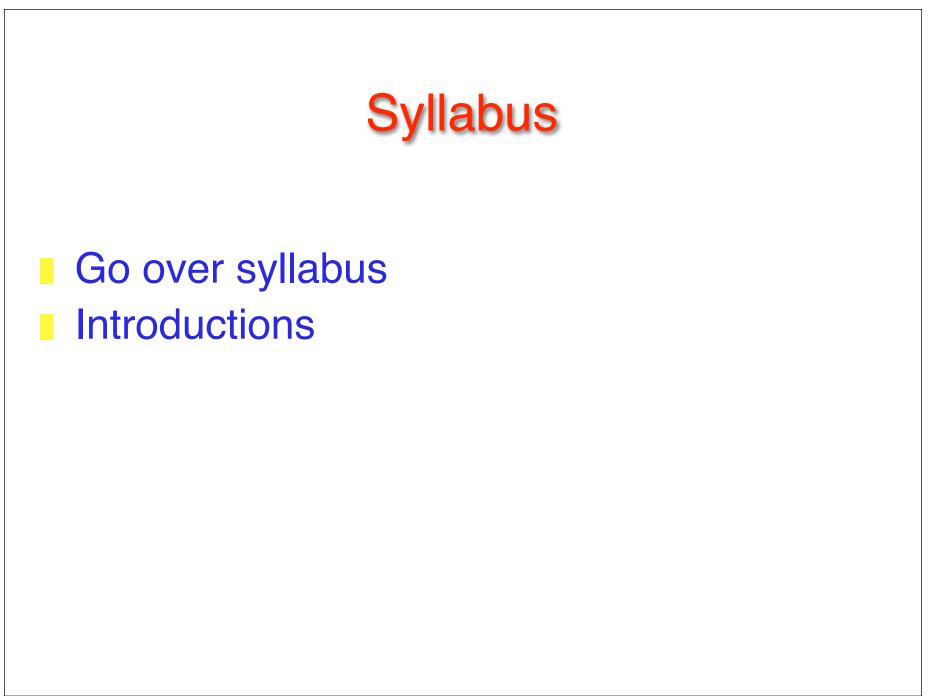
Introduction

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Formal Methods, Lecture 1

September 2008



Motivation

Instead of debugging a program, one should prove that it meets its specifications, and this proof should be checked by a computer program.

John McCarthy
"A Basis for a Mathematical Theory of Computation," 1961

Boyer-Moore Theorem Provers

1970's

- Edinburgh Pure Lisp Theorem Prover (1973)
- A Computational Logic (1978)

1980's

NQTHM (1981)

ACL2 (1989) A Computational Logic for Applicative Common Lisp

1990's-Present

- Kaufmman joins as developer
- Workshops (7 already); huge regression suite

2000's:

- ACL2 books
- Development environments (ACL2 Sedan, Dracula)
 - 2005 ACM Software System Award (Boyer, Kaufmann, Moore)

