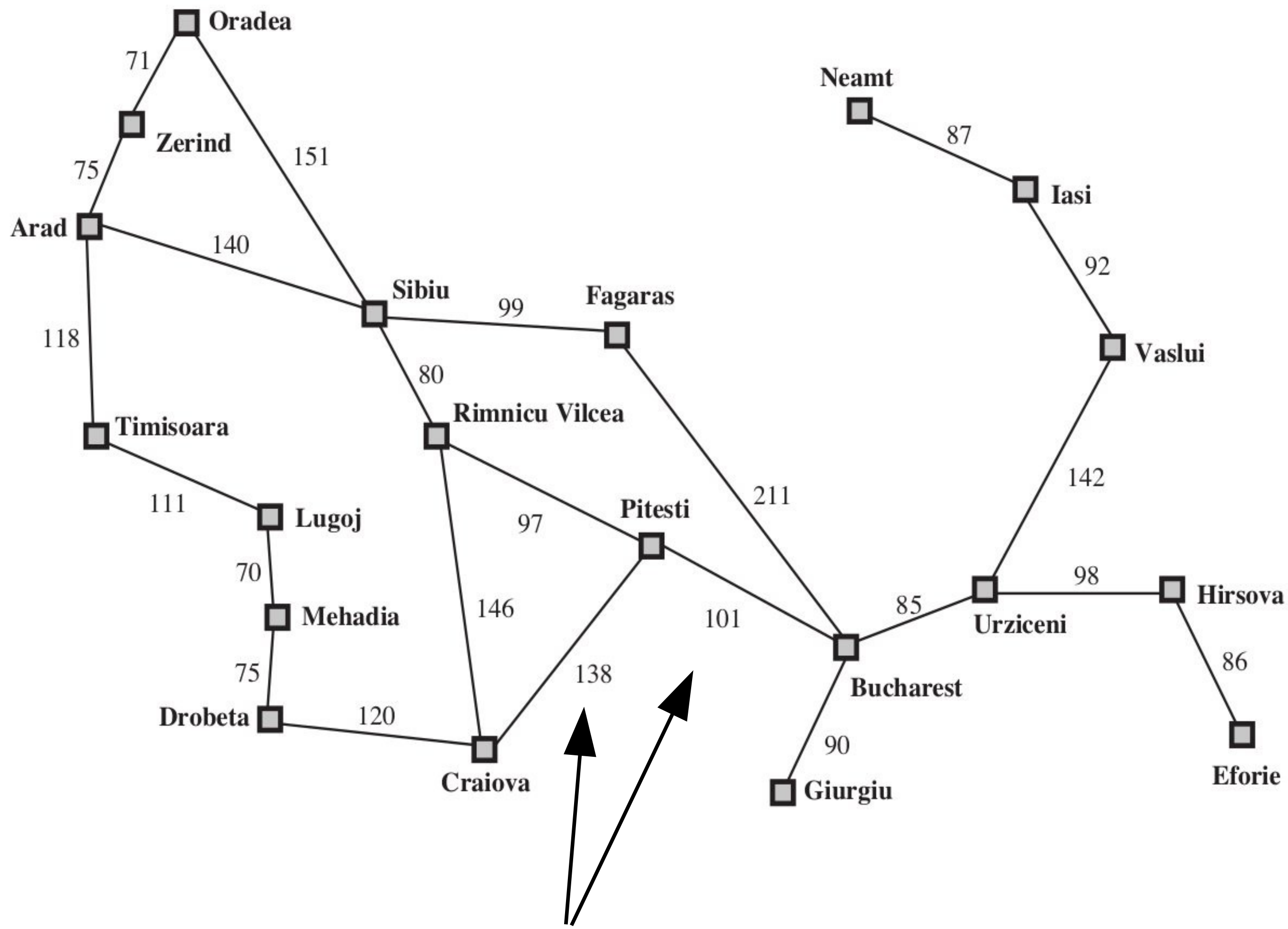
Notice the distances between cities
– does BFS take these distances into account?

Uniform Cost Search (UCS)



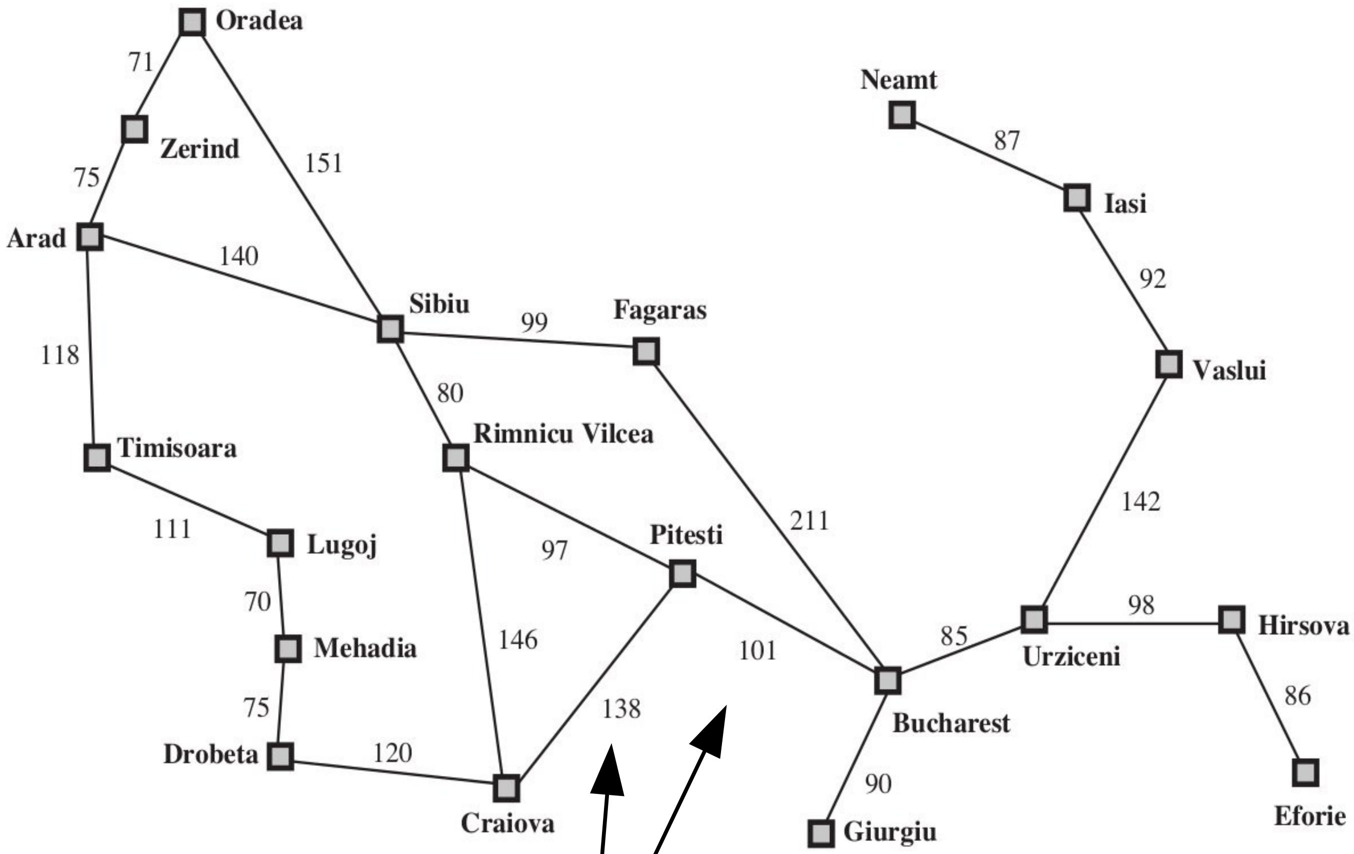
Notice the distances between cities
– does BFS take these distances into account?
– does BFS find the path w/ shortest milage?

