CS3000: Algorithms & Data Jonathan Ullman

Lecture 17: More Applications of Network Flow

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- Separate image into foreground and background
- We have some idea of:
 - whether pixel i is in the foreground or background
 - whether pair (i,j) are likely to go together

- Input:
 - an undirected graph G = (V, E); V = "pixels", E = "pairs"
 - likelihoods $a_i, b_i \ge 0$ for every $i \in V$
 - separation penalty $p_{ij} \ge 0$ for every $(i, j) \in E$
- Output:
 - a partition of V into (A, B) that maximizes

$$q(A,B) = \sum_{i \in A} a_i + \sum_{j \in B} b_j - \sum_{\substack{(i,j) \in E \\ \text{btw } A \text{ and } B}} p_{ij}$$

- Differences between SEG and MINCUT:
 - SEG asks us to maximize, MINCUT asks us to minimize

- Differences between SEG and MINCUT:
 - SEG allows any partition, MINCUT requires $s \in A, t \in B$

- Differences between SEG and MINCUT:
 - SEG has edges between A and B, MINCUT considers edges from A to B



- Differences between SEG and MINCUT:
 - SEG has terms for each node in A,B, MINCUT only has terms for edges from A to B



- How should the reduction work?
 - capacity of the cut should correspond to the function we're trying to minimize



Step 1: Transform the Input



Step 2: Receive the Output



Step 3: Transform the Output





• correctness?

• running time?

Densest Subgraph



- Want to identify communities in a network
 - "Community": a set of nodes that have a lot of connections inside and few outside

Densest Subgraph

- Input:
 - an undirected graph G = (V, E)
- Output:
 - a subset of nodes $A \subseteq V$ that maximizes $\frac{2|E(A,A)|}{|A|}$

• Different objectives

• maximize
$$\frac{2|E(A,A)|}{|A|}$$
 vs. minimize $|E(A,B)|$

• Suppose
$$\frac{2|E(A,A)|}{|A|} \ge \delta$$
 and see what that implies
 $\Leftrightarrow 2|E(A,A)| \ge \delta|A|$
 $\Leftrightarrow \sum_{v \in A} \deg(v) - |E(A,B)| \ge \delta|A|$
 $\Leftrightarrow \sum_{v \in V} \deg(v) - \sum_{v \in B} \deg(v) - |E(A,B)| \ge \delta|A|$
 $\Leftrightarrow 2|E| - \sum_{v \in B} \deg(v) - |E(A,B)| \ge \delta|A|$
 $\Leftrightarrow \sum_{v \in B} \deg(v) + \delta|A| + |E(A,B)| \le 2|E|$
 $\Leftrightarrow \sum_{v \in B} \deg(v) + \sum_{v \in A} \delta + \sum_{e \text{ from } A \text{ to } B} 1 \le 2|E|$

$\sum_{v \in B} \deg(v) + \sum_{v \in A} \delta + \sum_{e \text{ from } A \text{ to } B} 1 \le 2|E|$