

## CS1800 Day 19

### Admin:

- HW6 due today
- HW7 released today (due next Friday)
  - slightly shorter than most
    - more time to prep for exam2
    - will only count as 78% of other HWs with 100 points
- practice exam2 problems are out
  - slightly more induction examples
    - I want to make sure you have plenty to pull from in exam2

### Content:

- Induction with equalities and inequalities

## Induction (Weak):

Induction allows us to prove a never-ending sequence of statements:  $S(1)$ ,  $S(2)$ ,  $S(3)$ ,  $S(4)$ , ...

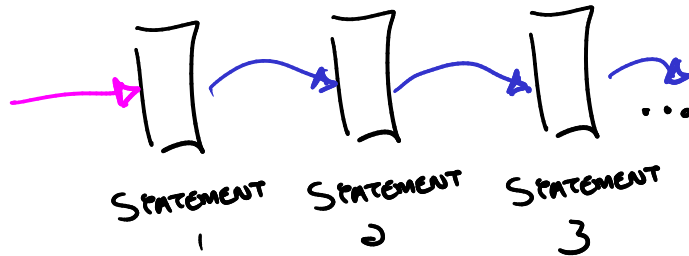
Process:

- Prove the first statement,  $S(n)$  for some  $n$
- Show that each statement implies the next statement:

Metaphor (Dominos):

To knock over all the dominos

- Push over the first one
- Place each other domino so that if the one behind it falls, it too will fall



EXAMPLE

# Summation Notation

$$1 + 2 + 4 + 8 + 16 + 32 + 64$$

$$= 2^0 + 2^1 + 2^2 + 2^3 + 2^4 + 2^5 + 2^6$$

$$= \sum_{k=0}^6 2^k$$

LAST VALUE OF K

STARTING VALUE OF K

"The sum of  $2^k$  where  $k$  goes from 0 to 6"

## In Class Activity: Summation Notation

Express each sum below in summation notation

$$9 + 11 + 13 + 15 + 17$$

$$9 + 2 \cdot 0 \quad 9 + 2 \cdot 2 \quad 9 + 2 \cdot 4$$

$$9 + 2 \cdot 1 \quad 9 + 2 \cdot 3$$

$$\sum_{k=0}^4 9 + 2k$$

$$1 + 2 + 3 + 4 + 5 + \dots + n$$

$$\sum_{k=1}^n k$$

Compute each sum below (the second one has a pattern and simplifies)

$$\sum_{k=10}^{12} 2k$$

$$= 2 \cdot 10 + 2 \cdot 11 + 2 \cdot 12$$

$$= 66$$

$$\sum_{k=0}^{101} (-1)^k = 0$$

$$\begin{array}{ccccccc} (-1)^0 & + & (-1)^1 & + & (-1)^2 & + & (-1)^3 & + & \dots \\ 1 & + & (-1) & + & 1 & + & (-1) & + & \dots \end{array}$$

## Algebraic Induction:

Show that the sum of the first  $n$  even integers is  $n(n+1)$

$$n=1$$

$$2 = 1(1+1)$$

$$n=2$$

$$2+4 = 2(2+1)$$

$$n=3$$

$$2+4+6 = 3(3+1)$$

## Algebraic Induction: Expressing our sum in summation notation

Show that the sum of the first  $n$  even integers is  $n(n+1)$

$$\begin{aligned} & 2 + 4 + 6 + 8 + \dots + 2n \\ = & 2 \cdot 1 + 2 \cdot 2 + 2 \cdot 3 + 2 \cdot 4 + \dots + 2n \end{aligned}$$

STATEMENT  $n$ :  $\sum_{k=1}^n 2k = n(n+1)$

## Induction Four Step Recipe: (AKA: how to not get turned around in a big induction proof)

1. Write out the statement for general  $n$
2. Specify the list of statements you're proving
  - your base case need not be  $n=1$
  - you're welcome to skip by more than 1 if you'd like  $S(3), S(5), S(7), S(9), \dots$
3. Prove the "Base case" (the smallest  $n$  for which your statement is true)
4. Prove the conditional: "If  $S(n)$  then  $S(n+1)$ "