Uniform Cost Search (UCS)



Notice the distances between cities
– does BFS take these distances into account?
– does BFS find the path w/ shortest milage?
– compare S-F-B with S-R-P-B. Which costs less?

Uniform Cost Search (UCS)



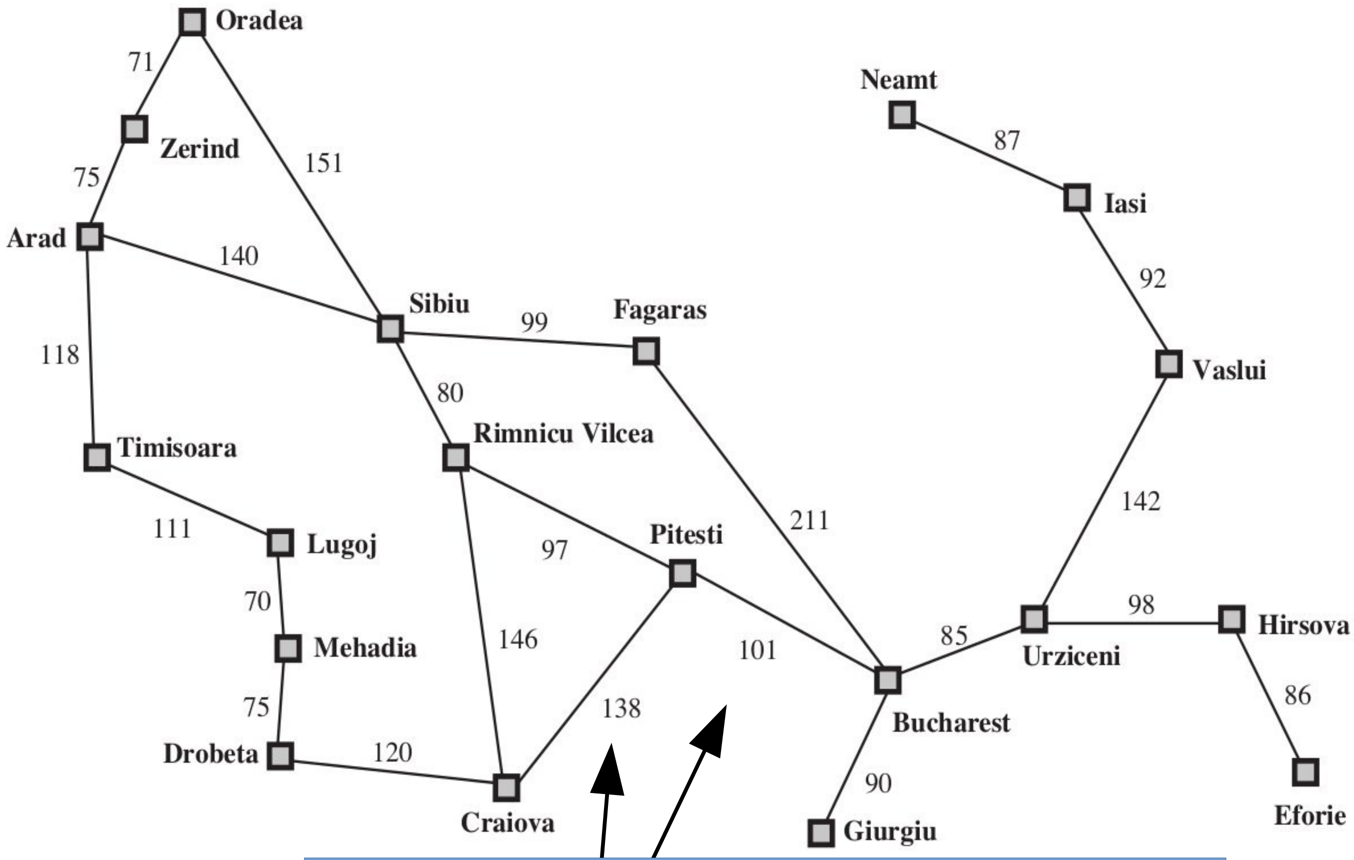
Notice
– doe
– doe
– com

How do we fix this?

nt?

less?

Uniform Cost Search (UCS)



Notice

– doe … nt?

– doe …

– com … less?

How do we fix this?
<u>UCS!</u>

Uniform Cost Search (UCS)

Same as BFS except: expand node w/ smallest path cost

Length of path

Uniform Cost Search (UCS)

Same as BFS except: expand node w/ smallest path cost

Length of path

Cost of going from state *A* to *B*: $c(A, B)$

Minimum cost of path going from start state to *B*: $g(B)$

Uniform Cost Search (UCS)

Same as BFS except: expand node w/ smallest path cost

Length of path

Cost of going from state *A* to *B:* $c(A, B)$

Minimum cost of path going from start state to *B:* $g(B)$

BFS: expands states in order of hops from start

UCS: expands states in order of $g(s)$

Uniform Cost Search (UCS)

