

CS1800

11/10 - Fri '11

Admin

- Hw6 due today 11:59
- Hw7 out, due 11/17 (78 pts)
- Exam #2 11/17
9AM-6PM
2 hour window
remote
review in
last recitation
next week

Agenda

1. Inaction overview
2. Example - mathy (inequality)
3. Examples - structural (graphs, sets)

1. Induction overview

- our "main" proof technique in CS1500
- used all the time (algo, s/w engineering)

When to use ...

Predicate $P(n)$

We could:

$P(1)$	$P(4)$	∞
$P(2)$	$P(5)$	
$P(3)$	$P(6) \dots$	

Induction Shortcut:

- Base case $P(1) \rightarrow$ specific value
↳ smallest value we care about
- Inductive step $P(k) \Rightarrow P(k+1)$

(IH) $P(k)$
use to show $P(k+1)$
 k is arbitrary

A couple mathy reminders:

- $(a+b)^2 = a^2 + 2ab + b^2$
- $(a+1)^2 = a^2 + 2a + 1$
- $a^{-b} = \frac{1}{a^b}$

$$\cdot 2^b \cdot 2^c = 2^{b+c}$$

$$\cdot (2^b)^c = 2^{bc}$$

$$\cdot n! = (n)(n-1)(n-2)\dots(1) \quad \text{by def}$$

$$\cdot \sum_{i=1}^n i = 1+2+3+\dots+n \quad \text{by def}$$

2. Induction Example - Inequality

- growth of functions (later in semester)
- ex: $3^n < n!$

3^n grows more slowly than $n!$

$$\text{Pred: } P(n) \quad 3^n < n!$$

$$\text{Logic: } \forall n \in \mathbb{Z} \quad \underline{n \geq 7} \Rightarrow P(n)$$

$$\text{Base case: } P(7) \quad 3^7 < 7!$$

$$2187 < 5040$$

✓ done!

Inductive Step

$$P(k) \Rightarrow P(k+1)$$

$$P(k) \quad 3^k < k! \quad \text{IH}$$

Goal:
 $P(k+1) \quad 3^{k+1} < (k+1)!$

- one step at a time!
- Only do two things

$$3^{k+1} = 3^k \cdot 3$$

$$< (k!) \cdot 3 \quad \text{by IH}$$

$$< (k!) \cdot \underline{(k+1)} \quad \hookrightarrow \text{bigger than 3}$$

$$= (k+1)! \quad \text{Done!}$$


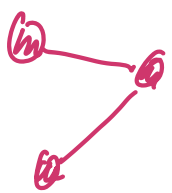
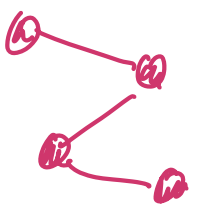
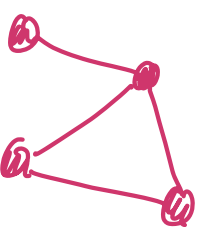
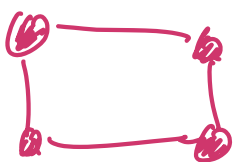
2. Induction Examples - graphs, sets

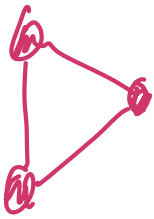
- be clear about what induction is on
 - # edges
 - # elements in set
 - # vertices
 - ...

- Inductive step: better to start at $k+1$ and remove something

$P(n)$
 \rightarrow # you chose

(ex) graph example (simple)

	<u># vertices</u>	<u># edges</u>	<u>total degree</u>
	2	1	2
	3	2	4
	4	3	6
	4	4	8
	4	4	8

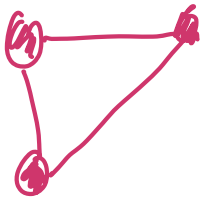


3

3

6

$\sum \text{#edges}$



add an edge, add 2 more total degree

Induction on $\# \text{edges}$

$P(n)$ a graph with n edges
has total degree $2n$

$\forall n \in \mathbb{Z}^+$ $P(n)$

Base case: $P(1)$



graph w/ one edge
by def, each vertex
has degree one

total degree = 2 ✓

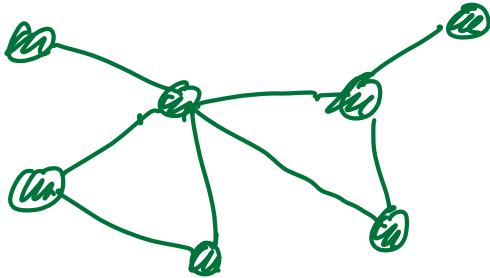
Inductive Step - $P(k) \Rightarrow P(k+1)$

