2 ADTs and Algebraic Specifications

One of the first questions we are faced with when trying to design a piece of software is “What is the data that the program must manipulate? What are the objects? For example, when designing a game or a simulation, likely data includes artifacts such as spaceships, planets, and stars. For an animation package, it may include pictures, characters, backgrounds. These are all fairly concrete data. Other forms of data are more abstract — for instance, trajectories for animated characters, or behaviors for interactive animations.

For the purpose of our discussion today, let’s focus on one kind of data that arises in an animation package, namely drawings. To keep the discussion simple, I will consider very simple kind of drawings, namely two-dimensional line drawings.

So how do we go about thinking about drawings? When designing software, it’s best to keep an open mind as to the exact form the software will take. We can simply assume that a drawing is an object without committing to anything else. In particular, we will not want to commit to a particular way of representing a drawing. Let’s embody this into a principle: you’ll like it, it’s all about not doing work.

The Principle of Least Commitment: Don’t commit yourself any more or sooner than necessary.

So drawings are objects of some kind. It does not really matter what objects are. What matters about objects is how they behave. So let’s ask the question: how should drawings behave?

That’s a vague question. Let’s refine it somewhat, and look for something specific. Behaviors are induced when objects are acted upon. So how do we want to act on drawings? In other words, what operations do we want to support on drawings?

Here are some sensible operations on drawings: creating new drawings, adding a line to a drawing, removing a line from a drawing, moving a in the drawing, loading and saving a drawing to disk, querying the content of a point in the drawing. You should be easily able to think of others.

Operations on drawings naturally divide into three kinds of operations. The first kind of operations are those that create a drawing, which we will call *creators* or sometimes constructors. The second kind of operations are those that extract information from drawings, which we will call *accessors* or *selectors*. When an accessor returns true or false we often
call the accessor a *predicate*. The final kind of operations are those that do not extract information or create new drawings, but rather effect a change in the real world, such as saving a drawing to disk and physically changing the content of the disk, which we will call *effectors*. Effectors are those operations that actually *do* something that you as a human being sitting at the computer notice. Printing out information on the screen for you to read is an effector operation. They cannot be undone.

For now, we will keep the design exceedingly small and simple. In particular, we’ll leave out some behaviors and operations that we will need later. But that’s okay. Part of the point of object-oriented design and programming is that it makes it easy to add behaviors to objects in the future.

I will also most likely make mistakes, something voluntarily, sometimes not. Between that and needing to add operations later on, we will certainly need to revise our design in the future. Revising a design is must less expensive that revising a program, because a program accumulates all sort of cruft that a design does not have, including implementation choices. Moreover, a good design is much easier to turn into correct code than a bad design. Worse, a buggy design makes it impossible to write correct code. Thus, it pays to get the design right.

The concept of data with an associated set of operations is important enough that we’ll give it a name: an *abstract data type* (ADT for short):

**Abstract Data Type**: An abstract data type is a set of data and a set of operations that can be performed on the data, along with a description of what those operations do.

Let’s consider an ADT for drawings. First, the operations. What operations should we want to support on drawings? First of all, we need to create drawings — we want creators. There are several choices possible, and here are mine. I want an operation `empty` that creates an empty drawing, that is, a drawing without any line in it. This is useful, in the same way that zero is useful in arithmetic. For the sake of simplicity, let’s assume that drawings have a fixed size. (It’s easy to make size a parameter.) I want an operation `oneLine` that creates a drawing containing a single line in it. I also want to create more complex drawings, recursively. To do so, I want an operation `merge` that creates a drawing by merging together two drawings, so that the lines in both drawings are included in the resulting drawing. (Other choices of creator operations are possible and equivalent to these three; as an exercise, try to come up with a few.)

Being able to create drawings is a good first step. But we may also want to extract information from drawings — we want accessors. Here are some operations that do so. First off, an operation `isEmpty` is useful to check if a drawing is empty. To extract the lines in a non-empty drawing, I want operations `firstLine` that returns the first line in a drawing, and `restDrawing` that returns a drawing containing all but the first line of a given drawing.
Signature of the Drawing ADT

A *signature* consists of the names of the operations together with their type. For the Drawing ADT, a reasonable signature would be as follows:

- `empty : -> Drawing`
- `oneLine: Line -> Drawing`
- `merge : Drawing Drawing -> Drawing`
- `isEmpty : Drawing -> boolean`
- `firstLine : Drawing -> Line`
- `restDrawing : Drawing -> Drawing`

This signature assumes that we have a type for lines which I will just assume is an ADT itself called `Line`, well as a type for Boolean values.

Notice that the creators all return a `Drawing` object, as expected, while all the accessors take a `Drawing` object as an argument.

A signature describes one aspect of the interface, namely, what shape the operations have, that is, what they expect as arguments and what they return as a result. The signature gives no clue as to how those operations are meant to behave, however. We need to remedy that situation. This is the second part of the definition of ADT.

Specification of the Drawing ADT

A *specification* (or spec, for short) is a kind of guarantee (or contract) between clients and implementors.

Clients

- depend on the behavior guaranteed by the spec, and
- promise not to depend on any behavior not guaranteed by the spec.