Same as BFS except: expand node w/ smallest path cost

Length of path

Cost of going from state *A* to *B:* $c(A, B)$

Minimum cost of path going from start state to *B:* $g(B)$

BFS: exp

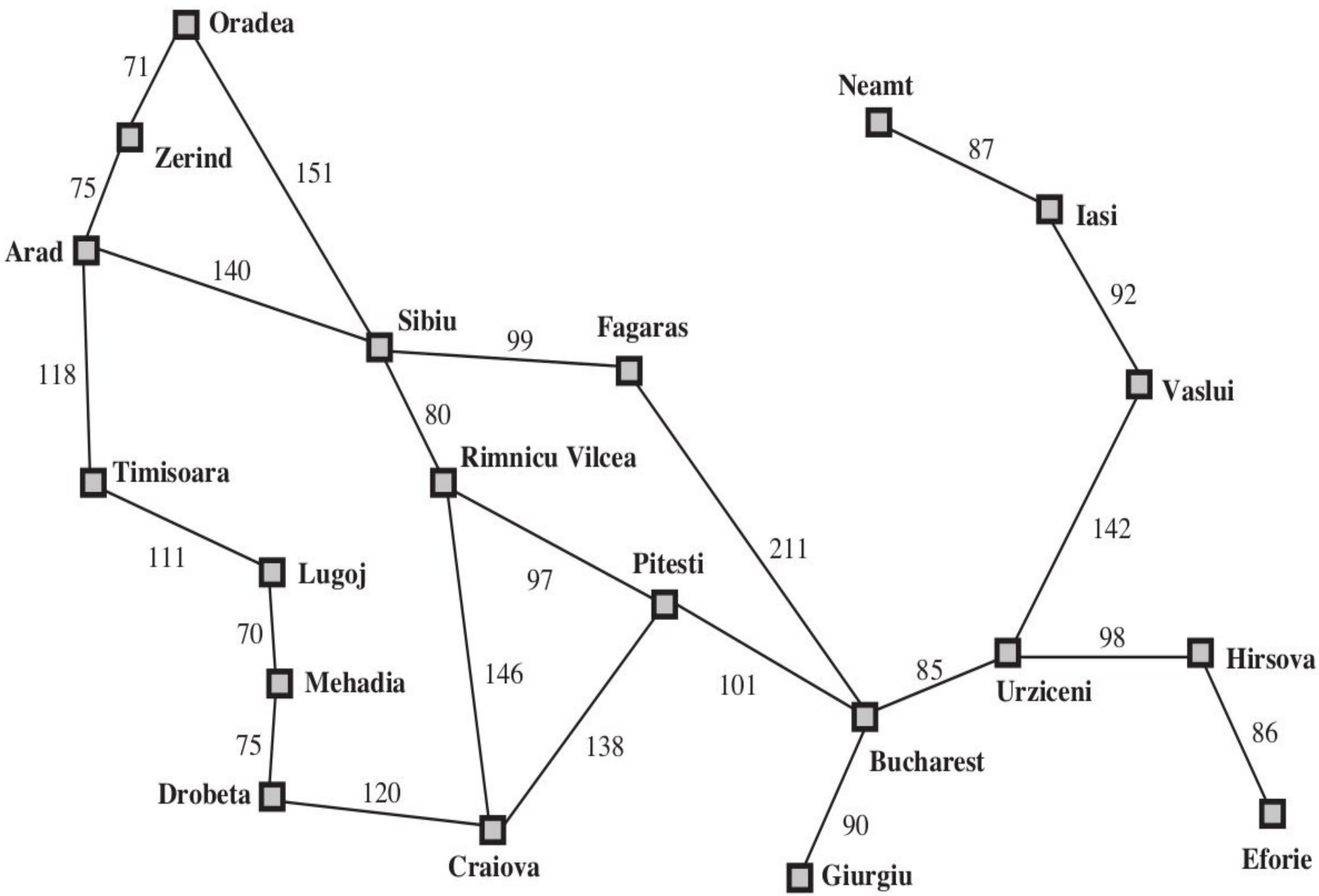
UCS: ex

How?

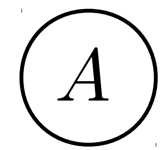
Uniform Cost Search (UCS)

Simple answer: change the FIFO to a priority queue
– the priority of each element in the queue is its path cost.

Uniform Cost Search (UCS)



UCS

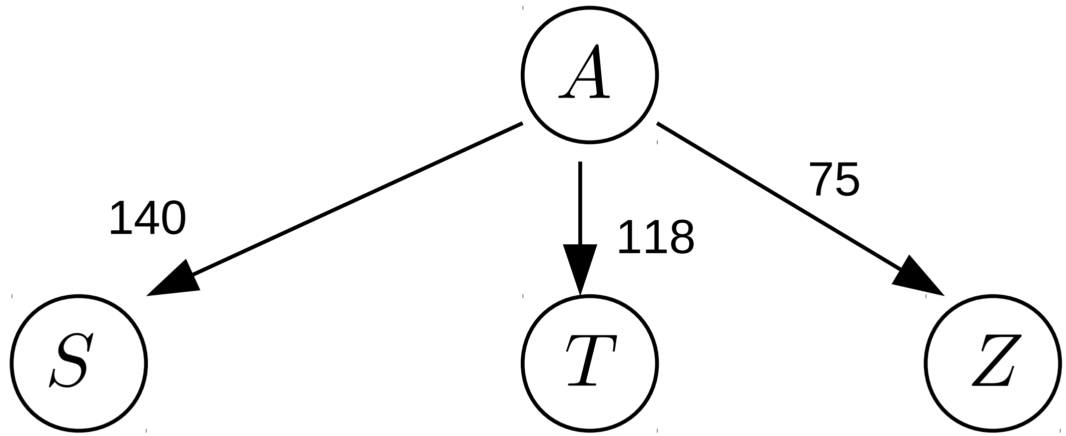


Fringe	Path Cost
A	0

Explored set:

UCS

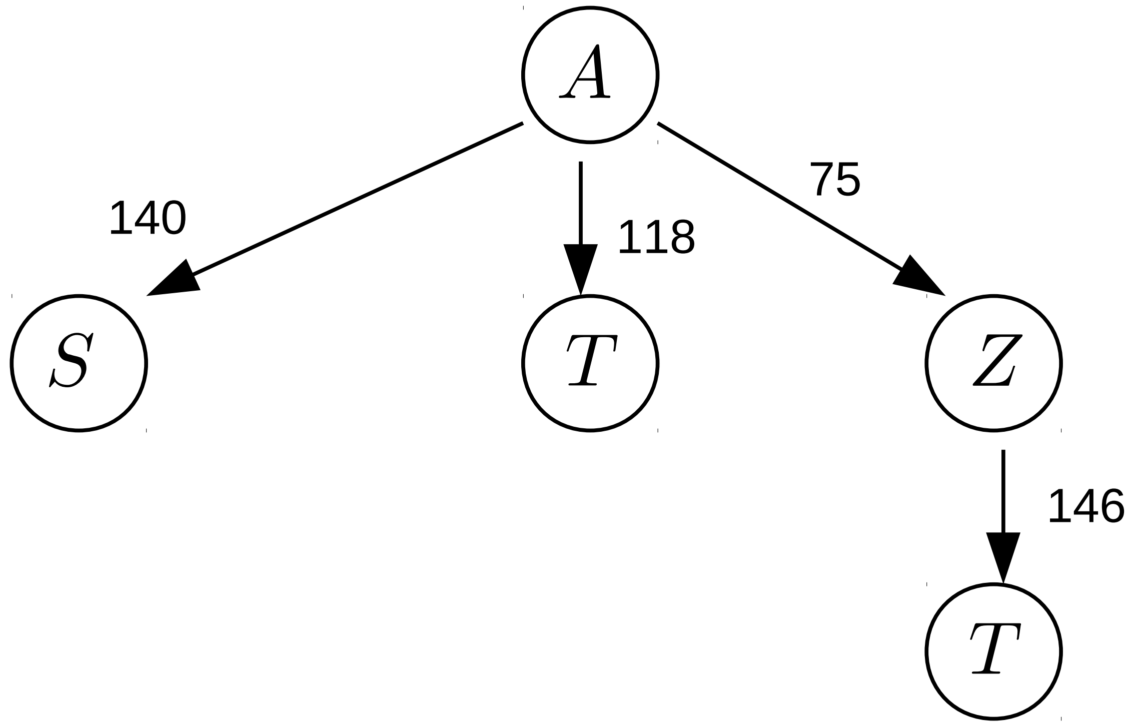
Fringe	Path Cost
~~A~~	~~0~~
S	140
T	118
Z	75



