You can turn in handwritten solutions to this part of the assignment. Please write clearly and use standard-sized (8.5 by 11in) paper. Solutions should be submitted at the beginning of class on the due date.

Problem 3: Type Soundness

We saw that the simply-typed $\lambda$-calculus ($\lambda^\to$) has a sound type system because it preserves types and guarantees progress of well-typed terms. Thus, well-typed terms do not get stuck (i.e., evaluation is safe). Let us add pairs to the call-by-value simply-typed $\lambda$-calculus. (Note that the syntax below is different from what we used in class. It matches the syntax in TAPL, Chapter 9.)

Types $\tau ::= \ldots \mid \tau_1 \times \tau_2$

Terms $e ::= \ldots \mid \{e_1, e_2\} \mid e.1 \mid e.2$

Values $v ::= \ldots \mid \{v_1, v_2\}$

New evaluation rules:

\[
\begin{align*}
\frac{e_1 \to e'_1}{\{e_1, e_2\} \to \{e'_1, e'_2\}} \quad (E-Pair1) \\
\frac{e_2 \to e'_2}{\{v_1, e_2\} \to \{v_1, e'_2\}} \quad (E-Pair2) \\
\frac{e \to e'}{e.1 \to e'.1} \quad (E-Fst) \\
\frac{e \to e'}{e.2 \to e'.2} \quad (E-Snd) \\
\frac{}{\{v_1, v_2\}.1 \to v_1} \quad (E-FstPair) \\
\frac{}{\{v_1, v_2\}.2 \to v_2} \quad (E-SndPair)
\end{align*}
\]

New typing rules:

\[
\begin{align*}
\frac{\Gamma \vdash e_1 : \tau_1 \quad \Gamma \vdash e_2 : \tau_2}{\Gamma \vdash \{e_1, e_2\} : \tau_1 \times \tau_2} \quad (T-Pair) \\
\frac{\Gamma \vdash e : \tau_1 \times \tau_2}{\Gamma \vdash e.1 : \tau_1} \quad (T-Fst) \\
\frac{\Gamma \vdash e : \tau_1 \times \tau_2}{\Gamma \vdash e.2 : \tau_2} \quad (T-Snd)
\end{align*}
\]

For this problem, you must extend the proofs of progress and preservation for STLC ($\lambda^\to$)—as well as the proofs of lemmas that these rely on—to demonstrate type soundness for this extended language ($\lambda^{\to\times}$).

(a) State the inversion lemma.

(b) State and prove the canonical forms lemma.

(c) State the permutation and weakening lemmas.

(d) State and prove the substitution lemma.

(e) Prove the progress and preservation lemmas; their statements are as follows:

**Lemma (Progress):** If $\vdash e : \tau$ then either $e$ is a value or there exists some $e'$ such that $e \to e'$.

**Lemma (Preservation):** If $\vdash e : \tau$ and $e \to e'$, then $\vdash e' : \tau$.

Note: When proving preservation, use induction on the derivation of $e \to e'$.

Note: For the proof portions only of parts (b), (d), and (e), you do not need to show the cases involving functions, application, and function types.