PREP WORK  
(NOT NEEDED IN FINAL PROOF)

→ EXPRESS  $S(n+1)$

→ START @ MORE COMPLEX SIDE



## Algebraic Induction:

Show that the sum of the first  $n$  even integers is  $n(n+1)$

STATEMENT  $n$ :  $\sum_{k=1}^n 2k = n(n+1)$

BASE CASE ( $n=1$ )

$$\sum_{k=1}^1 2k = 2 \cdot 1 = 2$$

$$= 1(1+1) \\ = n(n+1)$$

$$S(n) \rightarrow S(n+1):$$

ASSUME STATEMENT  $n$  IS TRUE

$$\begin{aligned} \sum_{k=1}^{n+1} 2k &= 2(n+1) + \sum_{k=1}^n 2k \\ &= 2(n+1) + n(n+1) \end{aligned}$$

$$= (n+1)(n+2)$$

STATEMENT  $n$ :  $\sum_{k=1}^n 2k = n(n+1)$

$$S(n+1) \equiv \sum_{k=1}^{n+1} 2k = (n+1)(n+2)$$

NOTE TO SELF

- EXPRESS  $S(n+1)$
- NEXT SLIDE

# SUMMATION NOTATION IN INDUCTION PROBLEMS

$$\sum_{k=1}^{n+1} 2^k = 2 \cdot 1 + 2 \cdot 2 + 2 \cdot 3 + \dots + 2 \cdot n + 2 \cdot (n+1)$$
$$= 2 \cdot (n+1) + \sum_{k=1}^n 2^k$$

Induction proofs will ask us to prove  $S(n) \rightarrow S(n+1)$ .

For sums, it's helpful to be able to "pull out" the sum relevant in  $S(n)$  from the sum relevant for  $S(n+1)$  in this way.

In Class Activity:

$$2^0 + 2^1 + 2^2 + 2^3 + \dots + 2^{n-1} = \sum_{k=0}^{n-1} 2^k = 2^n - 1$$

Show that the sum of the first  $n$  powers of 2 is equal to  $2^n - 1$

$$n=1$$

$$2^0 = 1 = 2^1 - 1 = 2^n - 1$$

$$n=2$$

$$2^0 + 2^1 = 1 + 2 = 3$$

$$2^2 - 1 = 4 - 1 = 3$$

→ EXPRESS  $S(n+1)$

→ START @ MORE COMPLEX SIDE

In Class Activity:

$$2^0 + 2^1 + 2^2 + 2^3 + \dots + 2^{n-1} = \sum_{k=0}^{n-1} 2^k = 2^n - 1$$

Show that the sum of the first n powers of 2 is equal to

$$S(n) = \sum_{k=0}^{n-1} 2^k = 2^n - 1 \quad S(n+1) = \sum_{k=0}^n 2^k = 2^{n+1} - 1$$

INDUCTIVE STEP  $S(n) \rightarrow S(n+1)$

BASE CASE ( $n=1$ )

ASSUME  $\sum_{k=0}^{n-1} 2^k = 2^n - 1$

$$\sum_{k=0}^1 2^k = 2^0 = 1 = 2^1 - 1 = 2^n - 1$$

$$\begin{aligned} \sum_{k=0}^n 2^k &= \\ &= 2^n + \\ &= 2^n - 1 + 1 \end{aligned} \quad \begin{aligned} &+ \\ &+ \\ &+ \sum_{k=0}^{n-1} 2^k \end{aligned}$$

$$S(n) \equiv \sum_{k=0}^{n-1} 2^k = 2^n - 1$$

$$S(n+1) = \sum_{k=0}^n 2^k = 2^{n+1} - 1$$

INDUCTIVE STEP  $S(n) \rightarrow S(n+1)$

BASE CASE ( $n=1$ )

Assume  $S(n)$

$$\begin{aligned} \sum_{k=0}^{n-1} 2^k &= 2^0 = 1 = 2^1 - 1 \\ &= 2^n - 1 \end{aligned}$$

$$\begin{aligned} \sum_{k=0}^n 2^k &= 2^n + \sum_{k=0}^{n-1} 2^k \\ &= 2^n + 2^n - 1 \\ &= 2^{n+1} - 1 \end{aligned}$$