Explored set: A

UCS

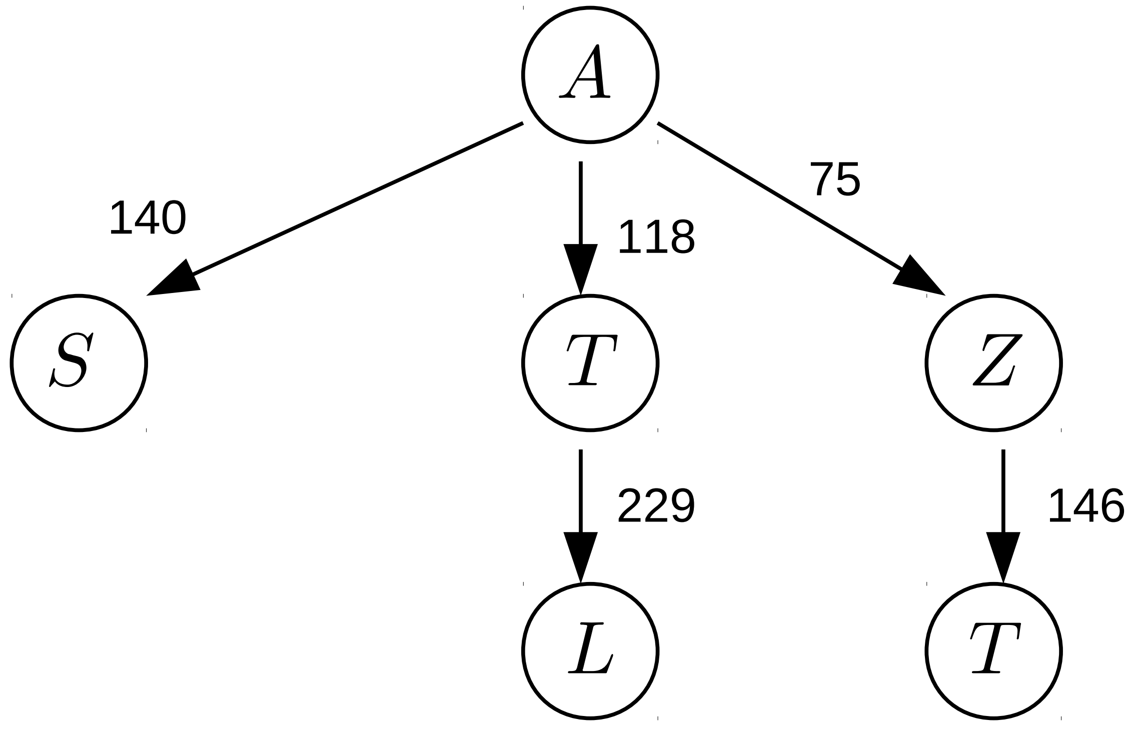
Fringe	Path Cost
~~A~~	~~0~~
S	140
T	118
~~Z~~	~~75~~
T	146



Explored set: A, Z

UCS

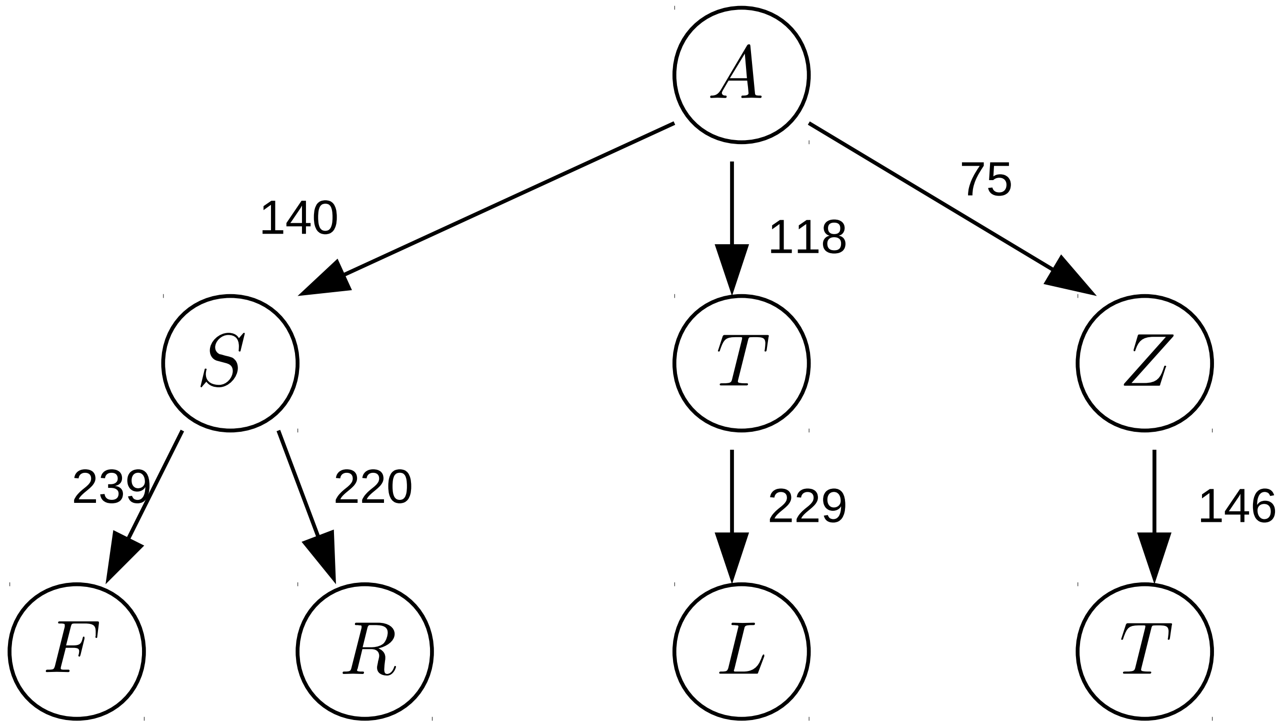
Fringe	Path Cost
~~A~~	~~0~~
S	140
~~T~~	~~118~~
~~Z~~	~~75~~
T	146
L	229



Explored set: A, Z, T

UCS

Fringe	Path Cost
~~A~~	~~0~~
~~S~~	~~140~~
~~T~~	~~118~~
~~Z~~	~~75~~
T	146
L	229
F	239
R	220