$P(k)$ a graph with k edges has degree $2k$

IH

Goal: $P(k+1)$
graph w/ $k+1$ edges
has degree $2(k+1)$

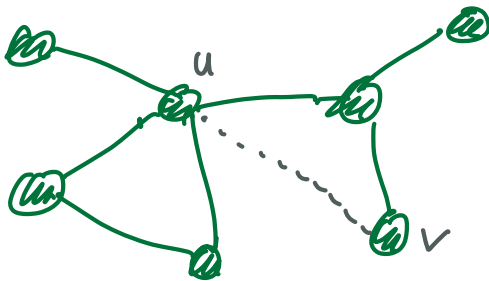
Graph with $k+1$ edges



(not real)

total degree: unknown

Remove an edge

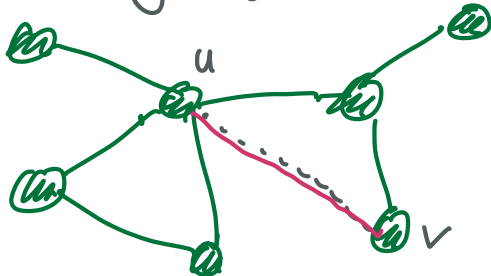


Graph with k edges (not real)

• has total degree $2k$

• removed edge (u, v)

Add missing edge back in



• added edge (u, v)

• Graph with $k+1$ edges

• added one degree for u

• added one degree for v

total degree = $2k + 2$

$= 2(k+1)$

done!

(ex) Induction on a set

• Subsets of size 2

$$A = \{a, b\}$$

Subsets of A w/ cardinality 2: $\{a, b\}$ (1)

$$A = \{a, b, c\}$$

Subsets of A w/ cardinality 2: $\{a, b\}, \{a, c\}, \{b, c\}$ (3)

$$A = \{a, b, c, d\}$$

Subsets of A w/ cardinality 2: $\{a, b\}, \{a, c\}, \{b, c\}, \{a, d\}, \{b, d\}, \{c, d\}$ (6)

A	#subsets
2	1
3	3
4	6

$$\begin{array}{l}
 P(2) \quad \frac{2 \cdot 1}{2} = 1 \quad P(4) \\
 P(3) \quad \frac{3 \cdot 2}{2} = 3 \quad \frac{4 \cdot 3}{2} = 6
 \end{array}$$

$P(n)$ Set with cardinality n

has $\frac{(n)(n-1)}{2}$ subsets of size 2

$$\forall n \in \mathbb{Z} \quad n \geq 2 \quad \Rightarrow P(n)$$

Base case $P(2)$

defn ✓
 $A = \{a_1, a_2\}$

has exactly one

subset of size 2, A

formula ✓

$$\frac{(2)(1)}{2} = 1$$

Inductive Step $P(k) \Rightarrow P(k+1)$

IH

$P(k)$

$|A| = k$, then # subsets
of size 2 is $\frac{(k)(k-1)}{2}$

Goal $P(k+1)$
 $|A| = k+1$, then
 $\frac{(k+1)(k)}{2}$

Let A be a set with $|A| = k+1$

$\{a_1, a_2, a_3, \dots, a_{k+1}\}$

• remove one element from A

• $A' = A - \{a_1\}$
 $= \{a_2, a_3, a_4, \dots, a_{k+1}\}$

$|A'| = k$

• # subsets of size 2 in A' is

$$\frac{(k)(k-1)}{2}$$

• add a_1 back in

now our set has cardinality $k+1$

• a_1 gets paired with every existing
element to create 2 subsets of size 2

total # subsets of size 2:

$$\frac{(k)(k-1)}{2} + k$$

IH

↳ everything a_i gets paired with

old subsets + new subsets

$$\frac{(k)(k-1)}{2} + k = \frac{(k)(k-1)}{2} + \frac{2k}{2}$$

$$= \frac{(k)(k-1) + 2k}{2}$$

$$= \frac{k^2 - k + 2k}{2}$$

$$= \frac{k^2 + k}{2}$$

$$= \frac{(k+1)(k)}{2} \quad \text{done!}$$