## Induction with inequalities: why? (preview a bit ahead, not necessary for exam2)

Suppose two algorithms both accomplish the same task but take a different number of "computes" to do so. For a list of size  $N$

Algorithm 1 takes:  $2^N$  computes

Algorithm 2 takes:  $N!$  computes

List Size	1	2	3	4	5	6	7
Algorithm 1 Computes	2	4	8	16	32	64	128
Algorithm 2 Computes	1	2	6	24	120	720	5040

*Note: In the original image, a hand-drawn arrow points to the column for List Size 4. In the 'Algorithm 1 Computes' row, the values 16, 32, 64, and 128 are circled. In the 'Algorithm 2 Computes' row, the values 1, 2, and 6 are circled.*

Goal: We want to show that algorithm 1 is faster for all lists sufficiently large ( $N$  is greater than some threshold)

## Induction with inequalities:

Prove that  $2^N < N!$  for all  $N$  above some threshold.

BASE CASE  $N=4$

$$2^N = 2^4 = 16 < 4! = N!$$

$$S(n) = "2^n < n!"$$

$$S(n+1) = 2^{n+1} < (n+1)!$$

INDUCTIVE STEP  $S(n) \rightarrow S(n+1)$

ASSUME  $2^n < n!$

$$(n+1)! = 1 \cdot 2 \cdot 3 \cdot 4 \cdot \dots \cdot n \cdot (n+1)$$

$$= n! \cdot (n+1)$$

$$> 2^n \cdot (n+1)$$

$$= 2^n \cdot 2 \cdot 2 \cdot \dots \cdot 2$$



## Algebra: Working with inequalities (1 of 3)

Move 1: add the same things to both sides, it preserves the inequality

$$\text{IF } 3 < 4 \quad \text{THEN} \quad 3 + 10 < 4 + 10$$

$$x < y \quad \rightarrow \quad x + c < y + c \quad \forall c \in \mathbb{R}$$

## Algebra: Working with inequalities (2 of 3)

Move 2: multiply by a positive value, it preserves the inequality

$$\text{IF } 3 < 4 \quad \text{THEN } 3 \cdot 10 < 4 \cdot 10$$

$$\text{IF } x < y \quad \text{THEN } xc < yc \quad \forall c \in \mathbb{R} \text{ with } c > 0$$

Move 3: multiply by a negative value, it swaps the inequality

$$\text{IF } 3 < 4 \quad \text{THEN } 3 \cdot -1 > 4 \cdot -1$$

$$\text{IF } x < y \quad \text{THEN } xc > yc \quad \forall c \in \mathbb{R} \text{ with } c < 0$$

## Algebra: Working with inequalities (3 of 3)

Move 4: sum two inequalities (large side together & small side together)

$$\begin{array}{l} \text{IF} \\ 3 < 4 \\ \text{AND} \\ 5 < 6 \\ \text{THEN} \\ 3 + 5 < 4 + 6 \end{array}$$

$$\begin{array}{l} \text{IF} \\ x < y \\ \text{AND} \\ w < z \\ \text{THEN} \\ x + w < y + z \end{array}$$

In Class Activity:

Show that  $n^2 > n + 10$  for all  $n$  above some threshold

BASE CASE



$n$	1	2	3	4	5	6
$n^2$	1	4	9	16	25	36
$n + 10$	11	12	13	14	15	16

In Class Activity:

STATEMENT  $N = N^2 > N+10$

Show that  $n^2 > n + 10$  for all  $n$  above some threshold

$$S(N+1) \equiv (N+1)^2 > N+11$$

BASE CASE  $N=4$

$$N^2 = 4^2 = 16 > 14 = 4+10 \\ = N+10$$

INDUCTIVE STEP  $S(n) \rightarrow S(n+1)$

ASSUME  $N^2 > N+10$

$$(N+1)^2 = N^2 + 2N + 1 \\ > N+10 + 2N + 1 \\ = N+11 + 2N \\ > N+11$$