Explored set: A, Z, T, S

UCS

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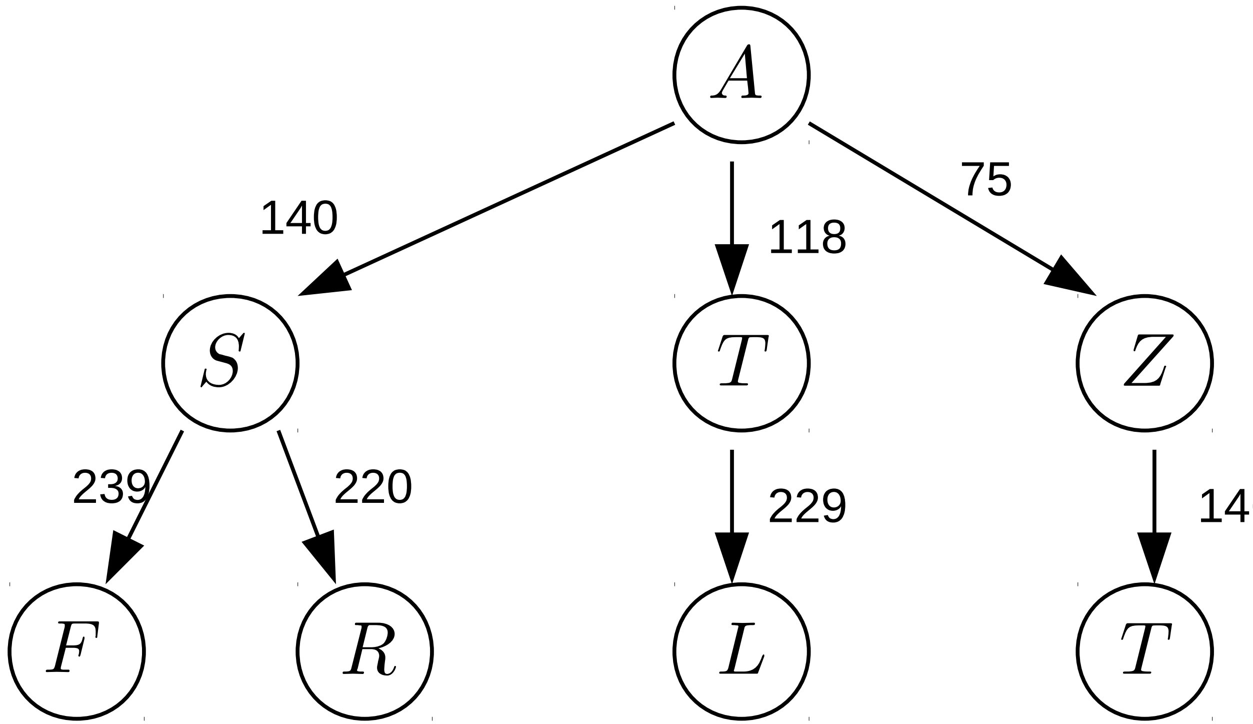

Explored set: A, Z, T, S

UCS

Fringe	Path Cost
~~A~~	~~0~~
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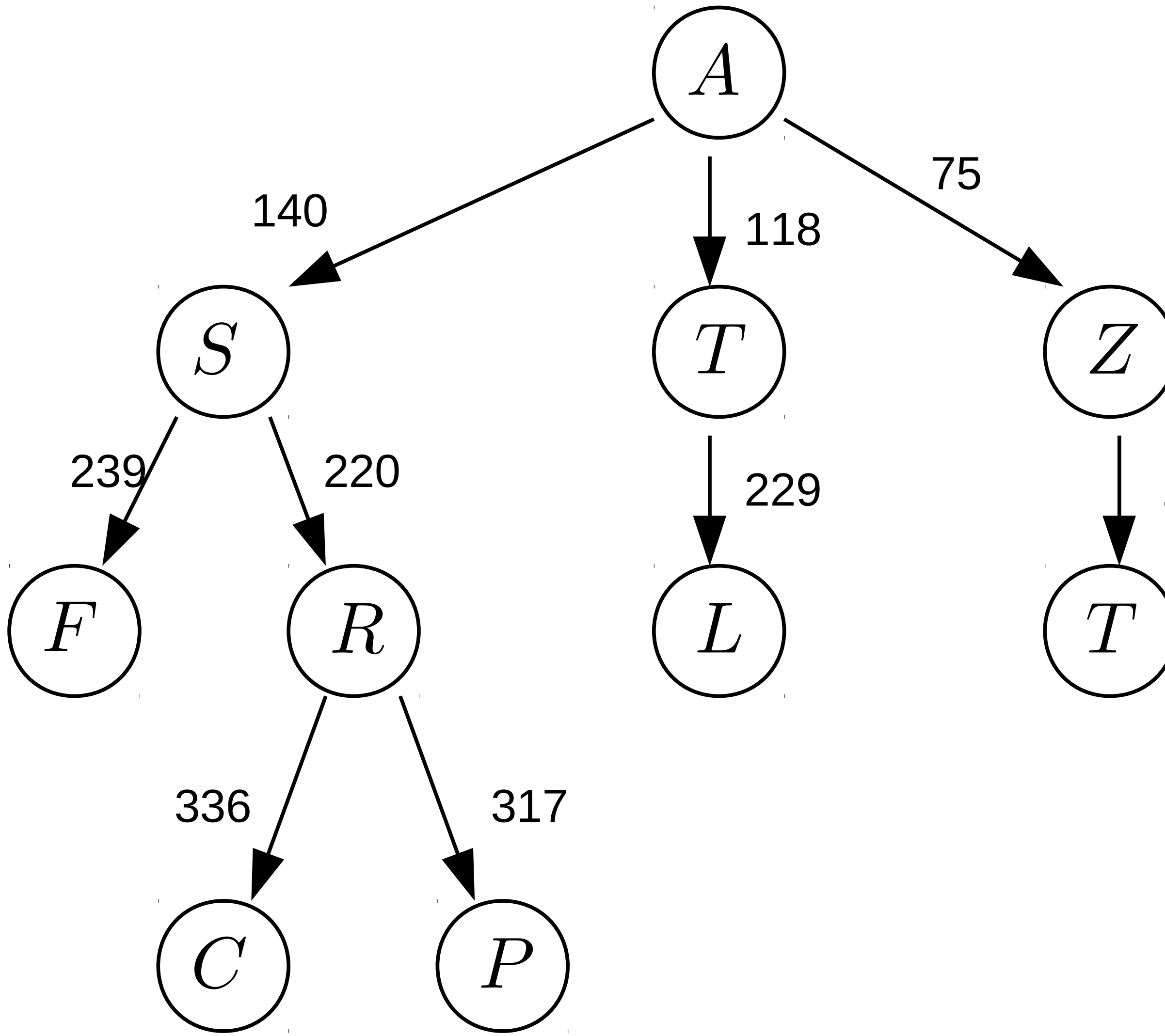
~~T~~	~~118~~
~~Z~~	~~75~~
~~T~~	~~146~~
L	229
F	239
~~R~~	~~220~~
C	336
P	317