Implementors

- guarantee that a provided abstraction behaves as specified by the spec, and
- do not guarantee any behavior not covered by the spec.

It is hard to specify how objects behave. Usually, this is done in English, in an informal way. But the resulting specification is often incomplete, incorrect, ambiguous, or confusing. You’ll see examples of those very often.
Let me introduce a formal way of specifying behavior, as a set of algebraic equations that the operations of the interface must obey. Because of that, we will call it an algebraic specification. (There are other ways of specifying behavior, which we may get to before the end of the course.)

The basic rule is to describe how each accessor works on any object constructed using the creators.

Specifying \texttt{isEmpty} is straightforward:

\begin{align*}
\texttt{isEmpty (empty ())} &= \text{true} \\
\texttt{isEmpty (oneLine (1))} &= \text{false} \\
\texttt{isEmpty (merge (d1,d2))} &= \begin{cases} 
\text{true} & \text{if isEmpty(d1)=true and isEmpty(d2)=true} \\
\text{false} & \text{otherwise}
\end{cases}
\end{align*}

Specifying \texttt{firstLine} is just a bit more interesting:

\begin{align*}
\texttt{firstLine (oneLine (l))} &= l \\
\texttt{firstLine (merge (d1,d2))} &= 
\begin{cases} 
\text{firstLine (d2)} & \text{if isEmpty(d1)=true} \\
\text{firstLine (d1)} & \text{otherwise}
\end{cases}
\end{align*}

First off, there is no equation describing how \texttt{firstLine} behaves when applied to an empty drawing. That’s on purpose: no behavior is specified, because it should be an error. The implementor is free to do as she wishes here. (In general, she will report an error through an exception or a similar mechanism, but it may make sense in some situations to silently ignore the error and proceed with the rest of the computation.)

Second, the equation for \texttt{firstLine} applied to a merged drawing is a conditional equation, because it depends on properties of the first drawing, namely whether it is empty or not.

Specifying \texttt{restDrawing} is similar to \texttt{firstLine}, with many of the same subtleties:

\begin{align*}
\texttt{restDrawing (oneLine (l))} &= \text{empty ()} \\
\texttt{restDrawing (merge (d1,d2))} &= 
\begin{cases} 
\text{restDrawing (d2)} & \text{if isEmpty(d1)=true} \\
\text{merge (restDrawing (d1),d2)} & \text{otherwise}
\end{cases}
\end{align*}

These equations seem only to specify the behavior of the accessors, but in fact, they describe the interaction between the accessors and the creators, and thereby implicitly also specify the behavior of the creators.

Let me emphasize again: the above specification describes how the operations behave, what they do. And in particular, they say a lot more than they seem to say. Look at how \texttt{firstLine} and \texttt{restDrawing} are defined. They assume that there is an order to the lines in
a drawing (what that order actually is is irrelevant), and that when you merge two drawings, 
the order of the lines is preserved, in the sense that the lines in the first drawing all come 
before the lines in the second drawing in the resulting merged drawing. Putting it another 
way, the description of the operations assumes that the lines in a drawing are put in a list, 
and merging two drawings corresponds to appending their list of lines. That may or may 
not be what you had in mind when I informally described how drawings behaved, but that’s 
irrelevant. The above is the specification, and the specification tells you the contract I expect 
the implementors of the Drawing ADT to abide to.

Exercise: Define a different Drawing ADT such that the operation `firstLine` returns the 
“upper leftmost” line in the drawing, and `restDrawing` returns the drawing obtained by 
removing the “upper leftmost” line. Take whatever sensible definition of “upper leftmost 
line” you can come up with. This is a very challenging exercise, and more intended to get 
you to think about what goes into a specification than anything else. And yes, you probably 
have to assume some things about the Line ADT as well.

Note that I have not said how drawings are implemented. And I don’t care at this point. I 
do not care how the operations do what they claim to do, I am just describing how I want 
them to behave. Figuring out how to implement objects, how to go from a description of 
their behavior to actual code that implements that behavior, is a decision that we will resist 
making for as long as possible, per the Principle of Least Commitment.

Despite this lack of description of how objects are implemented, the specification is still 
powerful enough to tell us the result of complex operations. For instance, suppose that I 
want to check what is the result of extracting the second line out of a drawing with three 
lines: (0,0)-(1,1), (1,1)-(0,1), and (0,1)-(0,0), in that order. (For simplicity, assume that a 
line is represented as (a,b)-(c,d).) Clearly, the result of that extraction should be the line 
(1,1)-(0,1). That is, we want to check that the result of the following expression is the line 
(1,1)-(0,1):

```
oneLine (restDrawing (merge (oneLine ((0,0)-(1,1)),
                        merge (oneLine ((1,1)-(0,1)),
                                oneLine ((0,1)-(0,0)))))
```

Well, I can use the equations above to replace equals by equals and simplify the above 
expression, just like you would do in algebra. This is were the name algebraic specification 
comes from, by the way. Here is one possible derivation. See if you can spot the equations I 
used at each step:
This is, of course, the result that we expected. But the point is, the whole point is, we can compute the result *without having ever said a word about how drawings are implemented!"