Boyer-Moore Theorems Proved

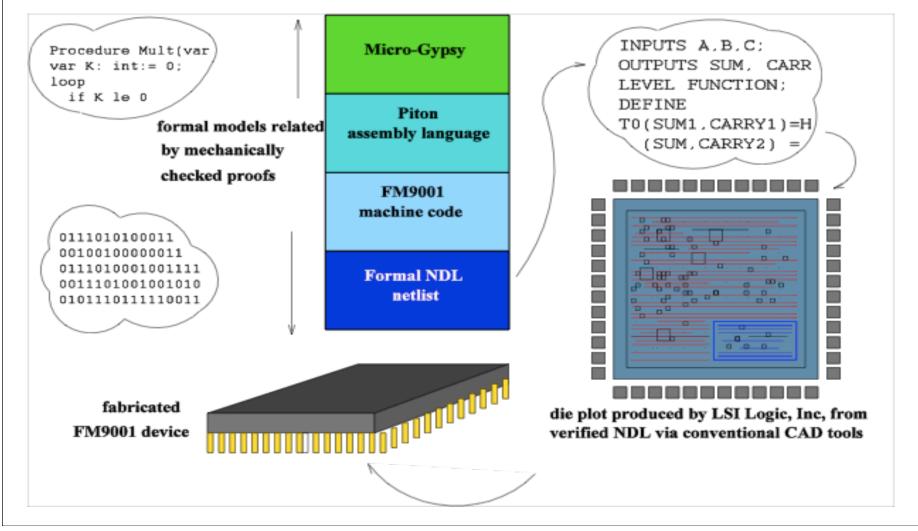
1970's: Simple List Processing

- Associativity of append
- Prime factorizations are unique

1980's: Academic Math & CS

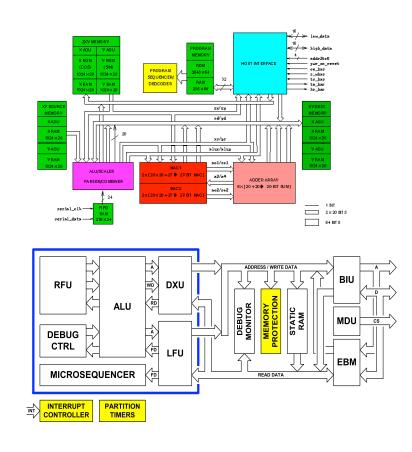
- Invertibility of RSA
- Undecidability of halting problem
- Gödel's First Incompleteness Theorem
- Gauss' Law of Quadratic Reciprocity
- CLI Stack:
 - Microprocessor
 - Assembler-linker-loader, Compiler, OS
 - High-level language

CLI Stack



1990's: Industrial Applications

- □ FDIV AMD Floating Point ...
- Motorola CAP DSP
 - □ Bit/cycle-accurate model
 - Run fasters than SPW model
 - Proved correctness of pipeline hazard detection in microcode
 - □ Verified microcode programs
- Rockwell Collins JEM1
- Rockwell Collins AAMP7
 - MILS EAL-7 certification from NSA for their crypto processor
 - Verified separation kernel



Hardware Verification: Motivation



International Technology Roadmap for Semiconductors, 2005 Edition.

Verification has become the dominant cost in the design process. In current projects, verification engineers outnumber designers, with this ratio reaching two or three to one

...

Without major breakthroughs, verification will be a non-scalable, showstopping barrier to further progress in the semiconductor industry.

• • •

The overall trend from which these breakthroughs will emerge is the shift from ad hoc verification methods to more structured, formal processes.

Hardware Verification Challenge

- Verification costs range from 30%-70% of design cost
- R&D for high-end CPU: 900+ team, costing ~ \$1B
- Bob Bently: FDIV bug would cost \$12B in 2005 terms
- The verification problem is getting worse
 - Nanotechnology: many inherently unreliable components
 - Multicores: concurrency, coherence, parallelism

ACL2s

ACL2 theorem prover

- Runs like a well-tuned race car in the hands of an expert
- Unfortunately, novices don't have the same experience
- Disseminate: wrote a book
- Not enough: undergrads
- ACL2s: The ACL2 Sedan
 - From race car to sedan
 - Self-teaching
 - Control a machine that is thinking about other machines
 - Visualize what ACL2 is doing
 - Levels & termination
 - Used in several classes
 - We'll use it in this class



ACL2 is ...



A programming language:

- Applicative, functional subset of Lisp
- Compilable and executable
- Untyped, first-order

ACL2 Universe

- The ACL2 universe, U consists of the following objectsAtoms
 - Numbers includes integers, rationals, and complex rationals
 - Examples include -1, 3/2, and #c(-1 2)
 - Characters represent the ASCII characters
 - Examples include #\2, #\a, and #\Space
 - Strings are finite sequences of characters
 - An example is "Hello World!"
 - Symbols consist of two strings:
 - a package name and a symbol name
 - For example, the symbol FOO::BAR has package name "FOO" and symbol name "BAR"

ACL2 Universe

- Conses are ordered pairs (binary trees) of objects
- The left component of a cons is called the car
- The right component is called the cdr
 - (1. "A")
 - (1 2 3)
 - ((A . 1) (B . 2) (C . 3))
- Notes
 - The symbols t and nil are used to denote true and false
 - Nil also denotes the empty list
 - Conses can be used to represent records and finite functions

ACL2 is ...



A programming language:

- Applicative, functional subset of Lisp
- Compilable and executable
- Untyped, first-order
- DEMO
- ACL2 Web page
- ACL2 Sedan Web page

ACL2 is ...



A programming language:

- Applicative, functional subset of Lisp
- Compilable and executable
- Untyped, first-order

A mathematical logic:

- First-order predicate calculus
- With equality, induction, recursive definitions
- Ordinals up to ε_0 (termination & induction)

Definitional Principle

- Functions are defined using defun
- Example: (defun succ (x) (+ x 1))
- The form of a defun is (defun $f \dots (x_1 \dots x_n)$ body), where:
 - *x*₁ ... *x*_n are distinct variable symbols
 - the free variables in body are in $x_1 \dots x_n$
 - ... are for documentation and declarations (optional)
 - *body* is a valid expression (history)
 - if *f* is recursive we must prove that it terminates
 - ACL2 logic is extended with the axiom $(f x_1 ... x_n) = body$