Explored set: A, Z, T, S, R

UCS

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~~S~~	~~140~~
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~~Z~~	~~75~~
~~T~~	~~146~~
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F	239
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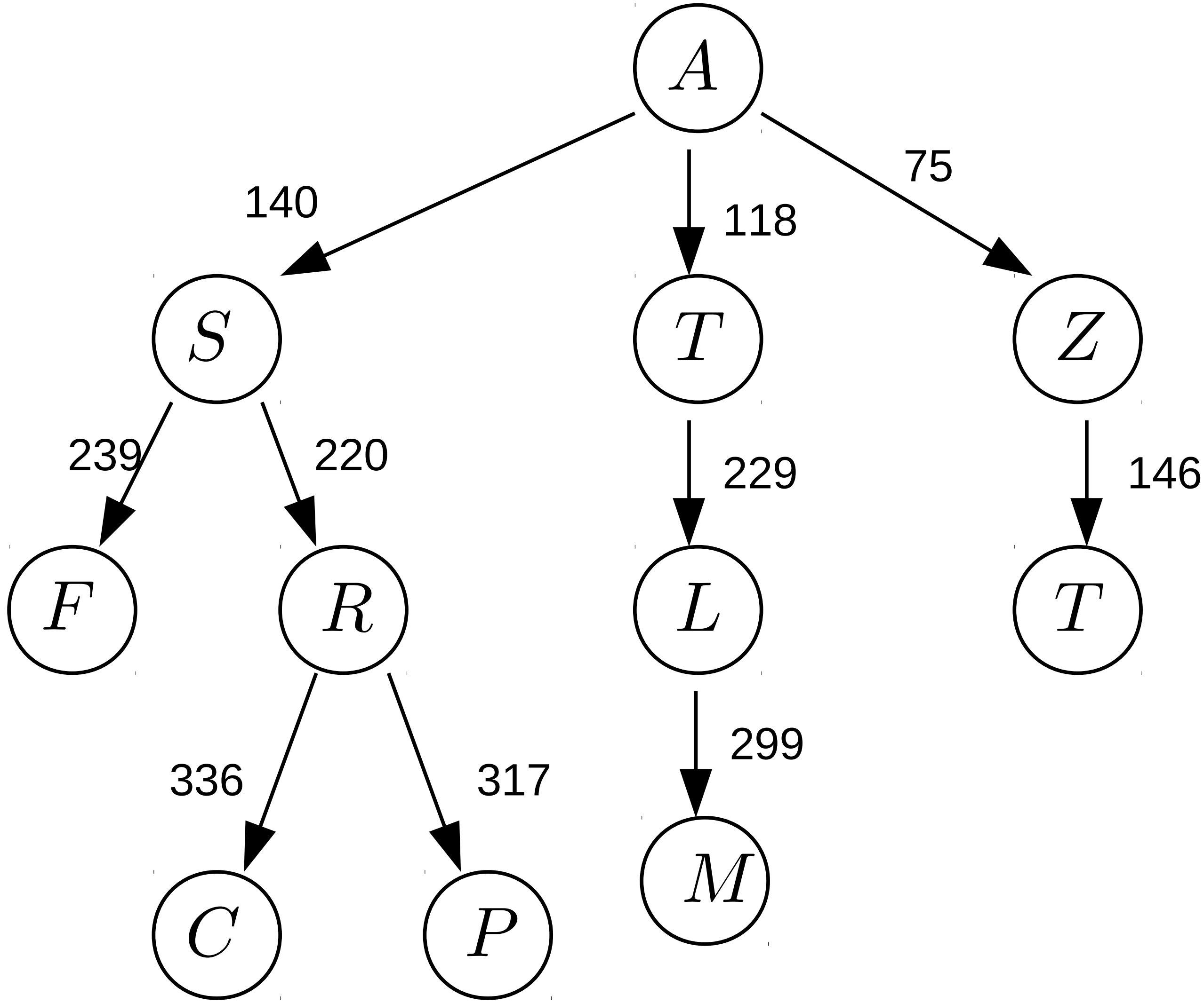
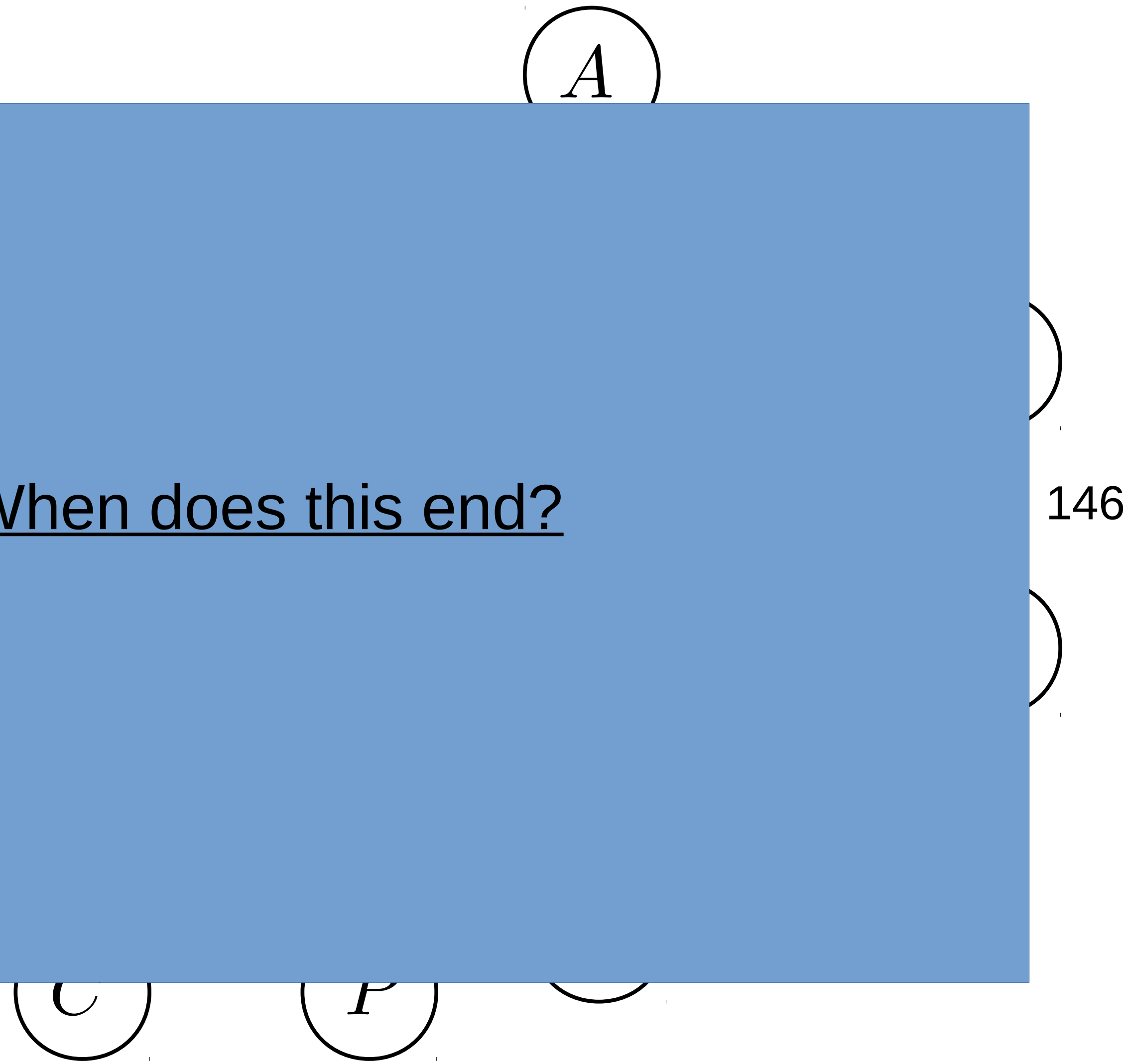
C	336
P	317
M	299



Explored set: A, Z, T, S, R, L

UCS



When does this end?

