CS7400

You can turn in handwritten solutions to this part of the assignment. Please write clearly and use standard-sized (8.5 by 11in) paper. Solutions should be submitted at the beginning of class on the due date.

Problem 3: Type Soundness We saw that the simply-typed λ -calculus (λ^{\rightarrow}) has a sound type system because it preserves types and guarantees progress of well-typed terms. Thus, well-typed terms do not get stuck (i.e., evaluation is *safe*). Let us add pairs to the call-by-value simply-typed λ -calculus. (Note that the syntax below is different from what we used in class. It matches the syntax in TAPL, Chapter 9.)

New evaluation rules:

$$\frac{e_1 \longrightarrow e'_1}{\{e_1, e_2\} \longrightarrow \{e'_1, e_2\}} \quad (E-PAIR1) \qquad \qquad \frac{e_2 \longrightarrow e'_2}{\{v_1, e_2\} \longrightarrow \{v_1, e'_2\}} \quad (E-PAIR2)$$
$$\frac{e \longrightarrow e'}{e.1 \longrightarrow e'.1} \quad (E-FST) \qquad \qquad \frac{e \longrightarrow e'}{e.2 \longrightarrow e'.2} \quad (E-SND)$$
$$\frac{\{v_1, v_2\}.1 \longrightarrow v_1}{\{v_1, v_2\}.2 \longrightarrow v_2} \quad (E-SNDPAIR)$$

New typing rules:

$$\begin{split} \frac{\Gamma \vdash e_1 : \tau_1 \qquad \Gamma \vdash e_2 : \tau_2}{\Gamma \vdash \{e_1, e_2\} : \tau_1 \times \tau_2} \quad (\text{T-PAIR}) \\ \\ \frac{\Gamma \vdash e : \tau_1 \times \tau_2}{\Gamma \vdash e.1 : \tau_1} \quad (\text{T-Fst}) \qquad \qquad \frac{\Gamma \vdash e : \tau_1 \times \tau_2}{\Gamma \vdash e.2 : \tau_2} \quad (\text{T-Snd}) \end{split}$$

For this problem, you must extend the proofs of progress and preservation for STLC (λ^{\rightarrow}) —as well as the proofs of lemmas that these rely on—to demonstrate type soundness for this extended language $(\lambda^{\rightarrow\times})$.

- (a) State the inversion lemma.
- (b) State and prove the canonical forms lemma.
- (c) State the permutation and weakening lemmas.
- (d) State and prove the substitution lemma.
- (e) Prove the progress and preservation lemmas; their statements are as follows:

Lemma (Progress): If $\vdash e : \tau$ then *either* e is a value *or* there exists some e' such that $e \longrightarrow e'$. **Lemma (Preservation):** If $\vdash e : \tau$ and $e \longrightarrow e'$, then $\vdash e' : \tau$.

Note: When proving preservation, use induction on the derivation of $e \longrightarrow e'$.

Note: For the proof portions only of parts (b), (d), and (e), you do not need to show the cases involving functions, application, and function types.