Fringe

- ~~A~~
- ~~S~~
- ~~T~~
- ~~Z~~
- ~~T~~
- ~~L~~
- F
- ~~R~~
- C
- P
- M

Explored set: A, Z, T, S, R, L

UCS

Fringe

~~A~~
~~S~~
~~T~~
~~Z~~
~~T~~
~~L~~
F
~~R~~
C
P
M

When does this end?

– when the goal state is removed from the queue

Explored set: A, Z, T, S, R, L

UCS

Fringe

~~A~~
~~S~~
~~T~~
~~Z~~
~~T~~
~~L~~
F
~~R~~
C
P
M

When does this end?

– when the goal state is removed from the queue

– NOT when the goal state is expanded

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Explored set: A, Z, T, S, R, L

UCS

```
function UNIFORM-COST-SEARCH(problem) returns a solution, or failure
  node ← a node with STATE = problem.INITIAL-STATE, PATH-COST = 0
  frontier ← a priority queue ordered by PATH-COST, with node as the only element
  explored ← an empty set
  loop do
    if EMPTY?(frontier) then return failure
    node ← POP(frontier)   /* chooses the lowest-cost node in frontier */
    if problem.GOAL-TEST(node.STATE) then return SOLUTION(node)
    add node.STATE to explored
    for each action in problem.ACTIONS(node.STATE) do
      child ← CHILD-NODE(problem, node, action)
      if child.STATE is not in explored or frontier then
        frontier ← INSERT(child, frontier)
      else if child.STATE is in frontier with higher PATH-COST then
        replace that frontier node with child
```

Figure 3.14 Uniform-cost search on a graph. The algorithm is identical to the general graph search algorithm in Figure 3.7, except for the use of a priority queue and the addition of an extra check in case a shorter path to a frontier state is discovered. The data structure for *frontier* needs to support efficient membership testing, so it should combine the capabilities of a priority queue and a hash table.

UCS Properties

Is UCS complete?
– is it guaranteed to find a solution if one exists?

What is the time complexity of UCS?
– how many states are expanded before finding a sol'n?
- – b: branching factor
- – C*: cost of optimal sol'n
- – e: min one-step cost
- – complexity = $O(b^{C^*/e})$

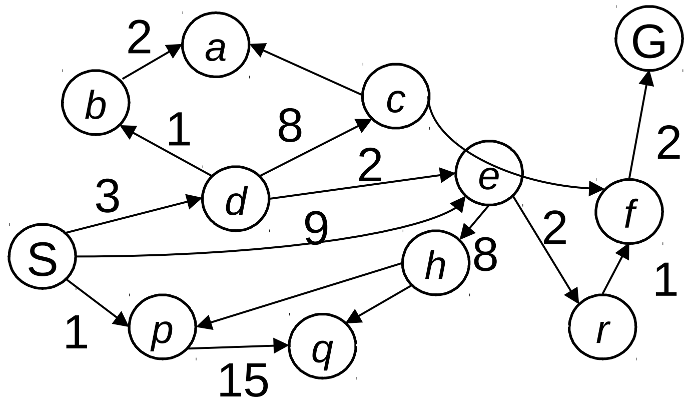
What is the space complexity of BFS?
– how much memory is required?
- – complexity = $O(b^{C^*/e})$

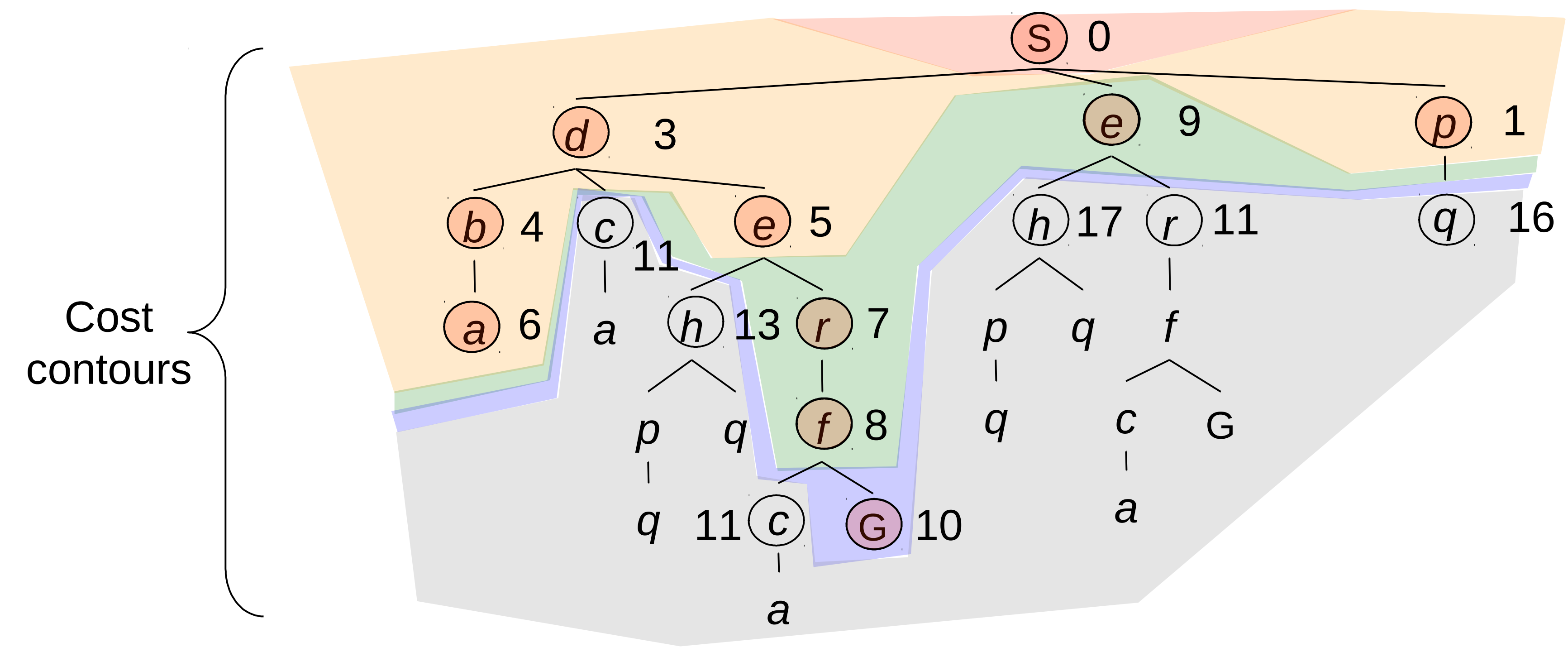
Is BFS optimal?
– is it guaranteed to find the best solution (shortest path)?

UCS vs BFS

Strategy: expand a cheapest node first:

Fringe is a priority queue (priority: cumulative cost)





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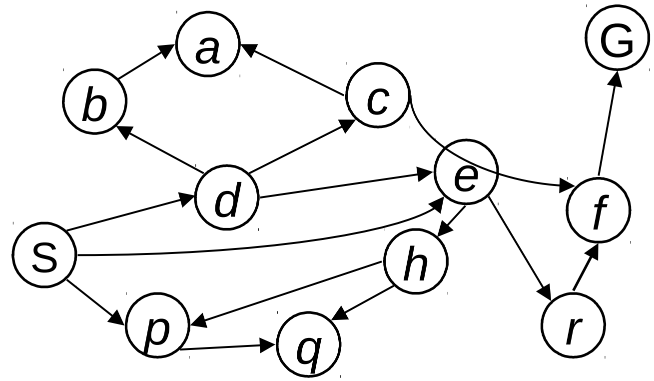
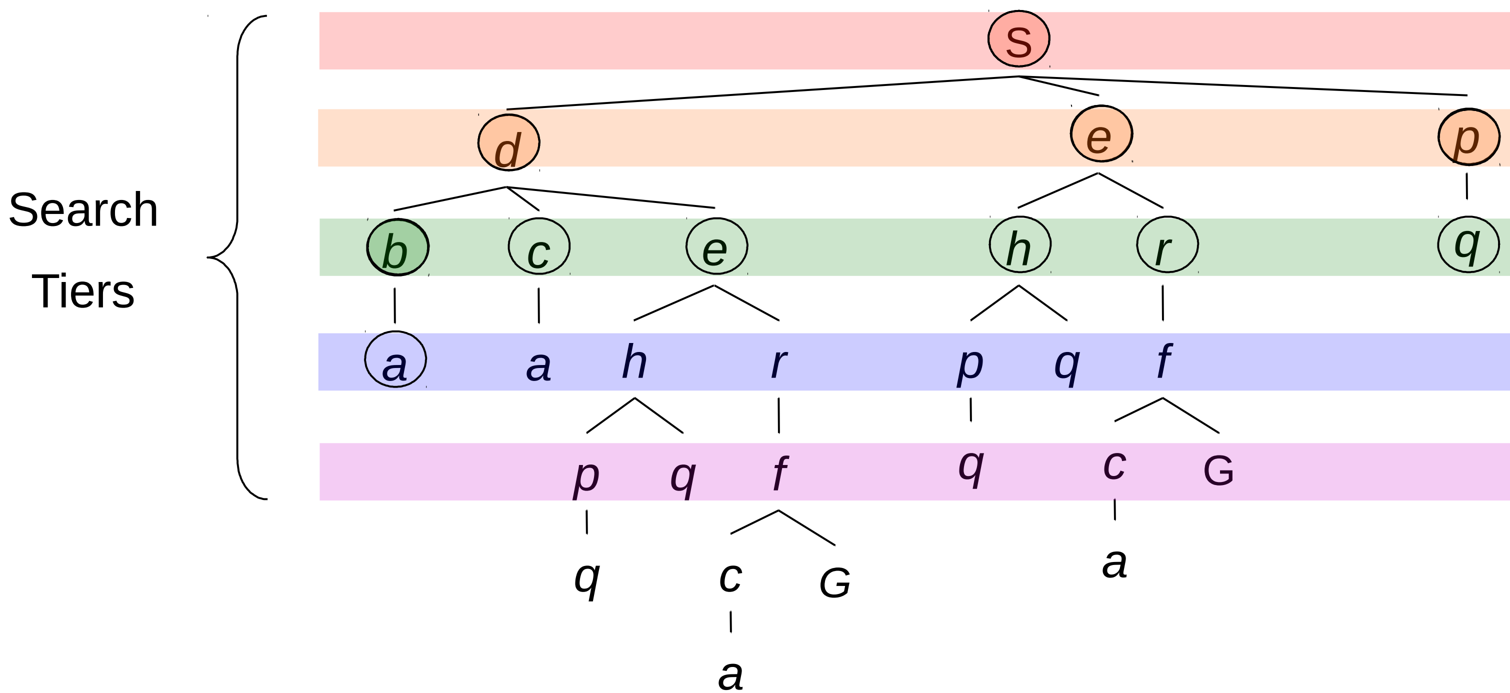
UCS vs BFS

Strategy: expand a shallowest node first

Implementation: Fringe is a FIFO queue





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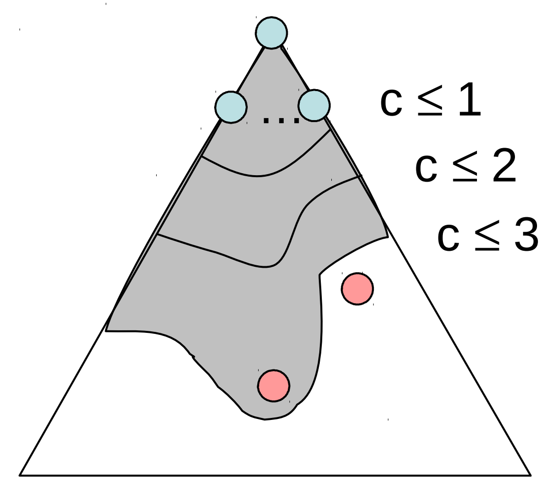
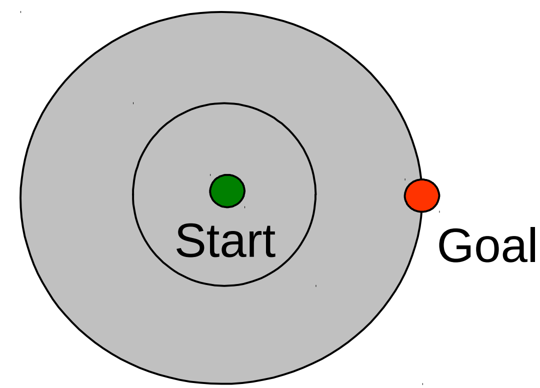
UCS vs BFS

- Remember: UCS explores increasing cost contours
- The good: UCS is complete and optimal!
- The bad:
  - Explores options in every “direction”
  - No information about goal location
- We’ll fix that soon!





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